

$$\int (\overset{\circ}{\frown} 5 \overset{\circ}{\frown}) dx$$

MA 839

Advanced Commutative Algebra

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A Quick Intro.

Setup: A ring is commutative with 1.

Let M be an R -module.

Observation: ① If M is cyclic, (say $M = \langle \alpha \rangle = \{a\alpha : a \in R\}$),
we get an R -linear map $R \rightarrow M$ which is onto.
 $a \mapsto a\alpha$

Then, $M \cong R/I$ where I is the kernel.
In this case, $I = \text{ann}_R(\alpha)$.

Thus, if M is cyclic, then M is a quotient of R .

② Suppose $\exists x, y \in M$ s.t. $M = \langle x, y \rangle = \{ax + by \mid a, b \in R\}$.
 $= \{ax + by \mid (a, b) \in R^{\oplus 2}\}$

Then, we get an onto R -linear map $R \oplus R \xrightarrow{\varphi} M$
 $e_1 \mapsto x$
 $e_2 \mapsto y$ } extend this
 $\{e_1, e_2\}$ is a basis
→ this lets us extend the map

In particular, $M \cong R^2 / \ker \varphi$.

Q. Is it necessary that we can actually write

$$M \cong \frac{R}{\langle \rangle} \oplus \frac{R}{\langle \rangle} ?$$

This has a positive answer: ① R is a field
② R is a PID

CAUTION: We won't include fields as PID.
That is, when we say "PID", we exclude fields **||**

③ Suppose M is a finitely generated (f.g.) R -module.

(That is, suppose $M = \langle \alpha_1, \dots, \alpha_n \rangle$.)

Then, M is a quotient of $R^{\oplus n}$.
→ \mathbb{A}^n

Then, M is a quotient of R^n .

way to get this

Define $R^n \xrightarrow{\varphi} M$ by $e_i \mapsto x_i$.

$$M \cong R / \ker \varphi.$$

④ In general, consider a free module with " M as basis", call it $F(M)$. Then $F(M)$ maps onto M .

Slightly more general: If $A \subset M$ is a generating set, i.e., $M = \langle A \rangle$,

then $F(A)$ maps onto M .

Thus, M can be written as a quotient of a free-module.

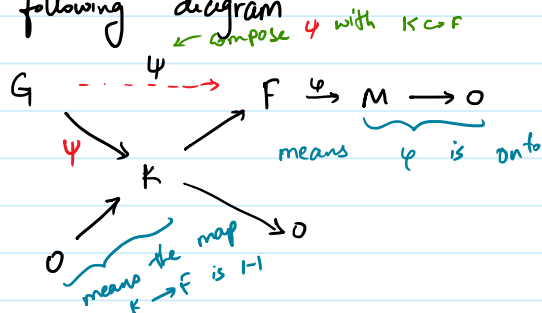
To Summarise: If M is an R -module, then M can be written as a quotient of a free R -module. Moreover, if M is f.g., then the free module can be assumed to have finite rank.

Free resolution of M (over R):

Let F be a free R -module mapping onto M with kernel K . That is, $\varphi: F \rightarrow M$ is onto R -linear and $\ker \varphi = K$.

Now, \exists a free R -module G and an onto map $\psi: G \rightarrow K$

We capture this in the following diagram



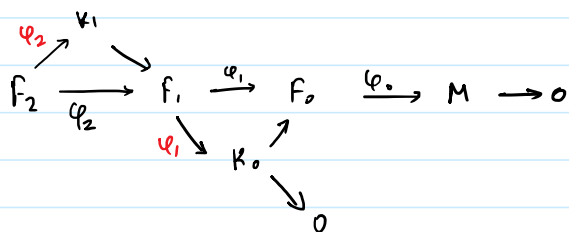
Note that $\text{im } \psi = K = \ker \varphi$.

Thus, we have $G \xrightarrow{\psi} F \xrightarrow{\varphi} M \rightarrow 0$.

- ① φ is onto and $\ker \varphi = \text{im } \psi$.
- ② G and F are free R -modules.

Note that we can repeat the above process with K instead of F .

Change notation: $F_0 := F$, $F_1 := G$, $K_0 := K$, $\varphi_0 := \varphi$, $\varphi_1 := \psi$.



Thus, we get free modules $\{F_n, \varphi_n: F_n \rightarrow F_{n-1}\}$ such that $\ker \varphi_{n-1} = \text{im } \varphi_n$ written as

$$\dots \rightarrow F_n \xrightarrow{\varphi_n} F_{n-1} \xrightarrow{\varphi_{n-1}} \dots \rightarrow F_1 \xrightarrow{\varphi_1} F_0 \xrightarrow{\varphi_0} M \rightarrow 0$$

F_i s are free, φ_0 is onto & $\ker \varphi_{n-1} = \text{im } \varphi_n$, $n \geq 1$

Often, we drop the 'n' and call

$$F: \dots \rightarrow F_n \xrightarrow{\varphi_n} F_{n-1} \rightarrow \dots \rightarrow F_1 \xrightarrow{\varphi_1} F_0 \rightarrow 0 \text{ as an}$$

R free resolution of M .

im $\varphi_1 = K$, this is not exact here. φ_1 not onto (rec.)

Q: ① If M is f.g.:

Can we get F_i s so that $\text{rank}(F_i) < \infty \forall i$.

② If yes, are $\text{rank}(F_i)$ s independent of construction?

③ Can you describe the maps?

④ Give explicit bases for F_i s s.t. the maps are "described nicely"?

Q: If two modules have "isomorphic" free resolutions, are they isomorphic?

$$\begin{array}{ccccccc} \dots & \xrightarrow{\varphi_3} & F_2 & \xrightarrow{\varphi_2} & F_1 & \xrightarrow{\varphi_1} & F_0 \rightarrow 0 \\ & & \downarrow \varphi_2' & & \downarrow \varphi_1' & & \downarrow \varphi_0' \\ \dots & \xrightarrow{\varphi_3'} & F_2' & \xrightarrow{\varphi_2'} & F_1' & \xrightarrow{\varphi_1'} & F_0' \rightarrow 0 \end{array} \quad \begin{array}{l} M \cong F_0 / \text{im } \varphi_1 \\ M' \cong F_0' / \text{im } \varphi_1' \end{array}$$

$$\varphi_1' \varphi_1 = \varphi_0' \varphi_1 \quad (*)$$

Claim. $\gamma_0(\text{im } \varphi_1) = \text{im } \varphi_1'$

(\Leftarrow) clear by (*)

(\Rightarrow) clear again since $\varphi_1' = \gamma_0 \varphi_1 \gamma_1^{-1}$

$$\text{Thus, } M_0 \cong \frac{F_0}{\text{im } \varphi} \cong \frac{\varphi_0(F_0)}{\varphi_0(\text{im } \varphi)} = \frac{F_0'}{\text{im } \varphi_1'} \cong M_0'$$

Lecture 1 (11-01-2021)

11 January 2021 11:07

Free modules: (Free modules)

As usual : R is a (commutative) ring (with 1).
 M is an R -module.

Def^m. ① Let $A \subset M$. A is said to be a **generating set** of M (as an R -module) if
 $\forall x \in M, \exists x_1, \dots, x_n \in A$ and $(a_1, \dots, a_n) \in R^n$ s.t.
 $x = a_1 x_1 + \dots + a_n x_n$.
(Note that A need not be finite.)

Notation : $M = \langle A \rangle$

If $A = \{x_1, \dots, x_n\}$ is finite, then $M = \langle x_1, \dots, x_n \rangle$
and M is said to be **finitely generated**.

②ⓐ Let $x_1, \dots, x_n \in M$. We say $\{x_1, \dots, x_n\}$ is **linearly independent** (over R) if for $(a_1, \dots, a_n) \in R^n$,

$$a_1 x_1 + \dots + a_n x_n = 0 \Rightarrow (a_1, \dots, a_n) = 0 \text{ in } R^n.$$

ⓑ A subset $A \subset M$ is **linearly independent** (over R) if every finite subset of A is linearly independent (over R).

③ A subset $A \subset M$ is a **basis** of M (over R) if $M = \langle A \rangle$ (over R) and A is linearly independent (over R).

④ M is **free** (over R) if M has a basis (over R).

REMARKS

- ① Not every R -module has a basis.
- ② A minimal generating set need not be lin. indep.
- ③ A maximal lin indep. set need not be a gen. set.

Q. If every R -module has a basis, is R a field?

(Yes. Take a non-field ring R and any non-trivial ideal $I \neq R$.
Then, R/I has no lin. indep. set over R .)

Q. If an R -module M has a basis, does every basis have the same cardinality?

Ans. Yes. This is called the Invariant Basis Number (IBN) property of R .

Remark. This is not true if R is non-commutative. (That is, we can find a counterexample of a non-commutative ring.)
If R is a division ring, then again we have IBN.

Defn

If M has a finite basis, say B , then we define

$$\text{rank}(M) := |B|.$$

If M is free with an infinite basis, $\text{rank}(M) := \infty$.

(Rank)
(When we do say "rank", we will usually mean "finite rank".)

EXAMPLES. ① $R^{(n)}$ is a free R -module of rank n
 $M_{m \times n}(R)$ of rank mn
 $R[x]$ of rank ∞

② Let A be a non-empty set and

$$F_0(A, R) = \{f: A \rightarrow R \mid f(a) = 0 \text{ for all but fin. many } a \in A\}.$$

Then, $F_0(A, R)$ is an R -module under pointwise operations.

In fact, $F_0(A, R)$ is a free R -module with basis $\{\chi_a\}_{a \in A}$,
where

$$\chi_a(b) = \begin{cases} 0 & ; b \neq a \\ 1 & ; b = a \end{cases}$$

To see where the above set is generating, given any $f \in F_0(A, R)$, we can write

$$f = \sum_{a \in A} f(a) \chi_a.$$

↑ the sum is actually finite since $f(a) = 0$ for all but finitely many a .
(it is to be understood that 0s are ignored.)

Q. What if we take $F(A, R)$? (All functions.)

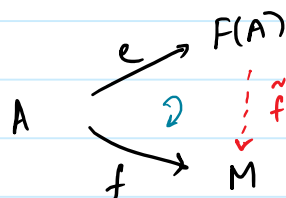
Universal Property of free modules: (Free R-module on A)

Def'n:

Given a non-empty set A , a free R-module on A is a pair $(F(A), e)$ where (i) $F(A)$ is an R-module, (ii) $e: A \rightarrow F(A)$ is a (set) function

satisfying:

Given an R-module M and a function $f: A \rightarrow M$, there exists a unique R-linear $\tilde{f}: F(A) \rightarrow M$ making the following diagram commute.



(That is, $\tilde{f}e = f$.)

REMARKS. ① Given $A = \emptyset$, a free R-module on A exists, and is unique up to isomorphism.

Moreover, $e: A \rightarrow F(A)$ is one-one and $F(A)$ is free with basis $\{e_a\}_{a \in A}$, where $e_a := e(a)$.

② If M is a free R -module, then $M \cong F(B)$, where B is any basis of M .

Thus, an R -module M is free iff $M \cong F(A)$ for some A .

What the universal property is really saying is that:
given a free R -module M with basis A , every R -linear
 $M \rightarrow N \rightarrow R\text{-module}$

is completely determined by its action on A .

(The above is in the sense that given any assignment of values on A , we do get an R -linear map.)

EXAMPLE: Given an R -module M , such that $M = \langle A \rangle$, we can write M as a quotient of $F(A)$.
(what we did last lec.)

Lecture 2 (12-01-2021)

12 January 2021 08:35

Weyl Algebra

Ex. k is a field, $k[x_1, \dots, x_d]$
 $\partial_1, \dots, \partial_d \rightarrow$ partial diff op.
 $Ad(k) = k[x_1, \dots, x_d, \partial_1, \dots, \partial_d]$ D -modules
 ↑ non-comm. How would you define products?

Tensor Product

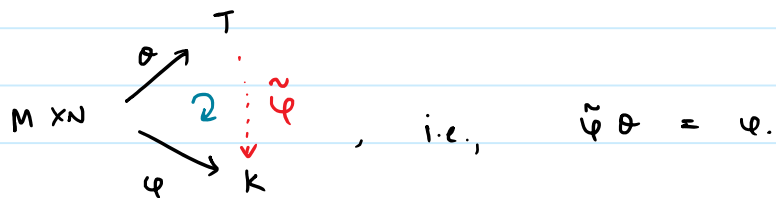
(Tensor product)

Tensor product (of two modules) essentially converts the study of bilinear maps to linear maps.

Defn.

Given R -modules M and N , the tensor product of M and N (over R) is a pair (T, θ) , where T is an R -module, $\theta: M \times N \rightarrow T$ is R -bilinear satisfying:

Given (K, φ) where K is an R module, $\varphi: M \times N \rightarrow K$ is R -bilinear, there exists a unique R -linear map $\tilde{\varphi}: T \rightarrow K$ making the following diagram commute



(We are using "with", but can use "and" and we prove $M \otimes_R N \cong N \otimes_R M$.)

Thm. A tensor of M with N exists and is unique, up to isomorphism.

Uniqueness follows by universal property.

Notation: $M \otimes_R N$

Construction:

Want $M \times N \xrightarrow{\varphi} T$ $\theta(x_1 + x_2, y) = \theta(x_1, y) + \theta(x_2, y)$

Step 1: Let $F = F(M \times N)$, the free module on the set $M \times N$.

We get a map $e: M \times N \rightarrow F(M, N)$
 $(x, y) \mapsto e(x, y)$

$\{e(x, y) : x \in M, y \in N\}$ is a basis for F .

Let G be the submodule of F generated by

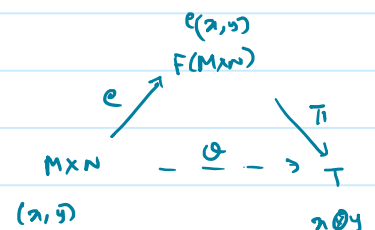
- $e(x_1 + x_2, y) - e(x_1, y) - e(x_2, y)$
- $e(x, y_1 + y_2) - e(x, y_1) - e(x, y_2)$
- $e(ax, y) - a e(x, y)$
- $e(x, ay) - a e(x, y)$

$\forall x, x_1, x_2 \in M, \forall y, y_1, y_2 \in N, \forall a \in R$

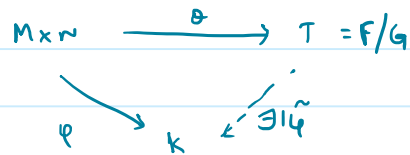
Step 2: Define $T = F/G$. Let $\pi: F \rightarrow T$ be the natural map.
 Set $\pi(e(x, y)) =: x \otimes y$.

Note that $\left. \begin{aligned} (x_1 + x_2) \otimes y &= x_1 \otimes y + x_2 \otimes y \\ x \otimes (y_1 + y_2) &= x \otimes y_1 + x \otimes y_2 \\ (ax) \otimes y &= a(x \otimes y) = x \otimes (ay) \end{aligned} \right\} \begin{aligned} \forall x, \dots \in M \\ \forall y, \dots \in N \\ \forall a \in R \end{aligned}$

Consider $\theta = \pi e: M \times N \rightarrow T$
 $(x, y) \mapsto x \otimes y$

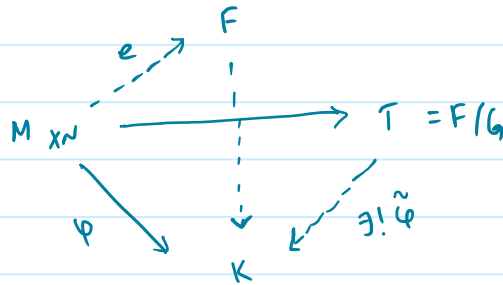


Step 3. Now, suppose we are given a bilinear $\varphi: M \times N \rightarrow K$. (K is some R -module.)

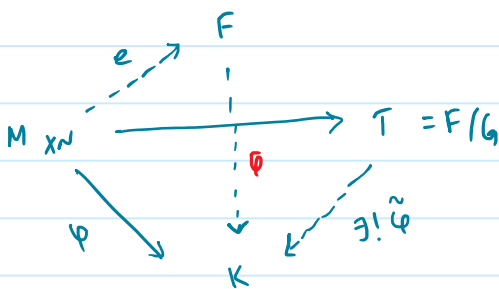


- Q1: Is there even a func. $\tilde{\varphi}$?
 Q2: Is it R -linear?
 Q3: Is it unique?

