Lattices

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Definition 1 (Upper bound)

If x and y belong to a poset P, then an upper bound of x and y is an element $z \in P$ satisfying $x \leq z$ and $y \leq z$.

Definition 2 (Least upper bound)

A least upper bound z of x and y is an upper bound such that every upper bound w of x and y satisfies $z \le w$.

Thus, if a least upper bound of x and y exists, then it is clearly unique due to antisymmetry of \leq . The element is denoted by $x \lor y$, read as "x join y" or "x sup y."

Dually, one can define a greatest lower bound of x and y.

Definition 3

A lattice is a poset for which every pair of elements has a least upper bound and greatest lower bound.

Let *L* be a lattice and $x, y \in L$.

One can verify that the following properties hold:

• The operations \land and \lor are associative, commutative and idempotent, that is, $x \land x = x \lor x = x$,

2
$$x \wedge (x \vee y) = x \vee (x \wedge y) = x$$
, and

$$3 \ x \wedge y = x \iff x \leq y \iff x \vee y = y.$$

In fact, one could even define a lattice axiomatically in terms of a set L with the operations \land and \lor satisfying the first two properties.

All finite lattices have $\hat{0}$ and $\hat{1}$. Indeed, if $L = \{x_1, \ldots, x_n\}$. Then $x_1 \land \cdots \land x_n$ and $x_1 \lor \cdots \lor x_n$ are well-defined elements of L and they are $\hat{0}$ and $\hat{1}$, respectively.

If *L* and *M* are lattices, then so are L^* , $L \times M$, $L \oplus M$. However, L + M will never be lattice unless $L = \emptyset$ or $M = \emptyset$. Indeed, if one takes $x \in L$ and $y \in M$, then there exists no meet of x and y in L + M. However, one can verify that $\widehat{L + M}$ is always a lattice.

Definition 4 (Meet semilattice)

If every pair of elements of a poset P has a meet, we say that P is a meet-semilattice.

Sometimes, it may be easy to check whether a finite poset is a meet-semilattice. The following proposition then helps us in determining whether the poset is also a lattice.

Proposition 1

Let P be a finite meet semilattice with $\hat{1}$. Then, P is a lattice.

Proof.

We just need to show that given $x, y \in P$, there exists a join of x and y. Towards this end, define $S := \{z \in P : x \le z, y \le z\}$. Then, S is finite as P is finite. Moreover, S is nonempty as $\hat{1} \in S$. Then, it can be seen that $x \lor y = \bigwedge_{z \in S} z$. The proof breaks for infinite posets as S defined earlier need not be finite and hence, its meet may not exist.

Analogously, one may define a join-semilattice and the corresponding proposition for a join-semilattice holds as well.

Definition 5

If every subset of L does have a meet and a join, then L is a called a complete lattice.

(The meet and join of a subset of a lattice have their natural meanings.) Clearly, a complete lattice has a $\hat{0}$ and $\hat{1}.$

Proposition 2

Let L be a finite lattice. The following are equivalent:

- L is graded and the rank generating function ρ of L satisfies $\rho(x) + \rho(y) \ge \rho(x \land y) + \rho(x \lor y)$ for all $x, y \in L$.
- **2** If x and y cover $x \land y$, then $x \lor y$ covers x and y.

We omit the proof.

A finite lattice satisfying either of the above (equivalent) properties is called a finite (upper) semimodular lattice.

A finite lattice L whose dual is semimodular is said to be lower semimodular.

A lattice which is both semimodular and lower semimodular is said to be modular.

Thus, a finite lattice is modular if and only if $\rho(x) + \rho(y) = \rho(x \land y) + \rho(x \lor y)$ for all $x, y \in L$.

This is the most important class of lattices from a combinatorial point of view.

Definition 6 (Distributive lattices)

A lattice L is said to be distributive if the following laws hold for all $x, y \in L$:

$$2 x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z).$$

Examples.

[*n*], B_n , D_n are distributive lattices. Π_n is not distributive for n > 2. Recall that the set of all order ideals of a poset P, denoted by J(P), ordered by inclusion, forms a poset.

Moreover, J(P) is closed under unions and intersections.

Thus, one can check that J(P) is a lattice as well with \land being \cap and \lor being \cup . Now, set theory tells us that J(P) is in fact, a distributive lattice as well.