Note also



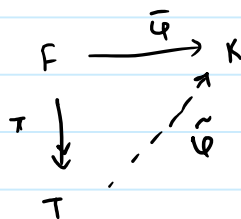
Note we have a set map $M \times N \xrightarrow{\theta} T = F/G$ which induces an R linear map $\bar{\varphi} : F \rightarrow K$. (UMP of free modules)



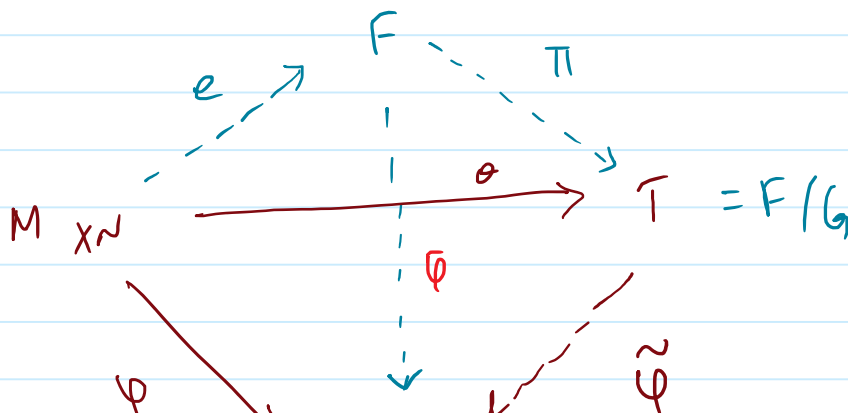
We now want to show that $\bar{\varphi}$ factors through T . It would suffice to show that $G \in \ker \bar{\varphi}$.

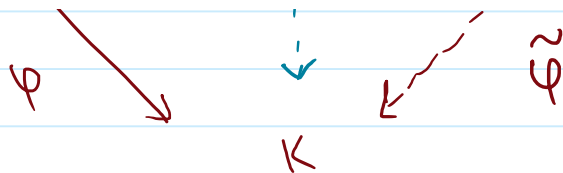
Using bilinearity of φ , it follows that all our (four types of) generators of G are in $\ker \bar{\varphi}$.

Thus, $\bar{\varphi}$ factors through quotient. That is, $\exists! R$ -linear $\tilde{\varphi} : T \rightarrow K$ s.t.



commutes. That is, $\bar{\varphi} = \tilde{\varphi} \pi$.





Can now verify $\tilde{\varphi} \circ \theta = \varphi$. (Use commutation of diff. triangles.)
 Can also verify that $\tilde{\varphi}$ is unique R -linear such.

Basic Properties:

- (1) ["Identity"] $R \otimes_R M \cong M$
- (2) [Commutativity] $M \otimes_R N \cong N \otimes_R M$
- (3) [Associativity] $(M \otimes_R N) \otimes_R L \cong M \otimes_R (N \otimes_R L)$
- (4) [Distributivity] $M \otimes_R (N \oplus L) \cong (M \otimes_R N) \oplus (M \otimes_R L)$

Q. How does get an R -linear map $M \otimes_R N \rightarrow L$?

A. Give an R -bilinear map $M \times N \rightarrow L$.

Pretty much the only way. $x \otimes y$ could be 0 even if $x, y \neq 0$.
 Thus, checking "well-defined"ness would become quite difficult.

Q. Let M be an R -module. $R \subset S$ subring.

Can you identify a natural S -module on $S \otimes_R M$?
 (Base change)

Lecture 3 (14-01-2021)

14 January 2021 09:31

Distributivity: Given R -modules L, M, N

$$L \otimes_R (M \oplus N) \cong (L \otimes_R M) \oplus (L \otimes_R N)$$

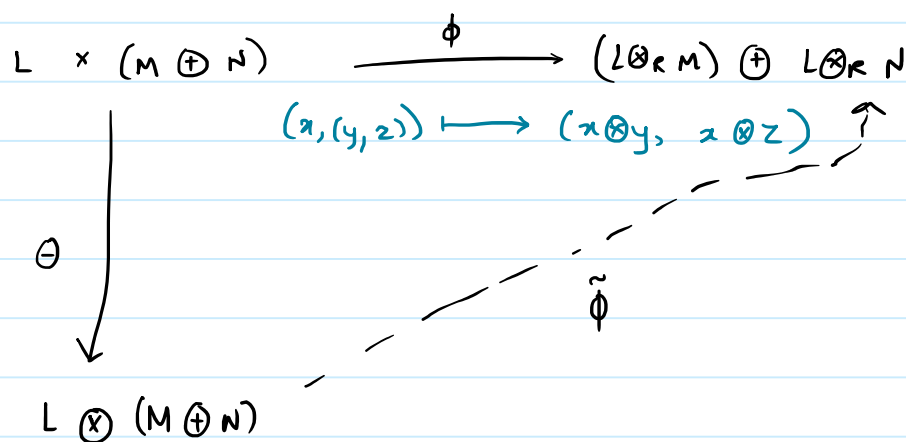
How do we show? Want something like:

$$x \otimes (y, z) \mapsto (x \otimes y, x \otimes z)$$

Note that elements of this form only GENERATE the tensor.

We now need to show the above map is well-defined. (as a function.)
To do so, we go back to $L \times (M \oplus N)$ and use the universal property.

This ϕ is well defined.
Every elt. here is uniquely written in the given form.



Note that ϕ is R -bilinear, thus an R -linear map $\tilde{\phi}$ (as indicated) which makes the diagram commute does exist.

To now show that is an isomorphism, we construct an inverse

$$\begin{aligned} \psi: (L \otimes_R M) \oplus (L \otimes_R N) &\longrightarrow L \otimes_R (M \oplus N) \\ (x \otimes y, 0) &\longmapsto x \otimes (y, 0) \\ (0, x \otimes z) &\longmapsto x \otimes (0, z) \end{aligned}$$

verify \rightarrow
these are \rightarrow
well defined
(again universal property)

Can verify now that ψ is the two-sided inverse of $\tilde{\phi}$.

inverse of $\tilde{\phi}$.

REMARK. Suppose M and N are R -modules.

① If $x = 0 \in M$, then $x \otimes y = 0 \quad \forall y \in N$.

However if $x \otimes y = 0$ for some $x \in M, y \in N$, we cannot conclude $x = 0$ or $y = 0$.

Example: Take $\mathbb{Z}/3\mathbb{Z} \otimes_{\mathbb{Z}} \mathbb{Z}$.

Thus $x \otimes y = 0 \quad \forall y \in N \not\Rightarrow x = 0$.
In fact, look at $\mathbb{Z}/3\mathbb{Z} \otimes_{\mathbb{Z}} \mathbb{Q} \leftarrow$ this is the zero module.

② $M \otimes_R N$ is generated by $\{x \otimes y : x \in M, y \in N\}$ as R -module.
In particular, if M and N are f.g., then so is $M \otimes_R N$.

Take fin. gen. sets S_M and S_N . Then

$$M \otimes_R N = \langle x \otimes y : x \in S_M, y \in S_N \rangle$$

③ If M and N are free, then so is $M \otimes_R N$. Identify a basis.

For finite rank: write

$$M = \underbrace{R \oplus \dots \oplus R}_m$$
$$N = \underbrace{R \oplus \dots \oplus R}_n$$

Then, $M \otimes_R N = (R \oplus \dots \oplus R) \otimes_R (R \oplus \dots \oplus R)$
 \searrow distribute and use $R \otimes_R R \cong R$.

④ It is possible that $M \neq 0 \neq N$ but $M \otimes_R N = 0$.
(See 1.)

Q. Given a simple tensor $x \otimes y$, how can we determine if it's 0?
Concrete ex: Is $2 \otimes 3 \in \mathbb{Z} \otimes_{\mathbb{Z}} \mathbb{Z}$ non-zero?

Q. Is it possible that $M \otimes_R M = 0$ even if $M \neq 0$?

Yes. Take $R = \mathbb{Z}$ and $M = \mathbb{Q}/\mathbb{Z}$.

$$1 \otimes 1 = \frac{1}{2} \otimes 2 = \frac{1}{2} \otimes 0 = 0$$

Yes. Take $R = \mathbb{Z}$ and $M = \mathbb{Q}/\mathbb{Z}$.

$$\left(\frac{a}{b} + \mathbb{Z}\right) \otimes \left(\frac{c}{d} + \mathbb{Z}\right) = \left(\frac{a}{db} + \mathbb{Z}\right) \otimes \underbrace{\left(c + \mathbb{Z}\right)}_0 = 0.$$

Tensor Algebra (Tensor Algebra)

The tensor algebra of M

$$R \oplus M \oplus T_2(M) \oplus T_3(M) \oplus \dots$$

$$T(M) = R \oplus M \oplus (M \otimes M) \oplus (M \otimes M \otimes M) \oplus \dots$$

$T(M)$ is clearly an additive group. One can define multiplication by "concatenation."

$$\underbrace{(x_1 \otimes \dots \otimes x_m)}_{m^{\text{th}} \text{ piece}} \cdot \underbrace{(y_1 \otimes \dots \otimes y_n)}_{n^{\text{th}} \text{ piece}} = x_1 \otimes \dots \otimes x_m \otimes y_1 \otimes \dots \otimes y_n.$$

Elements of $T(M)$ are written as formal sums: $z_0 + z_1 + \dots + z_n$
 Identify $T(R^{\otimes n})$.
 Some quotients of $T(M)$:
 $\underbrace{z_0}_R + \underbrace{z_1}_M + \dots + \underbrace{z_n}_{M^{\otimes n}}$
 $T_n(M)$

① Symmetric algebra (Symmetric Algebra)

Given $x, y \in M$, $x \otimes y \neq y \otimes x$.

(\neq^* : not necessarily equal)

$$\text{Define } \text{Sym}(M) = \frac{T(M)}{\langle x \otimes y - y \otimes x \mid x, y \in M \rangle} = R \oplus M \oplus \text{Sym}_2(M) \oplus \dots$$

$\text{Sym}(M)$ is now a commutative algebra.

② Exterior algebra [Wedge (M)]

(Exterior algebra)

$$\Lambda(M) = \frac{T(M)}{\langle x \otimes y + y \otimes x \rangle}$$

R and M
not affected.
only $M^{\otimes 2}$
onwards.

$$\langle x \otimes y + y \otimes x \rangle$$

$$\parallel$$

$$R \oplus M \oplus \Lambda^2(M) \oplus \dots$$

Q. What are $\text{Sym}(R^{\oplus n})$ and $\Lambda(R^{\oplus n})$?

(Trivial extension)

$$\textcircled{3} R \ltimes M = \frac{T(M)}{\langle x \otimes y \mid x, y \in M \rangle}$$

(Trivial extension or idealisation.)

In this algebra, M is an ideal, with $M^2 = 0$.
This is called an idealisation of M .

Q. What is $R \ltimes R^{\oplus n}$?

Lecture 4 (18-01-2021)

18 January 2021 10:33

Base change:

(Base change or extension of scalars)

Let R and S be rings and $\varphi: R \rightarrow S$ a ring homomorphism. Then, we say that S is an R -algebra "via φ ", i.e., S has an R -module structure defined by

$$a \cdot x = \varphi(a)x \quad \forall a \in R \quad \forall x \in S.$$

Two key examples:

- ① R is a subring of S (φ is $1 \mapsto 1$)
- ② S is a quotient of R (φ is onto)

(and their compositions)

Example. If $I \subset R[x_1, \dots, x_d]$ is an ideal, then

$$S = \frac{R[x_1, \dots, x_d]}{I} \text{ is an } R\text{-algebra}$$

(via the natural maps $R \hookrightarrow R[x_1, \dots, x_d] \twoheadrightarrow S$)

A consequence of the Hilbert Basis Theorem:

Thm. Every finitely generated algebra over a Noetherian ring is a Noetherian ring.

(Not saying Noetherian as an R -module! Hilbert's Basis Thm does not give that!)

Note. If S is an R -algebra (via φ) and M is an S -module, then M has a natural R -module structure (via φ).

Q. What about the reverse? If M is an R -module, is there a "natural" way to induce an S -module structure on it?
 M , in general. Take \mathbb{Z} as a \mathbb{Z} module and $\varphi: \mathbb{Z} \hookrightarrow \mathbb{Q}$.
Is there any "good" \mathbb{Q} -module structure on \mathbb{Z} ?

Example. Let M be an R -module and $I \subset R$ an ideal, $A \subset R$ m.c.s.

- ① When is M an R/I -module? } induced multiplication
 ② When is M an $R_{\mathfrak{p}}$ -module?

In general, we can create modules over R/I and $R_{\mathfrak{p}}$ starting from M : They are M/IM and $M_{\mathfrak{p}}$, respectively.

Obs. $M/IM \cong R/I \otimes_R M$ and $M_{\mathfrak{p}} \cong R_{\mathfrak{p}} \otimes_R M$
 (Isomorphic as R -modules.)

Base change (or extension of scalars):

Let S be an R -algebra (via φ) and M be an R -module.
 The R -module

$$S \otimes_R M$$

has a natural S -module structure defined by

$$a (b \otimes x) := (ab) \otimes x. \quad \forall a, b \in S \quad \forall x \in M$$

(Note that the above is only being defined for simple tensors.)

Note: ① If $M = \langle x_1, \dots, x_n \rangle$ over R , then

$$S \otimes_R M = S \langle 1 \otimes x_i : 1 \leq i \leq n \rangle.$$

② If $M \cong R^{\oplus n}$, $S \otimes_R M \cong S^{\oplus n}$ as R -modules.

\otimes distributes over \oplus

In fact, $S \otimes_R M \cong S^{\oplus n}$ as S -modules as well.

Q. Given an R -linear map $\psi: M \rightarrow N$, will this induce an S -linear map $\bar{\psi}: S \otimes_R M \rightarrow S \otimes_R N$?

Does this help in the above?

③ If $M = R[x]$, then $S \otimes_R M \cong S[x]$ as S -modules.

④ If M is a free R -module, then $S \otimes_R M$ is a free S -module.

Example of base change: "Mod p test for irreducibility of a polynomial in $\mathbb{Z}[x]$ "

Recall the test: Given $f = a_0 + a_1x + \dots + a_nx^n \in \mathbb{Z}[x]$

If we can find a prime p s.t. $f \pmod p$ is irred in $(\mathbb{Z}/p\mathbb{Z})[x]$, then f is irred.*

(*Need to take care of degree not falling.)

This is an example of base change with $R = \mathbb{Z}$ and $S = \mathbb{Z}/p\mathbb{Z}$.

Complexes and Homology

(Complexes and Homology)

Example: Construct a free resolution of $\mathbb{Z}/2\mathbb{Z}$ over \mathbb{Z} .

$$\begin{array}{ccccccc} \dots & \rightarrow & 0 & \rightarrow & 0 & \rightarrow & \mathbb{Z} & \xrightarrow{\quad} & \mathbb{Z} & \rightarrow & \mathbb{Z}/2\mathbb{Z} & \rightarrow & 0 \\ & & & & & & & & & & \downarrow \tau & & \\ & & & & & & & & & & \mathbb{Z}/2\mathbb{Z} & & \\ & & & & & & & & & & \downarrow \tau & & \\ & & & & & & & & & & 0 & & \end{array}$$

Diagram details: A sequence of maps $\mathbb{Z} \xrightarrow{\pi} \mathbb{Z}/2\mathbb{Z}$ with $\ker \pi = 2\mathbb{Z}$. A red arrow labeled τ points from $\mathbb{Z}/2\mathbb{Z}$ to the next $\mathbb{Z}/2\mathbb{Z}$ in the sequence.

Q1. What is a generating set for $\mathbb{Z}/2\mathbb{Z}$ over \mathbb{Z} ?

Ans. $\{\bar{1}\}$ ← singleton.

Thus, we map one copy of \mathbb{Z} onto $\mathbb{Z}/2\mathbb{Z}$.
That is,

$$\begin{array}{ccc} \mathbb{Z} & \rightarrow & \mathbb{Z}/2\mathbb{Z} \\ | & \mapsto & \bar{1} \end{array}$$

Q2. What is \ker ?

Ans. It is $2\mathbb{Z}$.

Q3. What can map onto $2\mathbb{Z}$?

Q3. What can map onto $2\mathbb{Z}$?

Ans. \mathbb{Z} . $x \mapsto 2x$

Q4. What is \ker ?

If is 0. The map is 1-1. This gives the diagram

Note that we could have also written

$$0 \rightarrow 2\mathbb{Z} \rightarrow \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$$

since $2\mathbb{Z}$ itself is a free \mathbb{Z} -module. The reason we did not do this is because we are sticking to writing free R -modules as copies of R .

Q. Let R be a ring and $a \in R$. Is the following a free resolution of $R/\langle a \rangle$ over R ?

$$0 \rightarrow R \xrightarrow{a} R \rightarrow R/\langle a \rangle \rightarrow 0$$

$z \mapsto az$

(finite free resolution)

Called a finite free resolution.

Obs. • Note that if some \ker is free, we can stop there. \uparrow
(That is basically saying that $R^{\oplus n} \rightarrow \ker$ will be 1-1.)

• The above does happen if R is a PID and M is fg. So, in that case, the resolution stops right away as above. (At Stage 1.)

$$\dots \rightarrow 0 \rightarrow 0 \rightarrow 0 \xrightarrow{\text{stage 1}} R^{\oplus m} \xrightarrow{\text{stage 0}} R^{\oplus n} \rightarrow M \rightarrow 0$$

$(m \leq n)$

Thus, every fg. module has a finite free resolution of "length" 1.

(Recall that a submodule of a fg. free module over a PID is free.)

• Question of "optimality" of free resolution

We know $0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$
is a free resolution of $\mathbb{Z}/2\mathbb{Z}$ over \mathbb{Z} .

We know $U \rightarrow \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$ is a free resolution of $\mathbb{Z}/2\mathbb{Z}$ over \mathbb{Z} .

$$0 \rightarrow \mathbb{Z} \xrightarrow{\begin{bmatrix} 0 \\ 1 \end{bmatrix}} \mathbb{Z}e_1 \oplus \mathbb{Z}e_2 \xrightarrow{\begin{bmatrix} 2 & 0 \\ 0 & 1 \end{bmatrix}} \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$$

$f \mapsto e_2$ $e_1 \mapsto 2$
 $e_2 \mapsto 0$

Note that we can go on and create an extra copy of \mathbb{Z} and make it longer.

$$0 \rightarrow \mathbb{Z}g \xrightarrow{\begin{bmatrix} 0 \\ 1 \end{bmatrix}} \mathbb{Z}f_1 \oplus \mathbb{Z}f_2 \xrightarrow{\begin{bmatrix} 0 & 0 \\ 1 & 0 \end{bmatrix}} \mathbb{Z}e_1 \oplus \mathbb{Z}e_2 \xrightarrow{\begin{bmatrix} 2 & 0 \\ 0 & 1 \end{bmatrix}} \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0.$$

$f_1 \mapsto e_2$ $e_1 \mapsto 2$
 $g \mapsto f_2 \mapsto 0$

Q. If $\text{rank}(F_i) \geq \text{rank}(F_{i-1})$, does that mean non-optimal?
 I don't think so. If a column is 0, then yes. But otherwise, don't think so.

Q. What is a free resolution of R over itself

$$0 \rightarrow R \xrightarrow{\text{id}} R \rightarrow 0. \quad \left(\begin{array}{l} \text{Dropping the module:} \\ 0 \rightarrow R \rightarrow 0 \end{array} \right)$$

In fact, the above is true for any free R -module F .

$$0 \rightarrow F \rightarrow F \rightarrow 0. \quad \left(\begin{array}{l} \text{Dropping the module} \\ 0 \rightarrow F \rightarrow 0. \end{array} \right)$$

Both are free resolutions of length 0.

Q. If every fg. R -module has a free resolution of length 0, what is R ?

Ans. length 0 $\Leftrightarrow M$ is free.
 Every R -module has a basis iff R is a field.
 (is free)

Q. If every fg. R -module has a free resolution of length 1, what is R ?

Ans. Field or PID.

One way implication is obvious. (Field or PID \Rightarrow length 1)
 (\Leftarrow) Suppose I is an ideal. Let R/I have a resolution of length 1.
 f.g.

$$0 \rightarrow F_1 \rightarrow F_0 \rightarrow R/I \rightarrow 0$$

We also have

$$0 \rightarrow I \rightarrow R \rightarrow R/I \rightarrow 0.$$

Have to conclude I is free. That would imply I is principal

INCOMPLETE.

Back to example: $0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$ is a \mathbb{Z} -free resolution of $\mathbb{Z}/2\mathbb{Z}$.

Tensor with $\mathbb{Z}/6\mathbb{Z}$ over \mathbb{Z} . \leftarrow We get a sequence of $\mathbb{Z}/6\mathbb{Z}$ modules via base change.

① Note that $\mathbb{Z}/2\mathbb{Z}$ is a $\mathbb{Z}/6\mathbb{Z}$ module. ($\bar{i}_{\mathbb{Z}/6\mathbb{Z}} \mapsto \bar{i}_{\mathbb{Z}/2\mathbb{Z}}$ gives a ring homom)

$$\textcircled{2} \quad \mathbb{Z}/6\mathbb{Z} \otimes_{\mathbb{Z}} \mathbb{Z}/2\mathbb{Z} \cong_{\mathbb{Z}} \mathbb{Z}/(6\mathbb{Z} + 2\mathbb{Z}) = \mathbb{Z}/2\mathbb{Z}$$

\hookrightarrow should be generated by $\{\bar{i} \otimes \bar{i}\}$ (recall tensor generated by tensor of gen.)

$$\textcircled{3} \quad \mathbb{Z}/6\mathbb{Z} \otimes_{\mathbb{Z}} \mathbb{Z} \cong_{\mathbb{Z}} \mathbb{Z}/6\mathbb{Z}$$

Thus, tensoring the free resolution with $\mathbb{Z}/6\mathbb{Z}$ gives

$$0 \rightarrow \mathbb{Z}/6\mathbb{Z} \xrightarrow[\bar{i}]{?} \mathbb{Z}/6\mathbb{Z} \xrightarrow[\bar{i}]{?} \mathbb{Z}/2\mathbb{Z} \rightarrow 0$$

$\bar{i} \mapsto \bar{i}$

Is this a free resolution of $\mathbb{Z}/2\mathbb{Z}$ over $\mathbb{Z}/6\mathbb{Z}$?

No! The map $\bar{1} \mapsto \bar{2}$ is not injective.
 $3\mathbb{Z}/6\mathbb{Z}$ is the kernel!
 \hookrightarrow not free \therefore

After base change, the free resolution does not remain
a resolution.

Lecture 5 (19-01-2021)

19 January 2021 11:31

Recall: $0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$ is a \mathbb{Z} -free resolution of $\mathbb{Z}/2\mathbb{Z}$

Tensoring with $\mathbb{Z}/6\mathbb{Z}$ does not give a $\mathbb{Z}/6\mathbb{Z}$ -free resolution of $\mathbb{Z}/6\mathbb{Z} \otimes_{\mathbb{Z}} \mathbb{Z}/2\mathbb{Z} \cong \mathbb{Z}/2\mathbb{Z}$.

It gives a free "complex" with "homology".

Note If S is an R -algebra (via φ), $M \xrightarrow{\psi} N$ is R -linear, then ψ induces a natural S -linear map $S \otimes_R M \rightarrow S \otimes_R N$.

Defⁿ A **complex** of R -modules is a sequence (finite or countable) of R modules with maps between them

$$\dots \rightarrow M_{n+1} \xrightarrow{\varphi_{n+1}} M_n \xrightarrow{\varphi_n} M_{n-1} \rightarrow \dots \quad \text{such that}$$

$$\ker \varphi_n \supset \operatorname{im} \varphi_{n+1}, \quad \text{i.e.,} \quad \varphi_n \circ \varphi_{n+1} = 0 \quad \forall n.$$

(Finite/infinite in one or both directions)

The complex is **exact** at n^{th} stage if $\ker \varphi_n = \operatorname{im} \varphi_{n+1}$.

\mathbb{Z} -modules: $0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z} \rightarrow 0$

(dropping the modules)

$\mathbb{Z}/6\mathbb{Z}$ -modules:
(after knowing)

$$0 \xrightarrow{\varphi_2} \mathbb{Z}/6\mathbb{Z} \xrightarrow{\varphi_1} \mathbb{Z}/6\mathbb{Z} \xrightarrow{\varphi_0} 0$$

$$\bar{i} \mapsto \bar{2}; \quad \bar{i} \mapsto 0$$

j^{th} stage

$$\ker \varphi_0 = \mathbb{Z}/6\mathbb{Z}$$

$$\operatorname{im} \varphi_1 = 2\mathbb{Z}/6\mathbb{Z}$$

NOT EXACT!

1^{st} stage

$$\ker \varphi_1 = 3\mathbb{Z}/6\mathbb{Z}$$

$$\operatorname{im} \varphi_2 = 0$$

NOT EXACT!

In both cases, we do have $\text{im} \subset \ker$. Thus, it is indeed a complex

Q: Note that $0 \rightarrow M \xrightarrow{\varphi} N \rightarrow 0$ is always a complex.
When is it exact?

A: Precisely when φ is an iso. $\left(\begin{array}{l} 0^{\text{th}} \text{ stage exact} \Leftrightarrow \varphi \text{ is onto} \\ 1^{\text{st}} \text{ stage exact} \Leftrightarrow \varphi \text{ is 1-1} \end{array} \right)$

Remark $0 \rightarrow M \xrightarrow{\varphi} N$ is exact $\Leftrightarrow \varphi$ is 1-1
 $M \xrightarrow{\varphi} N \rightarrow 0$ is exact $\Leftrightarrow \varphi$ is onto

Defⁿ A complex $\dots \rightarrow M_{n+1} \rightarrow M_n \rightarrow M_{n-1} \rightarrow \dots$ is exact if it is exact at each stage n .

Remark. Thus, if it is not exact, then it is not exact at some n .
That is, $\text{im } \varphi_{n+1} \not\subset \ker \varphi_n$.
Since we do have containment, the quotient $\ker \varphi_n / \text{im } \varphi_{n+1}$ makes sense. $\text{im } \varphi_{n+1} = \ker \varphi_n \Leftrightarrow \ker \varphi_n / \text{im } \varphi_{n+1} = 0$.

Defⁿ Given a complex C , we define the n^{th} homology of C as

$$H_n(C) = \frac{\ker \varphi_n}{\text{im } \varphi_{n+1}}$$

Thus, the homology is an "obstruction" to the complex being exact!

(Familiar examples: $\ker \varphi$ is an obstruction to φ being 1-1.
 $\text{coker } \varphi$ for onto. $[G_1, G_2]$ for G being abelian.)

Ex. Find the homologies in the previous examples.

Example. If $F: \dots \rightarrow F_2 \rightarrow F_1 \rightarrow F_0 \rightarrow 0$ is a free resolution of M , then

$$H_n(F_\bullet) = \begin{cases} 0 & ; \quad n > 0, \\ M & ; \quad n = 0. \end{cases}$$

(This was a complex, by construction; similar reason for $n > 0$ stages) being exact.

Remark. If S is an R -algebra, C_\bullet is an exact complex of R -modules, $S \otimes_R C_\bullet$ is a complex but not necessarily exact.

Notⁿ. An exact complex $0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0$ is called a **short exact sequence (s.e.s.)**.

This gives φ is 1-1, ψ is onto.
 \hookrightarrow can think of L as a submodule of M
 $\text{im } \varphi = \text{ker } \psi$ gives $M/L \cong N$.

Properties of submodules and quotients can be transferred to s.e.s.s usually.
 E.g. M is Noe (Art) $\Leftrightarrow N$ and L are Noe. (Art.)

Functorial properties of \otimes

① Fix an R -module K . Define $T(M) = K \otimes_R M$ for all R -modules M .

② Note: For any R -module M , $T(M)$ is also an R -module.

③ Given $\varphi: M \rightarrow N$, R -linear, we get an induced map
 $T(\varphi): T(M) \rightarrow T(N)$

$$\left[\begin{array}{l} T(\varphi): K \otimes_R M \rightarrow K \otimes_R N \text{ defined by} \\ [T(\varphi)](x \otimes y) = x \otimes \varphi(y). \end{array} \right] \text{ need to verify}$$

Notⁿ. $T(\varphi) = \text{id} \otimes \varphi$ or φ_* .

This assignment T has the following properties:

This assignment T has the following properties:

- (a) $T(\psi \circ \varphi) = T(\psi) \circ T(\varphi)$ $M \xrightarrow{\varphi} N \xrightarrow{\psi} L$
 $T(M) \xrightarrow{T(\varphi)} T(N) \xrightarrow{T(\psi)} T(L)$
- (b) $T(\text{id}_M) = \text{id}_{T(M)}$
- (c) $T(0) = 0$ (0 module or 0 map? Yes.)
- (d) $T(M \oplus N) = T(M) \oplus T(N)$

T above is a functor by (1) - (4). Denoted by $K \otimes_R -$.
It is a covariant functor, since arrows are in same direction.

Furthermore, $T(0) = 0$ show that if C is a complex of R -modules, then so is $T(C)$.

(Since compositions go to compositions and 0 goes to 0.)

Q. Does $K \otimes_R -$ preserve exactness?

No! We already have examples. Every one: $0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z}$ is exact.
Apply $\mathbb{Z}/2\mathbb{Z} \otimes_{\mathbb{Z}} -$ to get answer.

Note that if S is an R -algebra (via φ), then base change (i.e., $S \otimes_R -$) gives a functor from R -modules to S -modules. (Covariant or contravariant.)

Q. If $0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0$ is a s.e.s. of R -modules and K is a fixed R -module, what can you say about

$$K \otimes_R (0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0).$$

Examples of functors: ① Localisation: R -modules $\rightarrow R_A$ -modules.

② Forgetful functors: Group/Rings/etc. \rightarrow Set

③ Identity/inclusion functor

④ Linearisation: Set $\rightarrow R$ -modules ($A \mapsto F(A)$)

⑤ Fundamental group : $\text{Top.} \rightarrow \text{Grp}$
 $\text{Top} \rightarrow \text{Ring}$
 $X \mapsto \{ \text{continuous } f: X \rightarrow \mathbb{R} \}$

⑥ K a field.

Field extensions of $K \rightarrow K$ -vector spaces
Is this a functor?

(Note that we have to think of the morphisms as well.)

Lecture 6 (21-01-2021)

21 January 2021 09:30

Q. What can we say about $K \otimes_R (0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0)$?

A. We do get a complex

$$0 \rightarrow K \otimes_R L \xrightarrow{\varphi_*} K \otimes_R M \xrightarrow{\psi_*} K \otimes_R N \rightarrow 0$$

$$\varphi_* = \text{id}_K \otimes \varphi, \quad \psi_* = \text{id}_K \otimes \psi$$

Do we have exactness at all three points?

- ① Is φ_* 1-1?
- ② Is φ_* onto?
- ③ Is $\text{im } \varphi_* = \text{ker } \psi_*$?

① No. $K \otimes_R -$ does not take injective maps to injective maps.
 $\mathbb{Z}/6\mathbb{Z} \otimes_{\mathbb{Z}} (0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z})$ was seen.

② Yes. $K \otimes_R N$ is generated by $\{x \otimes z \mid x \in K, z \in N\}$
 If ψ is surjective, $\forall z \in N, \exists y \in M$ s.t. $\psi(y) = z$.
 Hence,

$$x \otimes z = x \otimes \psi(y) = \varphi_*(x \otimes y).$$

Thus, $K \otimes_R -$ takes surjections to surjections.

③ We already know it is a complex. Thus, $\text{im}(\varphi_*) \subset \text{ker}(\psi_*)$.

(\supset) Let $\sum x_i \otimes y_i \in \text{ker}(\psi_*)$. \leftarrow Doing this will almost never work.

$$\text{We prove the natural map } \frac{K \otimes_R M}{\text{im } \varphi_*} \xrightarrow{\pi} \frac{K \otimes_R M}{\text{ker } \psi_*}$$

is an isomorphism.

This would prove $\text{im } \varphi_* = \text{ker } \psi_*$.

this is the natural onto obtained because $\text{im} \subset \text{ker}$ is known

To do this, we prove that the map

$$\frac{K \otimes_R M}{\text{im } \varphi_*} \xrightarrow{\pi} \frac{K \otimes_R M}{\text{ker } \psi_*}$$

map induced by quotient

$$\begin{array}{ccc} k \otimes_R M & \xrightarrow{\psi_* \pi} & k \otimes_R N \rightarrow 0 \\ \text{im } \psi_* & & \end{array}$$

$\bar{\psi}_* : \frac{k \otimes_R M}{\ker \psi_*} \rightarrow k \otimes_R N$ is invertible.