The Fundamental Theorem of Finite Distributive Lattices (FTFDL) states that the converse is true when P is finite.

Theorem 1 (FTFDL)

Let L be a finite distributive lattice. Then, there is a unique (up to isomorphism) finite poset P for which $L \cong J(P)$.

The above theorem is also known as *Birkoff's Theorem*.

To prove this theorem, we first produce a candidate P and show that is indeed the case that $J(P) \cong L$. Towards this end, we define the following.

Definition 7 (Join-irreducible)

An element x of a lattice L is said to be join-irreducible if one cannot write $x = y \lor z$ with y < x and z < x.

Equivalently, the above condition says that if x is join-irreducible, then $x = y \lor z$ forces x = y or x = z.

Before carrying forward, we emphasise the following result from before as it will used often.

Proposition 3

Given an order ideal I of a finite poset P, there exists a corresponding antichain $A = \{x_1, \ldots, x_n\}$ where each x_i is a maximal element of I. We also write $I = \langle x_1, \ldots, x_n \rangle$. Moreover, I is the smallest order ideal containing A. It can also be verified that $\langle x_1, x_2 \rangle = \langle x_1 \rangle \cup \langle x_2 \rangle$. In fact, one has $\langle x_1, \ldots, x_n \rangle = \langle x_1 \rangle \cup \cdots \cup \langle x_n \rangle$. The following theorem will help us in coming up with a suitable candidate P and it will also help in showing the uniqueness of P claimed in FTFDL.

Theorem 2

An order ideal I of the finite poset P is join-irreducible in J(P) if and only if it is a principal order ideal of P.

Before giving a proof of this theorem, we observe that there's a natural one-to-one correspondence between principal order ideals of P and P. Namely, $\langle x \rangle \leftrightarrow x$. In fact, this correspondence is also an isomorphism as $\langle x \rangle \subset \langle y \rangle \iff x \leq y$.

Thus, if $J(P) \cong J(Q)$, then the set of join-irreducibles will also be isomorphic and in turn, $P \cong Q$. This shows us that the P mentioned in FTFDL, if it exists, is indeed unique.

Proof.

 (\implies) Suppose *I* is join-irreducible. Since *P* is finite, Proposition 3 tells us that there exists a corresponding antichain *A* that generates *I*. Let us assume that |A| > 1. Choose $a \in A$ and let $B := \{a\}$. Then, $\langle A \setminus B \rangle \cup \langle B \rangle = I$. By Proposition 3, $\langle A \setminus B \rangle \subsetneq I$ and $\langle B \rangle \subsetneq I$. However, this contradicts that *I* is join-irreducible. Thus, |A| = 1 and hence, *I* is principal.

(\Leftarrow) Suppose *I* is a principal order ideal. Then, there exists some $x \in P$ such that $\langle x \rangle = P$. Suppose $I = J \cup K$ for some $J, K \in J(P)$. Then $x \in J$ or $x \in K$. WLOG, we assume that $x \in J$. As *J* is an order ideal, we get that $\langle x \rangle \subset J$. But $\langle x \rangle = I$. Hence, we get that J = I. This proves that *J* is join-irreducible. What the theorem helps us with is the following -

Suppose that we are given an arbitrary (finite) poset Q and are told that $Q \cong J(P)$ for some poset P. The theorem has then shown that the poset P must be isomorphic to the set of the join-irreducibles of Q. (Or rather, the subposet obtained by inducing the structure of Q on the set of join-irreducibles of Q.)

In effect, it has given us a way of procuring an eligible candidate P to prove the Fundamental theorem that we wanted to prove.

Before proving the theorem, we shall prove the following lemma. Let L be a finite distributive lattice and let P be the set of all the join-irreducible elements of L.

Lemma 1

For $y \in L$, there exist $y_1, y_2, \ldots, y_n \in P$ such that $y = y_1 \vee y_2 \vee \cdots \vee y_n$. For n minimal, the expression is unique up to permutations.

Proof (Of existence).

If y is join-irreducible, then we are done.

Suppose $y \notin P$. Then, by definition, there exist $y_1, y_2 \in L$ such that $y = y_1 \lor y_2$ with $y_1 < y$ and $y_2 < y$. If one of y_1 or y_2 is not in P, then we can further "decompose" it. As L is finite and we keep getting smaller elements, this process must stop after a finite number of steps. Thus, given any $y \in L$, there does exist a representation as stated.