Construct the inverse as follows

$$\forall z \in k, z \in N, \text{ choose } y \in M \text{ s.t. } \psi(y) = z.$$

(Ex 1.) Verify $z \otimes z \mapsto z \otimes y + \text{im } \psi_*$ is a well-defined R -linear map. This is an inverse of $\bar{\psi}_* \pi$.

Thus, we have shown: If $0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0$ is an s.e.s, then $k \otimes_R L \xrightarrow{\psi_*} k \otimes_R M \xrightarrow{\psi_*} N \otimes_R M \rightarrow 0$ is exact and injective maps need not go to injective maps.

We say that $k \otimes_R -$ is not an exact functor but a right exact functor. In fact, our proof did not need injectivity of φ .

(Right exact functor, exact functor)

This is the defⁿ of right exactness. \rightarrow Thus, $L \rightarrow M \rightarrow N \rightarrow 0$ exact $\Rightarrow k \otimes_R (L \rightarrow M \rightarrow N \rightarrow 0)$ exact. Note the lack of $0 \rightarrow$ here.

Hom as a functor

Note that $k \otimes_R -$ and $- \otimes_R k$ are the same due to commutativity. However, as Hom is different.

Ex. Given R -modules M and N , is $\text{Hom}_R(M, N) \cong \text{Hom}_R(N, M)$?
No. Take $R = M = \mathbb{Z}$ and $N = \mathbb{Z}/2\mathbb{Z}$.

(1) Fix an R -module K .

Given an R -module M , $\text{Hom}_R(M, K)$ is an R -module.

Given an R -linear map $\varphi: M \rightarrow N$, we get a function

$$\begin{array}{ccc}
 M & \xrightarrow{\varphi} & N \\
 \swarrow & & \searrow \\
 & K &
 \end{array}
 \qquad
 \begin{array}{ccc}
 \text{Hom}(N, K) & \xrightarrow{\varphi_*} & \text{Hom}(M, K) \\
 \alpha & \longmapsto & \alpha \circ \varphi
 \end{array}$$

Note the reversal!

This association respects compositions (reverse), identity, and zero.

Thus, $\text{Hom}_R(-, K)$ is a **contravariant functor**.

(Contravariant functor)

Since it preserves compositions and zeroes, it preserves complexes. That is, given an s.e.s

$$0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0, \quad \text{we get a complex}$$

$$0 \rightarrow \text{Hom}_R(N, K) \xrightarrow{\psi_*} \text{Hom}_R(M, K) \xrightarrow{\varphi_*} \text{Hom}_R(L, K) \rightarrow 0$$

At what place(s) is it exact?

(2) Fix K . Given M , $\text{Hom}_R(K, M)$ is an R -module and given $\varphi: M \rightarrow N$, we get a map

$$\varphi_* = \text{Hom}(K, \varphi) : \text{Hom}(K, M) \rightarrow \text{Hom}(K, N).$$

(ex. 2.) Verify that $\text{Hom}_R(-, K)$ is a contravariant functor, and $\text{Hom}_R(K, -)$ is a covariant functor. Both are left exact.

(ex. 1.) First, we show $\mathcal{I} : K \times N \rightarrow \frac{K \otimes_R M}{\text{im } \varphi_*}$

$$(x, z) \mapsto x \otimes y + \text{im } \varphi_*$$

well-defined

$$0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0$$

Note that if $\varphi(y_1) = \varphi(y_2)$, then $y_1 - y_2 \in \ker \varphi = \text{im } \varphi$.

$$\begin{aligned} \text{Thus, } y_1 &= y_2 + \varphi(x') \text{ for some } x' \in L. \\ \Rightarrow x \otimes y_1 &= x \otimes y_2 + x \otimes \varphi(x') \\ \Rightarrow x \otimes y_1 &= x \otimes y_2 + \varphi_*(x \otimes x') \end{aligned}$$

Thus, $x \otimes y_1 + \text{im } \varphi_* = x \otimes y_2 + \text{im } \varphi_*$.

Thus, Φ is a well-defined map. That it is bilinear is clear.

Thus we get a well-defined map

$$\tilde{\Phi} : K \otimes_R N \rightarrow \frac{K \otimes_R M}{\text{im } \varphi_*} \text{ defined by}$$

$$x \otimes z \mapsto x \otimes y + \text{im } \varphi_*$$

It suffices to show that it is the left inv of $\tilde{\Psi}_* \pi$. ← already know it is onto

$$\frac{K \otimes_R M}{\text{im } \varphi_*} \xrightarrow{\pi} \frac{K \otimes_R M}{\ker \Psi_*} \xrightarrow{\tilde{\Psi}_*} K \otimes_R N \xrightarrow{\tilde{\Phi}} \frac{K \otimes_R M}{\text{im } \varphi_*}$$

on generators:

$$x \otimes y + \text{im } \varphi_* \mapsto x \otimes y + \ker \Psi_* \mapsto x \otimes \Psi(y) \mapsto x \otimes y + \text{im } \varphi_* \checkmark$$

(Ex. 2.)

$$L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0 \quad \text{exact}$$

↓

$$0 \rightarrow \text{Hom}(N, K) \xrightarrow{\Psi_*} \text{Hom}(M, K) \xrightarrow{\varphi_*} \text{Hom}(L, K) \quad \text{exact}$$

• Ψ_* 1-1

Let $\alpha \in \text{Hom}(N, K)$ s.t. $\Psi_* \alpha = 0$.

$$\begin{array}{ccc} N & \xrightarrow{\alpha} & K \\ \uparrow \psi & \nearrow & \\ M & \xrightarrow{\Psi_* \alpha} & \end{array} \quad \Psi_* \alpha = \alpha \circ \psi$$

Thus, $\alpha \circ \psi = 0$ map

$$\Rightarrow (\alpha \circ \psi)(m) = 0 \quad \forall m \in M$$

$$\Rightarrow \alpha(\psi(m)) = 0 \quad \forall m \in M$$

$$\Rightarrow \alpha(n) = 0 \quad \forall n \in N \quad (\because \psi \text{ is onto})$$

$$\Rightarrow \alpha = 0 \quad \text{map} \quad \square$$

• $\text{im } \varphi_* = \ker \Psi_*$

(\subseteq) since complex. $\in \text{Hom}(M, K)$

$$\begin{array}{ccc} N & & \\ \uparrow \psi & \dashrightarrow \alpha & \\ M & \xrightarrow{\varphi_*} & K \\ \uparrow \varphi & \nearrow \beta \circ \varphi & \\ L & & \end{array}$$

(2) Let $\beta \in \ker \varphi_*$. Then, $\varphi_* \beta = 0$ map
 $\Leftrightarrow \beta \circ \varphi = 0$ map
 $\Rightarrow \beta \circ \varphi(l) = 0 \quad \forall l \in L$

$$\boxed{\beta \in \text{im } \varphi_*} \\ \Rightarrow \beta = \alpha \circ \varphi$$

$$\Rightarrow \ker \beta \supset \text{im } \varphi = \ker \psi$$

Thus, if $\beta(m_1) = \beta(m_2)$, then
 $\psi(m_1) = \psi(m_2)$.

Thus, by UMP of quotients, $\alpha(n) = \alpha(\psi(m)) = \beta(m)$

is well defined and R -linear.

Thus, $\beta = \alpha \circ \psi = \varphi_* \alpha \in \text{im } \varphi_*$. \square

$$0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \quad \text{exact}$$

$$\Downarrow \\ 0 \rightarrow \text{Hom}(K, L) \xrightarrow{\varphi_*} \text{Hom}(K, M) \xrightarrow{\psi_*} \text{Hom}(K, N) \quad \text{exact}$$

• φ_* is 1-1

Let $\alpha \in \ker \varphi_*$. $\varphi_* \alpha = 0 \Rightarrow \psi \circ \alpha = 0$
 $\Rightarrow (\psi \circ \alpha)(k) = 0 \quad \forall k \in K$
 $\Rightarrow \alpha(k) = 0 \quad \forall k \in K \quad (\because \psi \text{ is 1-1})$
 $\Rightarrow \alpha = 0.$ \square

• $\text{im } \varphi_* = \ker \psi_*$

(\Leftarrow) clear.

(\Rightarrow) Let $\beta \in \ker \psi_*$. Then, $\psi_* \beta = 0$
 $\Rightarrow (\psi_* \beta)(k) = 0 \quad \forall k$
 $\Rightarrow \beta(k) \in \ker \psi = \text{im } \varphi \quad \forall k$

$$\begin{array}{ccc} & & L \\ & \alpha \nearrow & \downarrow \varphi \\ K & \xrightarrow{\beta} & M \\ & \searrow \psi \circ \beta & \downarrow \psi \\ & & N \end{array}$$

$\Rightarrow \forall k \exists l_k \in L$ s.t. $\varphi(l_k) = \beta(k)$.

$\because \varphi$ is 1-1, $\exists! l_k$

Moreover, $\alpha = (k \mapsto l_k)$ is R -linear.

Thus, $\alpha \in \text{Hom}(K, L)$ and

$$\beta(k) = \varphi(l_k) = (\varphi \circ \alpha)(k) \quad \forall k.$$

Thus, $\beta = \varphi_* \alpha = \varphi_* \alpha \in \text{im } \varphi_*$. \square

Lecture 7 (25-01-2021)

25 January 2021 10:35

Covariant Hom is **left exact**, i.e.,

$$0 \rightarrow L \rightarrow M \xrightarrow{\text{exact}} N \Rightarrow 0 \rightarrow \text{Hom}_R(K, L) \rightarrow \text{Hom}_R(K, M) \xrightarrow{\text{exact}} \text{Hom}_R(K, N)$$

Contra variant Hom is also **left exact**, i.e.,

$$L \rightarrow M \xrightarrow{\text{exact}} N \rightarrow 0 \Rightarrow 0 \rightarrow \text{Hom}_R(L, K) \rightarrow \text{Hom}_R(M, K) \xrightarrow{\text{exact}} \text{Hom}_R(N, K)$$

(Note the assumption of $\rightarrow 0$ on LHS.)

Usually, we will be more relaxed and start with an s.e.s. to begin with.

Q. Is $K \otimes_R -$ exact? No. $K = \mathbb{Z}/6\mathbb{Z}$ over $R = \mathbb{Z}$.
Is $K \otimes_R -$ exact for some K over some R ? Yes, $K = R$ for any R .

(1) Over any R , can you find a K s.t. $K \otimes_R -$ is not exact.
(2) Can you find a class of examples of K s.t. (a) $K \otimes_R -$ is not exact?
(b) $K \otimes_R -$ is exact?

Q. Can ask and (try to) answer similar questions about both the Hom functors.

Defⁿ A s.e.s. $0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0$ is **split exact** if \exists a map $\chi: N \rightarrow M$ which is a **splitting**, i.e., $\psi \circ \chi = \text{id}_N$.

(Split exact sequences)

(map will always refer to the appropriate morphisms.)

Ex. (1) $M = \varphi(L) \oplus \chi(N)$

② α is injective and hence $M \cong L \oplus N$.

① Split exact sequence captures the notion of \oplus .
(s.e.s. captures submodule and quotient.)

If $M = L \oplus N$, then $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$
is split exact with the natural maps.

② Let F be a functor from R -modules to R -modules which
is additive. If

(Definition at end)

$0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0$ is split exact,

then what can one say about

$0 \rightarrow F(L) \xrightarrow{\varphi_*} F(M) \xrightarrow{\psi_*} F(N) \rightarrow 0$?

→ The "splitness" is preserved. $\psi \alpha = \text{id}_N \Rightarrow \psi_* \alpha_* = \text{id}_{F(N)}$.
In particular, ψ_* remains surjective.

Is φ_* injective?

Actually, splitting gives a map $\pi: M \rightarrow L$ as well s.t.
 $\pi \circ \varphi = \text{id}_L$

Thus, $\pi_* \circ \varphi_* = \text{id}_{F(L)}$ and thus, φ_* is injective.

$$0 \begin{array}{c} \xrightarrow{\varphi} \\ \xleftarrow{\pi} \end{array} L \begin{array}{c} \xrightarrow{\varphi} \\ \xleftarrow{\pi} \end{array} M \begin{array}{c} \xrightarrow{\psi} \\ \xleftarrow{\alpha} \end{array} N \begin{array}{c} \xrightarrow{\psi} \\ \xleftarrow{\alpha} \end{array} 0$$

The below seq. is exact as well.

We also can say:

$0 \rightarrow N \xrightarrow{\alpha} M \xrightarrow{\pi} L \rightarrow 0$ is split exact
with $\varphi: L \rightarrow M$ being the splitting map.

To conclude:

$$0 \rightarrow F(L) \xrightarrow{\psi_L} F(M) \xrightarrow{\psi_M} F(N) \rightarrow 0$$

is split exact. (Verify at middle point. Do we need additivity?)

Thus, if $E = (0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0)$ is split exact, then, for any K : $K \otimes_R E$, $\text{Hom}_R(K, E)$, $\text{Hom}_R(E, K)$ are all split exact.

Defⁿ (Additive functor)

A functor $F: R\text{-Mod} \rightarrow R\text{-Mod}$ is called **additive** if

covariant

$F: \text{Hom}_R(M, N) \rightarrow \text{Hom}_R(F(M), F(N))$ is a group homomorphism.

Free resolutions (II)

$$\begin{array}{cccc} \mathbb{Z} & \mathbb{Z} & \mathbb{Z} & \mathbb{Z}/\mathbb{Z} = 0 \\ & \underline{\underline{=}} & \underline{\underline{=}} & \\ \mathbb{Z}/2\mathbb{Z} & \rightarrow & & \end{array}$$

$\rightarrow M, L_1, L_2$

Q! Given an s.e.s. $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$ of R -modules, if we know free resolutions of two of the three, can we construct a free resolution for the third?

(Given our experience, we might expect that a resolution of M gives for both. However, this is "hopeless".)

Example. Take N as any module, we know that \exists free F s.t. $F \rightarrow N \rightarrow 0$.
Take $0 \rightarrow \ker \rightarrow F \rightarrow N \rightarrow 0$.

know \leftarrow for this!

As an example

$$0 \rightarrow I \rightarrow R \rightarrow R/I \rightarrow 0.$$

\hookrightarrow this can't tell for both ends

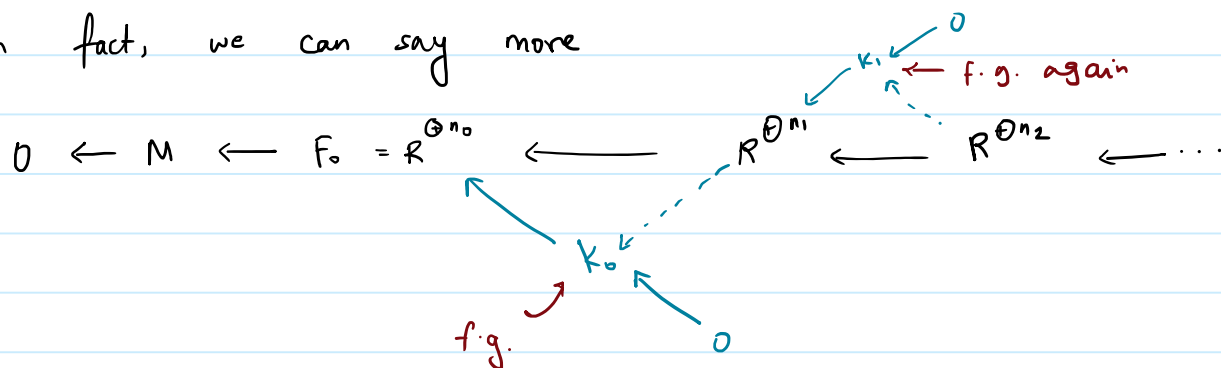
Q2. Suppose M is f.g. and F a finite rank free module s.t. $M \cong F/K$. Is K f.g.?

No. Take R to be non-noe., let $I \subseteq R$ be a non-f.g. ideal. Then, $F = R$, $M = R/I$, $K = I$ is a counterexample.

Q3. With same notation, give a condition of R which forces K to be f.g.

Ans. R is Noetherian. Then, $F = R^{\oplus n}$ and hence, $K \subseteq F$ is f.g.

In fact, we can say more



Thus, over a Noetherian ring R , a f.g. module M has a free resolution of the form

$$F : \dots \rightarrow F_2 \xrightarrow{\varphi_2} F_1 \xrightarrow{\varphi_1} F_0 \rightarrow 0 \quad \text{where}$$

$$F_i \cong R^{\oplus n_i}$$

Hence, fixing bases for F_i , we can write φ_i as matrices.

Defn

(Presentation)

A matrix representation of φ_i is called a **presentation** of M .

Q4. If M is f.g. and $M \cong F_1/K_1 \cong F_2/K_2$, where $F_1 \cong R^{\oplus n_1}$ and $F_2 \cong R^{\oplus n_2}$, then

(a) is $n_1 = n_2$? No. Have seen already. Can always pad more R s.

(b) Is it necessary that $K_1 \cong K_2$? $\rightarrow n_1 = n_2 = 1$, we know

(c) How are F_1 and F_2 related?

Think about examples: $\mathbb{Z}, \mathbb{K}[x] \rightarrow$, $\frac{\mathbb{K}[x,y]}{\mathcal{I}}$ where

Think of \mathbb{Z} as rings or \mathbb{Z} -modules $\rightarrow \mathbb{Z}/4\mathbb{Z}, \mathbb{Z}/6\mathbb{Z}$ \downarrow quotients

Construct modules over it \downarrow $\mathbb{Z}/2\mathbb{Z} \oplus \mathbb{Z}/3\mathbb{Z}$

$\mathcal{I} = \langle x, y \rangle, \langle x^2, y^2 \rangle, \langle x^2, xy \rangle, \langle x^2, xy, y^2 \rangle, \langle x \rangle$

$\frac{\mathbb{K}[x,y,z]}{\mathcal{I}} ; \mathcal{I} = \langle x, y \rangle, \langle x, y, z \rangle, \langle x^2, xy, y^2 \rangle, \dots$

Lecture 8 (26-01-2021)

26 January 2021 11:37

Optimality of free resolutions

Example Consider $M = \mathbb{Z}/6\mathbb{Z}$ as a \mathbb{Z} -module.
 $\mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/3\mathbb{Z}$

Then, $M = \langle (1,1) \rangle = \langle (1,0), (0,1) \rangle$.

$$(i) \quad 0 \longleftarrow M \longleftarrow \mathbb{Z} \xleftarrow{6} \mathbb{Z} \longleftarrow 0$$
$$(1,1) \longleftarrow 1$$

$$(ii) \quad 0 \longleftarrow M \longleftarrow \mathbb{Z}e_1 \oplus \mathbb{Z}e_2 \xleftarrow{\begin{bmatrix} 2 & 3 \end{bmatrix}} \mathbb{Z} \oplus \mathbb{Z} \longleftarrow 0$$
$$(1,0) \longleftarrow e_1$$
$$(0,1) \longleftarrow e_2$$

Which is more optimal? How could we make the first step optimal?

Pick gen. set which is least in cardinality?

(Note that both the gen sets above are minimal, since that is w.r.t. inclusion.)

Does this guarantee optimality in second step?

Recall. If (R, \mathfrak{m}, k) is local and M a f.g. module, then every minimal generating set of M has the same cardinality, denoted $\mu(M)$, where $\mu(M) = \dim_k (M/\mathfrak{m}M)$.

[Follows from Nakayama.]

$(\mu(M))$

We would also want the kernels to be f.g.

Thus, we work in the following setting:

(R, \mathfrak{m}, k) is Noetherian local, M f.g.

(In fact, some authors: local includes Noetherian and non-Noetherian local is quasi-local for them)

A way to ensure optimality: Let $\mu(M) = b_0$
 Map $F_0 = R^{\oplus b_0}$ onto M .

$$0 \leftarrow M \xleftarrow{\langle x_1, \dots, x_{b_0} \rangle} R^{\oplus b_0} \xleftarrow{e_i} K_0 \leftarrow 0$$

Put $b_1 := \mu(K_0)$ and map $F = R^{\oplus b_1}$ onto K_0 and continue.

$$0 \leftarrow R^{\oplus b_0} \xleftarrow{\varphi_1} R^{\oplus b_1} \xleftarrow{\varphi_2} R^{\oplus b_2} \leftarrow \dots$$

Here, $b_0 = \mu(M)$ and $b_i = \mu(\text{im } \varphi_i)$ for $i \geq 1$.

This is called a **minimal free resolution** of M over R .

(Minimal free resolution)

Q. If $\langle y_1, \dots, y_{b_0} \rangle = M$ and K_0' is the kernel obtained by mapping $e_i \mapsto y_i$, how are K_0 and K_0' related?

Is $\mu(K_0) = \mu(K_0')$? \rightarrow This guarantees b_1 is well-defined.

Doesn't guarantee anything for b_2 , however.

Would like to see: $K_0 \cong K_0'$? If yes, then everything would go well ad infinitum.

($n = b_0$)

Note that $y_j \in \langle x_1, \dots, x_n \rangle \quad \forall j$ and
 $x_i \in \langle y_1, \dots, y_n \rangle \quad \forall i$.

$$y_j = a_{j1}x_1 + \dots + a_{jn}x_n \quad j = 1, \dots, n$$

That is,

$$\begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix} = A \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

Similarly,

$$\begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = B \begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix}$$

where $A, B \in M_n(R)$.

Note

$$\begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = BA \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

Idea: Try to show A is invertible. → hope to show $K_i \cong K_i$ with this.

Note that: modulo \mathfrak{m} , $\overline{BA} = \overline{id}$ since $\{\overline{x}_1, \dots, \overline{x}_n\}$ is a basis and there is a unique way to express it in terms of itself.

Thus, $\det(\overline{BA}) \neq 0 \pmod{\mathfrak{m}}$

$\Rightarrow \det BA \notin \mathfrak{m}$

$\Rightarrow \det BA$ is a unit in R

$\Rightarrow BA$ is invertible in $M_n(R)$

$\Rightarrow A$ and B are invertible in $M_n(R)$.

) (R, \mathfrak{m}) local

Lecture 9 (28-01-2021)

28 January 2021 09:16

Setup: (R, \mathfrak{m}, k) local Noetherian

Would like the following: ① A minimal free resolution of M over R (say F_\bullet) is truly minimal in the following sense:

If G_\bullet is a free resolution of M , then

$$\text{rank}(F_i) \leq \text{rank}(G_i) \quad \forall i.$$

② If F_\bullet is of the form $\dots F_1 \xrightarrow{\varphi_1} F_0 \rightarrow M \rightarrow 0$
and G_\bullet is of the form $\dots G_1 \xrightarrow{\psi_1} G_0 \rightarrow M \rightarrow 0$

We would like to relate $\ker \varphi_1$ and $\ker \psi_1$. Would

like $\ker \varphi_1 \subset \ker \psi_1$.

All relations here are also here

③ For the next step:

$$F_2 \xrightarrow{\varphi_2} F_1 \rightarrow \ker \varphi_1 \rightarrow 0$$

$$G_2 \xrightarrow{\psi_2} G_1 \rightarrow \ker \psi_1 \rightarrow 0$$

What now? We only expect $\ker \varphi_1 \subset \ker \psi_1$.

How do $\text{rank } F_1$ and $\text{rank } G_1$ compare?

Is there a relation between $\ker \varphi_2$ and $\ker \psi_2$?

Note that we have defined "minimal" last time, least cardinality of generating set at each step. Want to know if it is truly minimal.

The following technical lemma takes care of it:

Lemma (Splitting Lemma) Let M and N be f.g. R -modules, where (R, \mathfrak{m}, k) is a local Noetherian ring. Let F and G

be free modules mapping onto M and $M \oplus N$, respectively. Further assume that $\text{rank } F = \mu(M)$. Then, F splits off G (i.e., \exists an R -module P s.t. $G \cong F \oplus P$) in a "natural way."

Moreover, if $\varphi: F \rightarrow M$ and $\psi: G \rightarrow M \oplus N$ are the given maps, then $\ker \varphi$ splits off $\ker \psi$.

$\text{rank}(F) = \mu(M)$ ensures the minimality.

Proof:

Consider the s.e.s.s

$$0 \rightarrow L \hookrightarrow G \xrightarrow{\psi} M \oplus N \rightarrow 0,$$

$$0 \rightarrow K \hookrightarrow F \xrightarrow{\varphi} M \rightarrow 0,$$

where $L = \ker \psi$ and $K = \ker \varphi$.

Note that $\text{rank}(F) \leq \text{rank}(G)$ since F maps minimally onto M . (This trivially gives a splitting of F off G , by the way.)

Let $\{e_1, \dots, e_r\}$ and $\{f_1, \dots, f_m\}$ be bases of F and G (over R), respectively. but we want a "natural" one!

Let $x_i = \varphi(e_i)$ and $y_j = \psi(f_j)$. Then, $\{x_1, \dots, x_n\}$ is a minimal gen set of M over R .

(That is, $\{\bar{x}_1, \dots, \bar{x}_n\}$ is a k -basis of $M/\mu(M)$.)

$$0 \rightarrow L \hookrightarrow G \xrightarrow{\psi} M \oplus N \rightarrow 0$$

$$0 \rightarrow K \hookrightarrow F \xrightarrow{\varphi} M \rightarrow 0$$

$\begin{matrix} \uparrow & \downarrow \pi \\ i & \end{matrix}$

Under the above inclusions and projection, we can "write y_j 's in terms of the x_i 's" and vice-versa.

More precisely:

$$\pi(y_1) = a_{11}x_1 + \dots + a_{1n}x_n$$

⋮

$$\pi(y_m) = a_{m1}x_1 + \dots + a_{mn}x_n$$

$$\begin{bmatrix} \pi(y_1) \\ \vdots \\ \pi(y_m) \end{bmatrix} = A_{m \times n} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

↓
keep in mind that the

Keep in mind that the elements of the columns are "vectors" themselves

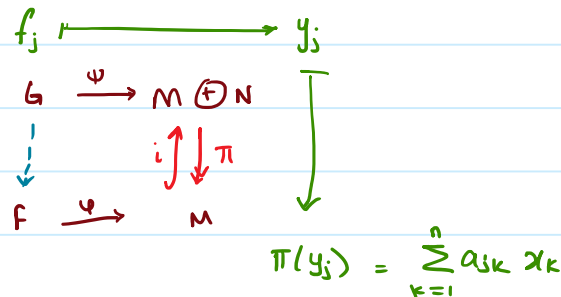
$$i(x_1) = b_{11}y_1 + \dots + b_{1m}y_m$$

$$\vdots$$

$$i(x_n) = b_{n1}y_1 + \dots + b_{nm}y_m$$

$$\begin{bmatrix} i(x_1) \\ \vdots \\ i(x_n) \end{bmatrix} = B_{n \times m} \begin{bmatrix} y_1 \\ \vdots \\ y_m \end{bmatrix}$$

So far, we have

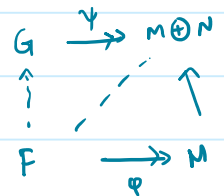
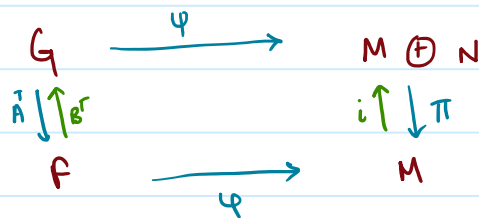


Want a blue map to make square commute.

Define $G \rightarrow F$ by $f_i \mapsto \sum_{k=1}^n a_{jk} e_k$.

Since G and F are free modules, we can represent it as a matrix. It is A^T .

Similarly, we have $F \rightarrow G$ given as B^T .



$$BA \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} c_{11} & \dots & c_{1n} \\ c_{21} & \dots & c_{2n} \\ \vdots & & \vdots \\ c_{n1} & \dots & c_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} a_{11}x_1 + \dots + c_{1n}x_n \\ \vdots \\ a_{n1}x_1 + \dots + c_{nn}x_n \end{bmatrix}_{n \times 1} = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

(The "inner" and "outer" squares commute.)

Note that $\begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = BA \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$ (Why? $\pi i(x_1) = x_1, \dots$)

Go modulo m_j :

$$\begin{bmatrix} \bar{x}_1 \\ \vdots \\ \bar{x}_n \end{bmatrix} = \overline{BA} \begin{bmatrix} \bar{x}_1 \\ \vdots \\ \bar{x}_n \end{bmatrix}, \text{ i.e.,}$$

$$(\bar{I} - \bar{B}A) \begin{bmatrix} \bar{x}_1 \\ \vdots \\ \bar{x}_n \end{bmatrix} = 0.$$

Since $\{\bar{x}_1, \dots, \bar{x}_n\}$ is a basis, we have $\bar{I} = \bar{B}A$.

Thus, $\det(\bar{B}A) = 1$ and hence, $\det(BA) \in 1 + \mathfrak{m}$.

Thus, $\det(BA) \in \mathcal{U}(R)$.

Hence BA is invertible.

Thus, so is $(BA)^T = A^T B^T$. } in $\mathcal{M}_n(R)$

Let $\chi : F \rightarrow F$ be s.t. $\chi A^T B^T : F \rightarrow F$ is idf.

Thus, $B^T : F \rightarrow G$ and $\chi A^T : G \rightarrow F$ are s.t.

$$(\chi A^T) B^T = \text{id}_F.$$

$$\begin{aligned} B^T \chi : F &\rightarrow G \\ A^T : G &\rightarrow F \\ G &= B^T(F) \oplus \ker A^T \end{aligned}$$

$$0 \rightarrow \ker(\chi A^T) \rightarrow G \begin{array}{c} \xrightarrow{\chi A^T} \\ \xleftarrow{B^T} \end{array} F \rightarrow 0$$

Thus, the above s.e.s. splits and hence

$$G = B^T(F) \oplus \ker(\chi A^T).$$

} show the naturality!

Next, we show: $L = B^T(K) \oplus (\ker(\chi A^T) \cap L)$

(Try it!)

• $B^T(K) \subset L$

Proof. Recall: $L = \ker \psi$ and $K = \ker \varphi$.

Let $x \in K = \ker \varphi$.

Then, $\varphi(x) = 0$ and thus, $i(\varphi(x)) = 0$.

But $i(\varphi(x)) = \psi(B^T(x))$.

$\therefore B^T(x) \in \ker \psi = L$. □

$$\begin{array}{ccccccc} 0 & \rightarrow & L & \rightarrow & G & \xrightarrow{\psi} & M \oplus N & \rightarrow & 0 \\ & & & & \uparrow B^T & \uparrow A^T & & & \\ 0 & \rightarrow & K & \rightarrow & F & \xrightarrow{\varphi} & M & \rightarrow & 0 \\ & & & & \downarrow \chi & & & & \end{array}$$

Lecture 10 (01-02-2021)

01 February 2021 10:31

Recall:

$$\begin{array}{ccccccc}
 0 & \rightarrow & L & \hookrightarrow & G & \xrightarrow{\psi} & M \oplus N \rightarrow 0 \\
 & & & & \beta \uparrow \downarrow \alpha & & \pi \downarrow \uparrow i \\
 0 & \rightarrow & K & \hookrightarrow & F & \xrightarrow{\varphi} & M \rightarrow 0
 \end{array}$$

($i\varphi = \psi\beta$)

Where β is multiplication by B^T and α by A^T .
 (After we fixed an appropriate basis.)

We showed: $G = \beta(F) \oplus \ker(\chi\alpha)$

Claim. $\beta(K)$ is a direct summand of L .

Proof. We show that: $L \cap \beta(F) = \beta(K)$

(\subseteq) Let $x \in K$. Then, $\varphi(x) = 0$ and $i\varphi(x) = 0$
 $\psi\beta(x)$

Thus, $\beta(x) \in \ker \psi = L$.

Thus, $\beta(K) \subseteq L \cap \beta(F)$.

(\supseteq) Suppose $y \in L \cap \beta(F)$. Then, $y = \beta(x)$ for some $x \in F$.

We show $x \in K$, i.e., $\varphi(x) = 0$.

Note that

$$\begin{aligned}
 i\varphi(x) &= \psi\beta(x) \\
 &= \psi(y) \hookrightarrow y \in L = \ker \psi \\
 &= 0
 \end{aligned}$$

$\therefore i\varphi(x) = 0$. Since i is 1-1, we get $\varphi(x) = 0$.