Proof (of uniqueness).

Now, suppose $y_1 \vee y_2 \vee \cdots \vee y_n = y = z_1 \vee \cdots \vee z_n$ for $y_i, z_i \in P$ for each $i \in [n]$ where n is the least number of elements of P required to be "joined" to get y. Note that given any $i \in [n], z_i \leq y$.

Thus,
$$z_i = z_i \land y = z_i \land \left(\bigvee_{j=1}^n y_j\right) = \bigvee_{j=1}^n (z_i \land y_j)$$
, by distributivity.

But $z_i \in P$ and thus, we get that $z_i = z_i \wedge y_j$ for some $j \in [n]$. This gives us that $z_i \leq y_j$.

Now, suppose it is the case that there exists $k \in [n]$ such that $k \neq i$ and $z_k \leq y_j$. We show that this leads to a contradiction. Since \lor is associative and commutative, we can assume that i = 1 and k = 2. As $y_j \leq y$, we get that $y_j \lor z_3 \lor \cdots z_n = y$, contradicting the minimality of n.

Thus, given any $i \in [n]$, there exists a unique $j \in [n]$ such that $z_i \leq y_j$. Similarly, we get an inequality in the other direction which proves the lemma.

Carrying on with the same notation, we define the following functions:

$$f: J(P) \to L$$
$$f(I) := \bigvee_{x \in I} x.$$

$$g: L o J(P)$$

 $g(y) := \bigcup_{i=1}^{n} \langle y_i \rangle,$

where $y = y_1 \lor \cdots \lor y_n$ in the unique way as described earlier. By our previous lemma and commutativity of union, we get that g is indeed well defined. By our previous lemma, it is also clear that f is surjective.

Proof of FTFDL

Now, we claim that g(f(I)) = I for every $I \in J(P)$. To see this, let A be the antichain corresponding to I. That is, let $A = \{y_1, \ldots, y_n\}$ where each y_i is a maximal element of I. Then, we get that $f(I) = \bigvee_{x \in I} x = y_1 \lor \cdots \lor y_n$, using the fact that each $a \lor b = b$ if $a \le b$. Now, $g(f(I)) = g(y_1 \lor \cdots \lor y_n) = \langle y_1 \rangle \cup \cdots \cup \langle y_n \rangle = \langle y_1, \ldots, y_n \rangle = I$. (*)

Thus, $g \circ f$ is the identity function on J(P) and hence, f is injective. This shows that f is bijective. As g is its one-sided inverse, it is also its two-sided inverse since f is a bijection. Now, we show that f is an isomorphism by showing that both f and g are order preserving.

Proof of FTFDL

1. *f* is order preserving. Suppose $I_1, I_2 \in J(P)$ with $I_1 \subsetneq I_2$. Then,

$$f(I_2) = \bigvee_{x \in I_2} x$$
$$= \left(\bigvee_{x \in I_1} x\right) \lor \left(\bigvee_{x \in I_2 \setminus I_1} x\right)$$
$$\ge \bigvee_{x \in I_1} x$$
$$= f(I_1)$$

As we already have seen that f is injective, we get that $f(I_1) < f(I_2)$, as desired.

Proof of FTFDL

2. g is order preserving. Suppose that for $I_1, I_2 \in J(P)$, we have $\bigvee x \leq \bigvee x$. We want to show that $I_1 \subset I_2$. $x \in I_1$ $x \in I_2$ Let $x \in I_1$ be given. We have that $x \leq \bigvee x$ and thus, $x \leq \bigvee y$. $x \in I_1$ $\implies x = \left(\bigvee_{y \in h} y\right) \land x$ $=\bigvee(y\wedge x)$ (By distributivity) $v \in b$

As x is join-irreducible, there exists some $y \in I_2$ such that $x = y \wedge x$, that is, $x \leq y$. As I_2 is an order ideal, this implies that $x \in I_2$. Thus, $I_1 \subset I_2$. The line marked (*) has a possible flaw. We are assuming that $g(y_1 \vee \cdots \vee y_n) = \langle y_1 \rangle \cup \cdots \cup \langle y_n \rangle$, that is, we are assuming that $y_1 \vee \cdots \vee y_n$ is indeed the minimal representation of f(I).

However, this is justified for if there were a shorter representation in terms of join-irreducibles, we would get a contradiction about the maximality of y_i s.