This finishes the proof. □

~~————— X —————~~

Restating the lemma:

Lemma

Let M and N be R -modules, $\varphi: F \rightarrow M$, $\psi: G \rightarrow M \oplus N$ be onto, $\ker \varphi = K$, $\ker \psi = L$, $F = R^{\oplus n}$, $G = R^{\oplus m}$, and $n = \mu(M)$. Consider the s.e.s'es

$$\begin{array}{ccccccc}
 0 & \rightarrow & L & \hookrightarrow & G & \xrightarrow{\psi} & M \oplus N \rightarrow 0 \\
 & & & & \uparrow \beta & & \uparrow i \\
 0 & \rightarrow & K & \hookrightarrow & F & \xrightarrow{\varphi} & M \rightarrow 0
 \end{array}$$

Then, $\exists \beta: F \rightarrow G$ s.t.

- ① $\psi\beta = i\varphi$.
- ② $\exists \gamma: G \rightarrow F$ onto s.t. β is a splitting.
- ③ $\beta|_K: K \rightarrow L$ is a splitting.
(Part of ③ that $\psi(K) \subset L$.)

Notation

$K \parallel L$ that K is (isomorphic to) a direct summand of L .

Note:

$$\begin{array}{ccccccc}
 0 & \rightarrow & \overset{\alpha}{\ker \gamma} & \longrightarrow & G & \xrightarrow[\beta]{\gamma} & F \rightarrow 0 \\
 & & & & & & \\
 0 & \rightarrow & \ker \gamma \cap L & \longrightarrow & L & \xrightarrow[\beta]{\gamma} & K \rightarrow 0
 \end{array}$$

Some observations:

① Free modules have a lifting property:

$$\begin{array}{ccccccc}
 G & \xrightarrow{\psi} & M \oplus N & \longrightarrow & 0 \\
 \uparrow \beta & & \nearrow i\varphi & & \\
 F & & & & \\
 \{e_1, \dots, e_n\} & & & &
 \end{array}$$

constructed this $\rightarrow \beta$

Suppose $e_i \xrightarrow{i\varphi} z_i \in M \oplus N$.
Let $z'_i \in G$ be s.t.
 $\psi(z'_i) = z_i$.

Then, we define $\beta(e_i) := z'_i$.

Can do this for every e_i .

Then, this defines a map $f: F \rightarrow G$ since F is free with basis $\{e_i\}_i$.

All we really used is: ① ψ is onto.

② F is free.

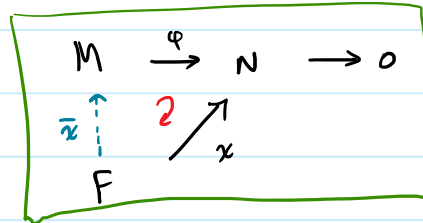
More generally, we have: (Lifting property)

Let ψ be onto and F free.

Then, $\exists \bar{x}: F \rightarrow M$ s.t.

$$\psi \bar{x} = x.$$

(rank $F < \infty$ not necessary.)



Def. This defines an R -module being **projective**.

(One that has the lifting property as above.)

(Projective modules)

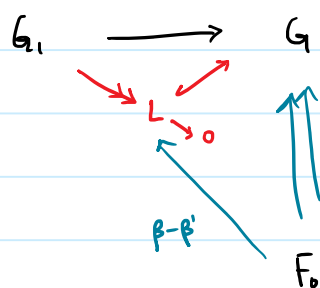
Eg. Free modules are projective.

① Let β, β' be two lifts. How are they related?

$\beta - \beta'$ must be in \ker :

$$\begin{array}{ccccccc} 0 & \rightarrow & L & \rightarrow & G & \rightarrow & M \oplus N \rightarrow 0 \\ & & \beta - \beta' & \searrow & \beta \uparrow \beta' \uparrow & & \\ 0 & \rightarrow & K & \rightarrow & F & \rightarrow & M \rightarrow 0 \end{array}$$

Suppose we have the free resolutions:



By the lifting property, \exists a map $F_0 \rightarrow G_1$, do's

$$\begin{array}{ccccc}
 G_1 & \longrightarrow & G_0 & \longrightarrow & \\
 & & \uparrow & & \\
 & \nearrow & & & \\
 f_1 & \longrightarrow & F_0 & \longrightarrow &
 \end{array}$$

This idea will define a "homotopy" between two maps of chain complexes

③ Definition of chain maps:

$$\begin{array}{ccccccc}
 0 & \longrightarrow & K & \longrightarrow & F_0 & \longrightarrow & M \longrightarrow 0 \\
 & & \nearrow^{f_1} & & \vdots & & \downarrow \cong \\
 0 & \longrightarrow & L & \longrightarrow & G_0 & \longrightarrow & M \oplus N \longrightarrow 0 \\
 & & \searrow_{g_1} & & \vdots & & \\
 & & & & & &
 \end{array}$$

This idea helps to define maps between chain complexes.

Lecture 11 (02-02-2021)

02 February 2021 11:33

Consequences of the proof of the Splitting Lemma :

① if $N=0$, then $L = \beta(K) \oplus \ker(\alpha)$.

② if $m=n$, then ($N=0$ and) $L = \beta(K)$.

In particular, $L \cong K$.

Consequence of the Splitting Lemma:

Thm. Let (R, \mathfrak{m}, K) be a ^{Noetherian} local ring, M a f.g. R -module and $F \rightarrow M$ be a minimal R -free resolution of M . If $G \rightarrow M$ is any R -free resolution of M , then $\forall i \geq 0$, \exists injective maps $\beta_i: F_i \rightarrow G_i$ satisfying:

$$\begin{array}{ccccccc} (1) & 0 & \leftarrow & M & \xleftarrow{\varphi_0} & F_0 & \xleftarrow{\varphi_1} & F_1 & \leftarrow & \dots \\ & & & \text{id}_M \parallel & & \downarrow \beta_0 & & \downarrow \beta_1 & & \\ & 0 & \leftarrow & M & \xleftarrow{\varphi_0} & G_0 & \xleftarrow{\varphi_1} & G_1 & \leftarrow & \dots \end{array}$$

$$\varphi_i \beta_i = \beta_{i-1} \varphi_i \quad \text{or} \quad \varphi \beta = \beta \varphi \quad \text{or} \quad \text{each square commutes}$$

$$\textcircled{2} \quad \beta_i(F_i) \mid G_i, \quad \text{i.e.,} \quad G_i = \beta_i(F_i) \oplus \dots$$

In particular, $\text{rank}(F_i) \leq \text{rank}(G_i)$. (Since β_i is 1-1.)

Proof. We use induction on i ($=:n$) to show that

$$\exists \beta_n: F_n \rightarrow G_n \text{ satisfying } \textcircled{1} \text{ and } \textcircled{2} \\ \text{and } \textcircled{3} \quad \ker \varphi_n = \beta(\ker(\varphi_n)) \oplus \dots$$

The base case $n=0$ is the splitting lemma.

By induction, assume that $\forall i \leq n$, we have

$$\beta_i: F_i \rightarrow G_i$$

satisfying ①, ②, and ③.

$$\begin{array}{ccccccc}
 0 & \rightarrow & L_n & \rightarrow & G_n & \xrightarrow{\psi_n} & G_{n-1} & \rightarrow & \dots \\
 & & & & \beta_n \uparrow & & \downarrow & & \uparrow \beta_{n-1} \\
 0 & \rightarrow & K_n & \rightarrow & F_n & \xrightarrow{\varphi_n} & F_{n-1} & \rightarrow & \dots
 \end{array}$$

We know that $L_n = \beta(K_n) \oplus \dots$

The splitting lemma applied to $0 \rightarrow \ker \psi_{n+1} \rightarrow G_{n+1} \rightarrow L_n \rightarrow 0$
 $0 \rightarrow \ker \varphi_{n+1} \rightarrow F_{n+1} \rightarrow K_n \rightarrow 0$

gives β_{n+1} satisfying ①, ② and ③.

Remark

The compatibility of β with φ and ψ shows that every free resolution of M contains a minimal free resolution.

(Can think of $\beta_i \subseteq$ as inclusions $F_i \hookrightarrow G_i$ and $\psi_i \subseteq$ restrict to maps $F_i \rightarrow F_{i-1}$ where it becomes φ_i .)

Next consequence:

Thm

If F and G are two minimal resolutions of M over R , then $\text{rank}(F_i) = \text{rank}(G_i)$.

Remark

In fact, the two resolutions are "isomorphic".

(We have chain maps which are isomorphisms. ← definition pending)

Def

Let F be a minimal free resolution (m.f.r.) of a f.g. module M over a Noetherian local ring R .

Then,

① the i^{th} Betti number of M over R , denoted $\beta_i^R(M) = \text{rank}_R(F_i)$.

② $\ker \varphi_i = \text{im } \varphi_{i+1}$ is called the $(i+1)^{\text{st}}$ syzygy module

of M over R , denoted $\Omega_{i+1}^{\wedge}(M)$.
 (Note the shift of index, $\chi_0 = \Omega_1^{\wedge}(M)$.)

The above is well-defined, by the discussion above.

Note

① Thus, f breaks into short exact sequences (with $\Omega_0(M) = M$):

$$0 \rightarrow \Omega_{i+1}(M) \rightarrow R^{\oplus p_i} \rightarrow \Omega_i(M) \rightarrow 0.$$

$\Omega_{i+1}(M)$ is the first syzygy of $\Omega_i(M)$ since f_i chosen minimally

② $p_i = \mu(\Omega_i(M))$.

Remark

The above also makes sense in the category of graded modules over graded rings.

(There's a notion of graded local, graded NAK, graded free resolutions, graded free resolution.)

Testing minimality:

Q. What does it mean for $\{x_1, \dots, x_n\}$ to be a minimal generating set of M ?

(Notation: Let $\varphi: F = \bigoplus_{i=1}^n Re_i \rightarrow M$ with $e_i \mapsto x_i$. Assume it actually is a gen. set.)

$\{x_1, \dots, x_n\}$ is a minimal gen. set of M over R

$\Leftrightarrow \{\bar{x}_1, \dots, \bar{x}_n\}$ is a K -basis of $M/\mathfrak{m}M$

Lecture 12 (04-02-2021)

04 February 2021 09:31

Q. What does it mean for $\{x_1, \dots, x_n\}$ to be a minimal generating set of M ?

(Notation: Let $\varphi: F = \bigoplus_{i=1}^n Re_i \rightarrow M$ with $e_i \mapsto x_i$. Assume it actually is a gen. set.)

$\{x_1, \dots, x_n\}$ is a minimal gen. set of M over R

$\Leftrightarrow \{\bar{x}_1, \dots, \bar{x}_n\}$ is a K -basis of $M/\mathfrak{m}M \Leftrightarrow \text{rank } F = \mu(M)$

observe:

$$\begin{array}{ccc} F & \xrightarrow{\varphi} & M \\ \pi \downarrow & & \downarrow \pi \\ F/\mathfrak{m}F & \xrightarrow{\bar{\varphi}} & M/\mathfrak{m}F \end{array}$$

$\varphi: F \rightarrow M$ induces an onto map $\bar{\varphi}: F/\mathfrak{m}F \rightarrow M/\mathfrak{m}M$ of K -vector spaces. (Can verify manually or observe this as tensor $\otimes_R K$.)

Now: $\{x_1, \dots, x_n\}$ is a minimal gen set of $M \Leftrightarrow \bar{\varphi}$ is an iso.

(\Rightarrow) Can use right-exactness to show $\bar{\varphi}$ is onto.

However, $\dim_K(F/\mathfrak{m}F) = \dim_K(M/\mathfrak{m}F) = n < \infty \therefore \bar{\varphi}$ iso. \square

(\Leftarrow) If $S = \{x_1, \dots, x_n\}$ is not minimal, then $\exists A \subsetneq S$ minimal.

But then \bar{A} is a basis with $< n$ elements.

Contradiction since \bar{F} has dim. n . \square

Q. What does this say about $\ker \varphi$?

Ans. $\ker \varphi \subset \mathfrak{m}F$.

Proof. Let $y \in \ker \varphi$. Write $y = \sum a_i e_i$.

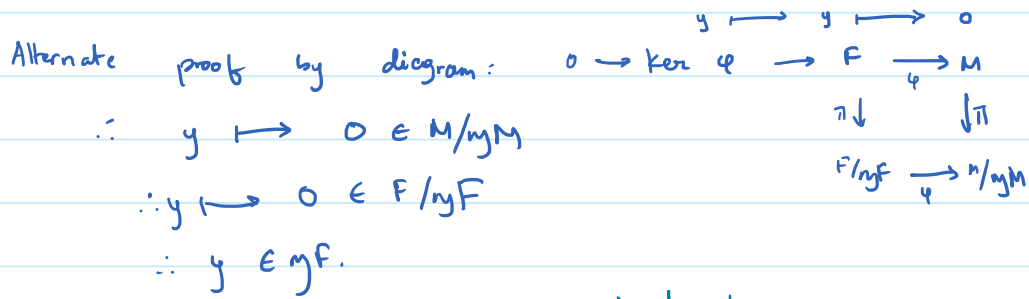
Claim: $a_i \in \mathfrak{m} \forall i$. (This would prove $\ker \varphi \subset \mathfrak{m}F$.)

Proof. $\sum a_i x_i = 0$ in M over R

$\Rightarrow \sum \bar{a}_i \bar{x}_i = 0$ in $M/\mathfrak{m}M$ over R/\mathfrak{m}

$\Rightarrow \bar{a}_i = 0 \forall i$ in R/\mathfrak{m}

$$\Rightarrow a_i \in \mathfrak{m} \quad \forall i \text{ in } R$$



The above is again equivalent. \rightarrow do element wise or use the diagram

Thus, we have: (R, \mathfrak{m}, k) local Noetherian, M f.g.,
 $M = \langle x_1, \dots, x_n \rangle.$

$$F = R e_1 \oplus \dots \oplus R e_n, \quad \varphi : F \rightarrow M \quad \text{where } e_i \mapsto x_i.$$

Then TFAE

- (1) $\{x_1, \dots, x_n\}$ is a minimal gen. set of M .
- (2) $\{\bar{x}_1, \dots, \bar{x}_n\}$ is a basis of $M/\mathfrak{m}M$ (over k)
- (3) $\text{rank } F = \mu(M)$
- (4) $\bar{\varphi} : F/\mathfrak{m}F \rightarrow M/\mathfrak{m}M$ is an iso.
- (5) $\ker \varphi \subset \mathfrak{m}F$

Consider $F = F_0$ and F_1 a free module mapping onto $\ker \varphi$.
 Then, $\text{im } F_1 \subset \mathfrak{m}F_0$. Equivalently, if $\text{im } F_1 \subset \mathfrak{m}F_0$, then
 $\ker \varphi \subset \mathfrak{m}F_0$ and hence, $\{x_1, \dots, x_n\}$ is a min. gen. set of M .

Thus, we have the following: With M and R as above,
 let $F. \rightarrow M$ be a free resolution of M over R .
 Then, $F.$ is minimal $\Leftrightarrow \text{im } \varphi_i \subset \mathfrak{m}F_{i-1} \quad \forall i \geq 1$

$$\dots \rightarrow F_{i+1} \rightarrow F_i \xrightarrow{\varphi_i} F_{i-1} \xrightarrow{\varphi_{i-1}} \dots$$

$\swarrow \text{ker } \varphi_{i-1}$
 \searrow

$\text{im } \varphi_i = \ker \varphi_{i-1} \subset \mathfrak{m}F_{i-1}$ since F_i maps minimally onto $\ker \varphi_{i-2}$.

\Leftrightarrow (writing φ_i as a matrix) the entries of φ_i are in \mathfrak{m} .

Q. Let $G. \rightarrow M$ be a free resolution of M over R .
(R local Noetherian, M f.g.). If $\text{rank}(G_i) = n_i$, what can we
say about the Betti numbers of M ?

Q. If $G.$ is minimal after one stage, then can we write it as
 $(F.) \oplus N$?

Can we drop local-ness?

Lecture 13 (08-02-2021)

08 February 2021 10:36

Note: Over a Noetherian ring, our convention has and will be that the free modules in a free resolution have finite ranks.

Remarks (R, \mathfrak{m}, K) is Noetherian local, M is a f.g. R -module

- ① Bounds on the Betti numbers of M . How to get?
- Upper bound: Construct any free resolution of M over R .
By defⁿ, we will have $\beta_i^R(M) \leq \text{rank}(F_i)$.

Obs. $\beta_0^R(M) = \mu(M) = \dim_K(M/\mathfrak{m}M)$.

- Lower bound: Find a complex of free modules which can be "put inside" a minimal free resolution of M .
(Easier said than done.)

(Open) Conjecture (Buchsbaum - Eisenbud - Horrocks) Given M , $\exists c$ (an invariant of M) such that $\forall i$ $\beta_i^R(M) \geq \binom{c}{i}$.
not stated precisely

(Why binomial coefficients? See Koszul complexes.)

② Poincaré series: $P_M^R(t) = \sum_{t \geq 0} \beta_i^R(M) t^i$

(polynomial iff minimal resolution is finite length)

For $R = \frac{k[x]}{\langle x \rangle}$ and $M = k$, we had seen $P_M^R(t) = \sum_{t \geq 0} t^i$
(in tut.) $= \frac{1}{1-t}$.

Q. Can ask: Is $P_M^R(t) \in \mathbb{Z}(t)$?

↳ rational polynomial

was asked by Serre

Serre proved: For $M = k \leftarrow R/\mathfrak{m}$ the Poincaré series is

term-wise bounded above by a rational function (which comes from a certain Koszul complex) and asked the above Q for $M = k$.

Example 1 $R = \frac{k[x, y]}{\langle x, y \rangle^2}$, $M = k$. Q. Is $P_k^R(t)$ rational?

↳ Noe. local

② $R = \frac{k[x, y]}{\langle x^2, xy \rangle}$ $M = k$. ↗

Anick proved that : $P_k^R(t)$ is not necessarily rational.

Bjurgvad used "idealisation" to show that this fails even for nice rings (Gorenstein rings).

(See A. Kustin's write-up "Georgia Southern University talk".)

③ Suppose S is an R -algebra (via φ) and $F_i \rightarrow M \rightarrow 0$ is a free resolution of M over R .

⊗ Recall $S \otimes_R F_i$ is indeed a complex of free S -modules where

$$S \otimes_R F_i \rightarrow S \otimes_R F_{i-1} \rightarrow S \otimes_R M \rightarrow 0 \text{ is exact.}$$

We have already seen an example where $S \otimes_R F_i$ is not free.

Q. When will it give a free resolution?

A. Tensoring with S should preserve exactness.

Defn.

(Flat modules) An R module K is flat if $K \otimes_R -$ is exact.

Eg. If $A \subset R$ is an m.c.s., take $K = R_A$.

Remark.

$K \otimes_R -$ being exact says that s.e.s.es go to s.e.s.es.

However, if C_\bullet is any exact complex of R -modules, then above implies $K \otimes_R C_\bullet$ is exact.

(Reason: Every long exact sequence (l.e.s.) breaks up into

s.e.ses.)

Thus, if S is an R -algebra which is flat as an R -module, then $S \otimes_R F. \rightarrow S \otimes_R M \rightarrow 0$ is an S -free resolution of $S \otimes_R M$.

(b) What about minimality?

Suppose $F. \rightarrow M$ is a minimal R -free resolution.

Q. Does $S \otimes_R F.$ give a minimal S -free resolution of $S \otimes_R M$?

To talk about that, we assume that S is Noetherian and local.

Let $(S, \mathfrak{m}, \mathfrak{l})$ be local, $\varphi: R \rightarrow S$ a ring map, and $F. \rightarrow M \rightarrow 0$ a min'l R -free resolution of M .

Is $S \otimes_R F. \rightarrow S \otimes_R M \rightarrow 0$ a min'l free resolution of $S \otimes_R M$ over S ?

Note that: ① Just flatness is not enough. (Find an example.)

(Take $R = \mathbb{Q}[x]$ and $S = \text{field of fracs.} = \mathbb{Q}(x)$.)

② If $\varphi(\mathfrak{m}) \subset \mathfrak{m}$, then the entries of the matrices in $S \otimes_R F.$ lie in \mathfrak{m} and hence, $S \otimes_R F. \rightarrow S \otimes_R M \rightarrow 0$ is a min'l free resolution.

Thus, we need that $\varphi: R \rightarrow S$ is a flat, local map.

Def'n. (Flat, local map) A ring map $\varphi: R \rightarrow S$ is flat if S is a flat R -module (via φ) and local if (S, \mathfrak{m}) is local with $\varphi(\mathfrak{m}) \subset \mathfrak{m}$.

Example If $A \subset R$ is an m.c.s., then $S = R_A$ is a flat R -algebra, local if $A = R \setminus \mathfrak{p}$ for some $\mathfrak{p} \in \text{Spec}(R)$ but may not satisfy $\varphi_{\mathfrak{p}}: R \rightarrow R_A$ is local.

Ex. Suppose $\varphi: R \rightarrow S$ is flat.

Suppose $F_\bullet \rightarrow M$ is an R -free resolution. Suppose $S \otimes_R F_\bullet \rightarrow S \otimes_R M$ is a min'l S -free resolution. Then, $F_\bullet \rightarrow M$ was minimal.

Example There is the notion of "completion in the \mathfrak{m} -adic topology," denoted \hat{R} and the natural map $\varphi: R \rightarrow \hat{R}$ is flat, local.

Aside Standard reduction in many problems: WLOG, assume R is a complete local domain.

$$R \xrightarrow{\text{localise}} (R, \mathfrak{m}) \xrightarrow{\text{complete}} (\hat{R}, \mathfrak{m}\hat{R}) \xrightarrow[\text{suitable prime}]{\text{go mod}} \text{complete local domain?}$$

Complexes and Homology

Defn ① Given complexes C_\bullet and D_\bullet , a map of complexes $\alpha_\bullet: C_\bullet \rightarrow D_\bullet$ is a collection of R -linear maps $\alpha_i: C_i \rightarrow D_i$ such that each sequence below commutes. Also called a chain map.

$$\begin{array}{ccccccc} \dots & \rightarrow & C_i & \xrightarrow{\partial^C} & C_{i-1} & \rightarrow & \dots \\ & & \alpha_i \downarrow & \wr & \downarrow \alpha_{i-1} & & \\ \dots & \rightarrow & D_i & \xrightarrow{\partial^D} & D_{i-1} & \rightarrow & \dots \end{array} \quad (\alpha \partial^C = \partial^D \alpha)$$

② C_\bullet and D_\bullet are isomorphic if $\exists \alpha_\bullet: C_\bullet \rightarrow D_\bullet$ s.t. $\alpha_i: C_i \rightarrow D_i$ is an isomorphism for all i .

③ Given $\alpha_\bullet, \beta_\bullet: C_\bullet \rightarrow D_\bullet$ chain maps, we say α_\bullet is homotopic to β_\bullet \exists maps $\delta_i: C_i \rightarrow D_{i+1}$ s.t.

$$\alpha_i - \beta_i = \delta_{i-1} \partial_i^C + \partial_{i+1}^D \delta_i$$

$$\begin{array}{ccccccc} \dots & \rightarrow & C_i & \xrightarrow{\partial^C} & C_{i-1} & \rightarrow & \dots \\ & \searrow \delta_i & \alpha_i \downarrow & \wr & \downarrow \beta_i & \swarrow \delta_{i-1} & \\ \dots & \rightarrow & D_{i+1} & \rightarrow & D_i & \rightarrow & \dots \end{array}$$

(Chain maps, homotopic maps)

Observations:

- ① (a) A chain map $\alpha_\bullet : C_\bullet \rightarrow D_\bullet$ induces a map $H_*(\alpha) : H_*(C_\bullet) \rightarrow H_*(D_\bullet)$, i.e., we get maps $H_*(\alpha_i) : H_i(C_\bullet) \rightarrow H_i(D_\bullet)$.

This follows since $\alpha_i(\ker(\partial_i^C)) \subset \ker(\partial_i^D)$ and $\alpha_i(\text{im}(\partial_{i+1}^C)) \subset \text{im}(\partial_{i+1}^D)$.

$$\begin{array}{ccc} \ker(\partial_i^C) & \xrightarrow{\alpha} & \ker(\partial_i^D) \\ \downarrow & & \downarrow \\ \frac{\ker(\partial_i^C)}{\text{im}(\partial_{i+1}^C)} = H_i(C) & \xrightarrow{\quad} & H_i(D) = \frac{\ker(\partial_i^D)}{\text{im}(\partial_{i+1}^D)} \end{array}$$

- ② If $\alpha_\bullet = \text{id}$, then $H_*(\alpha_\bullet) = \text{id}$.

- ③ If $C_\bullet \xrightarrow{\alpha_\bullet} C'_\bullet \xrightarrow{\beta_\bullet} C''_\bullet$ are chain maps, then $C_\bullet \xrightarrow{\beta_\bullet \circ \alpha_\bullet} C''_\bullet$ is a chain map and

$$\begin{array}{ccc} H_*(C_\bullet) & \xrightarrow{H_*(\alpha_\bullet)} & H_*(C'_\bullet) \\ & \searrow H_*(\beta_\bullet \circ \alpha_\bullet) & \downarrow H_*(\beta_\bullet) \\ & & H_*(C''_\bullet) \end{array} \quad \text{Commutative.}$$

- ④ If α_\bullet and β_\bullet are homotopic, then $H_*(\alpha_\bullet) = H_*(\beta_\bullet)$, i.e., homotopic chain maps induce the same maps on homology.

Lecture 14 (09-02-2021)

09 February 2021 11:35

Remark H_* is a functor from the category of chain complexes to chain complexes. We think of $H_*(C_*)$ as a chain complex with zero differential (i.e., all maps are the zero maps).

$$\dots \rightarrow H_n(C_*) \xrightarrow{0} H_{n-1}(C_*) \rightarrow \dots$$

Putting 0 maps ensures that homology is again $H_*(C_*)$.

Prop' If α_* and β_* are homotopic, then $H_*(\alpha_*) = H_*(\beta_*)$.

Proof. Denote $H_i(\alpha_*)$ as $\alpha_{i*} : H_i(C_*) \rightarrow H_i(D_*)$.

Claim. $\forall i : \alpha_{i*} = \beta_{i*}$.

$$\begin{array}{ccccccc} \dots & \rightarrow & C_{i+1} & \xrightarrow{\partial^C} & C_i & \xrightarrow{\partial^C} & C_{i-1} \rightarrow \dots \\ & & \searrow \gamma_i & & \downarrow \alpha_i & & \downarrow \beta_i \\ & & & & & & \downarrow \gamma_{i-1} \\ \dots & \rightarrow & D_{i+1} & \xrightarrow{\partial^D} & D_i & \xrightarrow{\partial^D} & D_{i-1} \rightarrow \dots \end{array}$$

$$\alpha_i - \beta_i = \gamma_{i-1} \partial^C + \partial^D \gamma_i$$

Notation: $Z_i(C) = \ker \partial_i^C$, $B_i(C) = \text{im } \partial_{i+1}^C$, $H_i(C) = Z_i(C) / B_i(C)$.

Let $x \in Z_i$. Consider $\bar{x} \in H_i(C)$. Want: $\alpha_{i*}(\bar{x}) = \beta_{i*}(\bar{x})$.

That is want to show

$$\alpha_i(x) - \beta_i(x) \in B_i(D).$$

However,

$$\begin{aligned} (\alpha_i - \beta_i)(x) &= \gamma_{i-1} \partial^C(x) + \partial^D \gamma_i(x) \\ &= \partial^D \gamma_i(x) \quad \stackrel{=0}{=} \text{since } x \in Z_i \in \ker \partial^C \\ &\in \text{im } \partial_{i+1}^D = B_i(D). \quad \square \end{aligned}$$

③ Let C and D be complexes. We say that C is homotopic to D (denoted $C \simeq D$) if \exists chain maps $\alpha: C \rightarrow D$, $\beta: D \rightarrow C$ such that $\alpha \circ \beta \simeq \text{id}_D$ and $\beta \circ \alpha \simeq \text{id}_C$.

(homotopic complexes)

If $C \simeq D$, then $H(C) = H(D)$.

Q. If C and D are homotopic, are they isomorphic?

No.

Comment.

Let $F: \dots \rightarrow F_2 \rightarrow F_1 \xrightarrow{\varphi_1} F_0 \rightarrow 0$ be a free resolution of M .
 Think of M as a complex $\dots \rightarrow 0 \rightarrow 0 \rightarrow M \rightarrow 0$
 as a complex $\dots \rightarrow 0 \rightarrow 0 \rightarrow M \rightarrow 0$
 $M = \text{ker } \varphi_0$

(See "derived categories".)

Koszul complex

$$\text{Boundary} \left(\begin{array}{c} e_1 \\ \triangle \\ e_2 \quad e_3 \end{array} \right) = e_1 e_2 + e_2 e_3 + e_3 e_1 \quad (\text{formal sum})$$

$$= e_1 e_2 + e_2 e_3 - e_1 e_3$$

$$\text{Boundary} (e_1, e_2, e_3)$$

$$\partial(e_1 \wedge e_2 \wedge e_3) = e_2 \wedge e_3 - e_1 \wedge e_3 + e_1 \wedge e_2$$

$\uparrow \qquad \qquad \uparrow \qquad \qquad \uparrow$
 excluding e_1 excl. e_2 exc. e_3

$$\partial(e_2 \wedge e_3) = e_3 - e_2$$

$$\partial(e_1 \wedge e_3) = e_3 - e_1$$

$$\partial(e_1 \wedge e_2) = e_2 - e_1$$

$$\begin{aligned} \therefore \partial^2(e_1 \wedge e_2 \wedge e_3) &= (e_3 - e_2) - (e_3 - e_1) + (e_2 - e_1) \\ &= 0 \end{aligned}$$

$$0 \rightarrow \mathbb{Z}(e_1 \wedge e_2 \wedge e_3) \xrightarrow{\partial} \begin{array}{c} \mathbb{Z}(e_2 \wedge e_3) \\ \oplus \\ \mathbb{Z}(e_1 \wedge e_3) \\ \oplus \\ \mathbb{Z}(e_1 \wedge e_2) \end{array} \xrightarrow{\partial} \mathbb{Z}e_1 \oplus \mathbb{Z}e_2 \oplus \mathbb{Z}e_3 \rightarrow 0$$

A different boundary map:

Let $\varphi: \mathbb{Z}e_1 \oplus \mathbb{Z}e_2 \oplus \mathbb{Z}e_3 \rightarrow \mathbb{Z}$ be a map.

Define

$$\partial(e_1 \wedge e_2 \wedge e_3) = \varphi(e_1)e_2 \wedge e_3 - \varphi(e_2)e_1 \wedge e_3 + \varphi(e_3)e_1 \wedge e_2.$$

Koszul complex on $a_1, \dots, a_n \in R$. Define $M = Re_1 \oplus \dots \oplus Re_n$


and $\varphi: M \rightarrow R$ by $e_i \mapsto a_i$.

Write $\wedge^k M$ as:

$$0 \rightarrow \wedge^n M \xrightarrow{\partial} \wedge^{n-1} M \rightarrow \dots \rightarrow \wedge^2 M \xrightarrow{\partial} \wedge^1 M \xrightarrow{\varphi} R \rightarrow 0$$

where

$$\partial(e_{i_1} \wedge \dots \wedge e_{i_k}) = \sum_{1 \leq j \leq k} (-1)^{j-1} \varphi(e_{i_j}) e_{i_1} \wedge \dots \wedge \tilde{e}_{i_j} \wedge \dots \wedge e_{i_k}$$



Verify that the above is a complex.

Assignment. ① Explicitly write down the modules in $K_*(\mathfrak{a})$.

$$\mathfrak{a} = a_1, \dots, a_n.$$

② Identify $H_j(K_*(\mathfrak{a}))$ at least when $j = 0, 1, n$.

③ For $n = 1, 2, 3$ try many examples.

④ Is $\sum_r I_r$ a free resolution of $H_0(K_*(\mathfrak{a}))$?

[When is it?]
↳ when exact

Lecture 15 (11-02-2021)

11 February 2021 09:40

(Direct sums and tensor products of complexes)

Q. We had discussed "chain complexes" as a category.
Can you think of direct sums and tensor products (of two complexes)?

a) Can verify direct sum works in natural way. Define

$$(C \oplus C')_i = C_i \oplus C'_i$$

$$\downarrow \partial = (\partial_i^C, \partial_i^{C'})$$

$$(C \oplus C')_{i-1} = C_{i-1} \oplus C'_{i-1}$$

and verify $\partial^2 = 0$.

b) $(C \otimes C')_i = C_i \otimes C'_i$

$$\downarrow \partial = \partial^C \otimes \partial^{C'}$$

$$(C \otimes C')_{i-1} = C_{i-1} \otimes C'_{i-1}$$

This does define a chain complex but we won't consider this as the tensor product! (Informal: Note that $\partial^2 = 0$ above is because we get $\partial^2 = (\partial^C)^2 \otimes (\partial^{C'})^2 = 0 \otimes 0 = 0$. However just one component is enough for \otimes to be 0. This is "overskill".)

Example: $C : \dots \rightarrow C_n \rightarrow C_{n-1} \rightarrow \dots \rightarrow C_1 \rightarrow C_0 \rightarrow 0$ R -complex

$C' : \dots \rightarrow 0 \rightarrow S \rightarrow 0$ $S = R$ -alg.

The above product gives: $0 \rightarrow S \otimes C_0 \rightarrow 0$

Would want something like:

$$\rightarrow S \otimes C_n \rightarrow \dots \rightarrow S \otimes C_1 \rightarrow S \otimes C_0 \rightarrow 0.$$

Q. Define the complex $(C \otimes C')$. Identify a "good" differential.

$$\bigoplus_{i=0}^n C_i \otimes C'_{n-i} ?$$

Snake lemma

Given two short exact sequences of R -modules with compatible maps between them

$$\begin{array}{ccccccccc} 0 & \longrightarrow & L_1 & \xrightarrow{\varphi_1} & M_1 & \xrightarrow{\psi_1} & N_1 & \longrightarrow & 0 \\ & & \alpha \downarrow & & \beta \downarrow & & \downarrow \gamma & & \\ 0 & \longrightarrow & L_2 & \xrightarrow{\varphi_2} & M_2 & \xrightarrow{\psi_2} & N_2 & \longrightarrow & 0, \end{array}$$

we get an exact sequence

$$0 \rightarrow \ker \alpha \xrightarrow{\varphi_1} \ker \beta \xrightarrow{\psi_1} \ker \gamma \xrightarrow{\delta} \operatorname{coker} \alpha \xrightarrow{\overline{\varphi_2}} \operatorname{coker} \beta \xrightarrow{\overline{\psi_2}} \operatorname{coker} \gamma \rightarrow 0.$$

($\varphi_1, \psi_1, \overline{\varphi_2}, \overline{\psi_2}$ above are the induced maps.)

δ is called the **connecting homomorphism**.

[Tool: If $0 \rightarrow K \rightarrow 0$ is exact, then $K = 0$.]

Consequences:

- ① If two out of α, β, γ are isomorphisms, then so is the third.
- ② β injective $\Rightarrow \alpha$ injective
 β surjective $\Rightarrow \gamma$ surjective
- ③ β isomorphism $\Rightarrow \alpha$ injective, γ surjective AND $\ker \gamma \stackrel{\delta}{\cong} \operatorname{coker} \alpha$

Proofs φ is 1-1 $\Leftrightarrow \ker \varphi = 0$

φ is onto $\Leftrightarrow \operatorname{coker} \varphi = 0$

$0 \rightarrow A \rightarrow B \rightarrow 0$ is exact $\Leftrightarrow A \cong B$

Q. If N_1, N_2 are submodules of M such that $M/N_1 \cong M/N_2$,
 is $N_1 \cong N_2$?

Look at $0 \rightarrow N_1 \rightarrow M \rightarrow M/N_1 \rightarrow 0$

$$\begin{array}{ccccccc} \text{Look at} & 0 & \rightarrow & N_1 & \rightarrow & M & \rightarrow & M/N_1 & \rightarrow & 0 \\ & & & \downarrow & & \downarrow \text{id} & & \downarrow \cong & & \\ & 0 & \rightarrow & N_2 & \rightarrow & M & \rightarrow & M/N_2 & \rightarrow & 0 \end{array}$$

"Proof" Using the lemma, $N_1 \cong N_2$.

Attempt at Tensor product:

Consider D_i defined by

$$D_i = \bigoplus_{j=0}^i C_j \otimes C_{i-j}$$

①

$$D_i \xrightarrow{d} D_{i-1}$$

$$\partial(C_j \otimes C'_{i-j}) = \underbrace{(\partial C_j)}_{C_{j-1}} \otimes C'_{i-j} + C_j \otimes \underbrace{\partial C'_{i-j}}_{C'_{i-j-1}} \in D_{i-1}$$

$$\partial^2(C_j \otimes C'_{i-j}) = \partial^2 C_j \otimes C'_{i-j} + C_j \otimes \partial^2 C'_{i-j} \rightarrow \text{NOT ZERO!}$$

②

$$D_i \xrightarrow{d} D_{i-1}$$

$$\partial(C_j \otimes C'_{i-j}) = \underbrace{(\partial C_j)}_{C_{j-1}} \otimes C'_{i-j} - C_j \otimes \underbrace{\partial C'_{i-j}}_{C'_{i-j-1}} \in D_{i-1}$$

$$\partial^2(C_j \otimes C'_{i-j}) = -\partial^2 C_j \otimes C'_{i-j} - C_j \otimes \partial^2 C'_{i-j} \rightarrow \text{NOT ZERO!}$$

③

$$D_i \xrightarrow{d} D_{i-1}$$

$$\partial(C_j \otimes C'_{i-j}) = \underbrace{(\partial C_j)}_{C_{j-1}} \otimes C'_{i-j} + (-1)^j C_j \otimes \underbrace{\partial C'_{i-j}}_{C'_{i-j-1}} \in D_{i-1}$$

$$\partial^2 (c_j \otimes c_{i-j}) = (-1)^{i-1} \partial^c c_j \otimes \partial^c c_{i-j} + (-1)^i \partial^c c_j \otimes \partial^c c_{i-j} = 0.$$

③ works?

Lecture 16 (16-02-2021)

16 February 2021 11:35

Snake Lemma (a more general version)

Thm

Let

$$\begin{array}{ccccccc}
 & & \downarrow & & \downarrow & & \downarrow \\
 0 & \rightarrow & A_{n+1} & \xrightarrow{\varphi_{n+1}} & B_{n+1} & \xrightarrow{\psi_{n+1}} & C_{n+1} \rightarrow 0 \\
 & & \downarrow d_{n+1} & & \downarrow \varphi_{n+1} & & \downarrow \psi_{n+1} \\
 0 & \rightarrow & A_n & \xrightarrow{\varphi_n} & B_n & \xrightarrow{\psi_n} & C_n \xrightarrow{\exists \delta_n} 0 \\
 & & \downarrow d_n & & \downarrow \varphi_n & & \downarrow \psi_n \\
 0 & \rightarrow & \cdots \rightarrow A_{n-1} & \xrightarrow{\varphi_{n-1}} & B_{n-1} & \xrightarrow{\psi_{n-1}} & C_{n-1} \rightarrow 0 \\
 & & \downarrow d_{n-1} & & \downarrow \varphi_{n-1} & & \downarrow \psi_{n-1}
 \end{array}$$

\downarrow differentials
 \rightarrow rows exact
 all squares commute

be given. (δ_n is not actually a map $C_n \rightarrow A_{n-1}$. It will be on homology!)

Then, we have an exact sequence on homology:

$$\cdots \rightarrow H_{n+1}(C) \xrightarrow{\delta_{n+1}} H_n(A) \xrightarrow{\varphi_*} H_n(B) \xrightarrow{\psi_*} H_n(C) \xrightarrow{\delta_n} H_{n-1}(A) \rightarrow \cdots$$

δ is called the connecting homomorphism. (We already know $\exists \varphi_*, \psi_*$.)

That is, a short exact sequence of complexes induces a long exact sequence on homology.

Proof Step 1. Construction of δ_n .

Want: $\delta_n : H_n(C) \rightarrow H_{n-1}(A)$

Let $z \in \ker \psi_n$. (Then, $\exists \bar{z} \in H_n(C)$.)

Want to define: $\delta(\bar{z}) \in H_{n-1}(A)$.

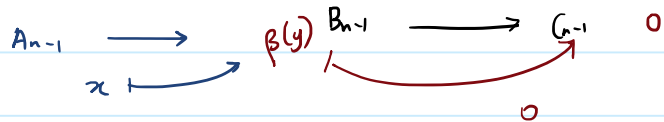
That is, we want $z \in \ker d$ s.t. $\bar{z} \in H_{n-1}(A)$ can be defined as $\delta(\bar{z})$.

$\exists y$ s.t. $\psi(y) = z$ $\nearrow \psi$ is onto

$$\begin{array}{ccc}
 y & \xrightarrow{\quad} & z \\
 \downarrow \psi & & \downarrow \psi \\
 B_{n-1} & \xrightarrow{\quad} & C_{n-1}
 \end{array}$$

$\exists y$ s.t. $\psi(y) = z$

$$\psi(\beta(y)) = \gamma(\psi(y)) = 0$$

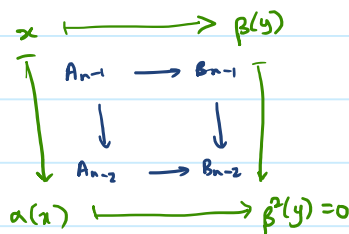


$\therefore \beta(y) \in \ker \psi = \text{im } \varphi$. $\therefore \exists! x \in A_{n-1}$ s.t. $\varphi(x) = \beta(y)$. φ is one-one

Want: $\bar{z} \mapsto \bar{x}$. Does \bar{x} even make sense?

- Check:
- ① $x \in \ker \alpha$
 - ② \bar{x} should not depend on (choice of) y
 - ③ same comment for z also.
- can then define $\delta(\bar{z}) = \bar{x}$

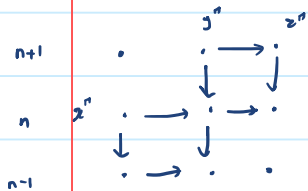
① $\varphi_{n-2}(\alpha(x)) = \beta(\varphi_{n-1}(x)) = \beta(\beta(y)) = 0$
 and hence, $\alpha(x) = 0$ since φ_{n-2} is 1-1.



② and ③: Given $z, z' \in \ker \gamma$ with $y, y' \in B_n$ s.t. $\psi(y) = z, \psi(y') = z'$ and $x, x' \in A_{n-1}$ s.t. $\varphi(x) = \beta(y)$ and $\varphi(x') = \beta(y')$. We show that if $\bar{z} = \bar{z}'$, then $\bar{x} = \bar{x}'$.

Proof. $z - z' = \gamma(z'')$ for some $z'' \in B_{n+1}$.
 $\exists y'' \in B_{n+1}$ s.t. $\psi(y'') = z''$

Now, $\psi(y - y') = z - z' = \gamma(z'') = \gamma(\psi(y''))$
 $= \psi(\beta(y''))$



$$\Rightarrow y - y' - \beta(y'') \in \ker \psi = \text{im } \psi$$

$$\Rightarrow y - y' - \beta(y'') = \psi(x'') \quad \text{for } x'' \in A_n.$$

$$\Rightarrow \beta(y) - \beta(y') - \beta^2(y'') = \beta \psi(x'')$$

ψ is H
 $\therefore \alpha(x'') = z - z'$

$$\Rightarrow \psi(x - x') = \beta \psi(x'') = \psi \alpha(x'')$$

$$\Rightarrow x - x' \in \text{im } \alpha, \quad \text{as desired. } \square$$

Thus, we can now define $\delta: H_n(C) \rightarrow H_{n-1}(A)$ as follows:

→ Take $z \in \ker \gamma$.

→ $\exists y \in B_n$ s.t. $\psi(y) = z$.

→ $\exists x \in A_{n-1}$ s.t. $\varphi(x) = \beta(y)$.

→ Define $\delta(z) = x$.

} This is well-defined.

To be completed. Need to verify that the long sequence is exact. Assignment!

$$\text{koszul on } (a_1) \hookrightarrow K(a_1, a_2) \hookrightarrow K(a_1, a_2, a_3) \dots$$

Lecture 17 (18-02-2021)

18 February 2021 09:02

One application of Snake Lemma is the following:

Let $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$ be a s.e.s. of R -modules and $F. \rightarrow L$ and $G. \rightarrow N$ be free resolutions.

$$\begin{array}{ccccccccc} 0 & \rightarrow & L & \rightarrow & M & \rightarrow & N & \rightarrow & 0 \\ & & \uparrow & & \uparrow & & \uparrow & & \\ & & F. & & F. \oplus G. & & G. & & \end{array}$$

We will prove that $F. \oplus G. \rightarrow M$ is a free resolution.
(Horse-shoe lemma)

Rewrite this as

$$(*) \quad \begin{array}{ccccccccc} & & \downarrow & & \downarrow & & \downarrow & & \\ 0 & \rightarrow & F_{n+1} & \rightarrow & F_{n+1} \oplus G_{n+1} & \rightarrow & G_{n+1} & \rightarrow & 0 \\ & & \downarrow & & \downarrow & & \downarrow & & \\ 0 & \rightarrow & F_n & \rightarrow & F_n \oplus G_n & \rightarrow & G_n & \rightarrow & 0 \\ & & \downarrow & & \downarrow & & \downarrow & & \end{array}$$

Beware: The maps are not always the usual ones!

1. In the proof of horseshoe lemma, it would be enough to prove $F. \oplus G. \rightarrow M$ is a complex. The left and right are exact, which will give exactness of middle. Thus, it would become a resolution.

2. Let F be an additive functor on R -modules.

Apply F on $(*)$. Then,

$$0 \rightarrow F(F.) \rightarrow F(F. \oplus G.) \rightarrow F(G.) \rightarrow 0$$

is a s.e.s. (in fact, split exact) since the rows were

split exact.

↳ part of the horse-shoe lemma

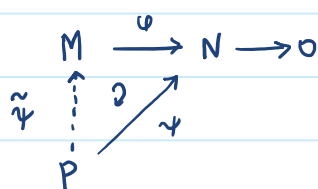
Then we can use Snake Lemma again.

Lecture 18 (22-02-2021)

22 February 2021 10:33

Projective modules : A module which has the lifting property of free modules.

Defⁿ. We say that R -module P is **projective** if given R -linear maps $\varphi : M \rightarrow N$ and $\psi : P \rightarrow N$ with φ onto, then $\exists \tilde{\psi} : P \rightarrow M$ which lifts ψ , i.e.,



$$\varphi \circ \tilde{\psi} = \psi.$$

Examples. (1) Every free module is projective.

Q. Is every projective module free?

Note that saying P is projective is equivalent to saying that $\text{Hom}_R(P, -)$ is exact.

Reason: Projective $\Leftrightarrow \text{Hom}_R(P, -) : \text{Hom}_R(P, M) \rightarrow \text{Hom}_R(P, N)$ is onto!
where $\varphi : M \rightarrow N$ is onto R -linear

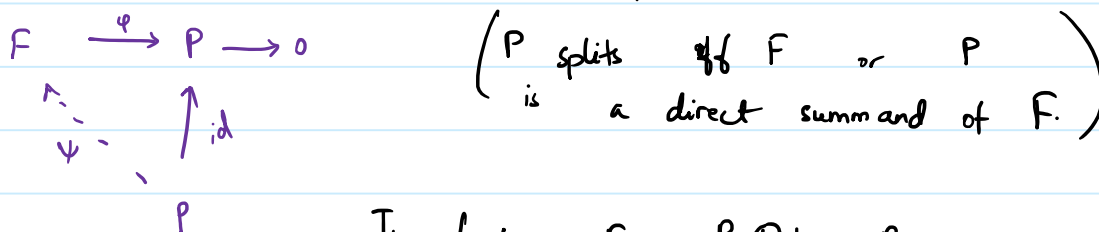
And we already knew $\text{Hom}_R(P, -)$ is always left exact.

Aside. Some definitions:

① An R module K is called **flat** if the functor $K \otimes_R -$ is exact.

② An R module K is called **injective** if $\text{Hom}_R(-, K)$ is exact.

Obs. Let P be projective, F a free R -module mapping onto P (via φ).
 $\exists \psi: P \rightarrow F$ s.t. $\varphi\psi = \text{id}_P$, i.e., $P \mid F$.



In fact, $F = P \oplus \ker \varphi$.

① P is a direct summand of a free module.

② We never used F was free. Thus, every onto map splits.
 In other words:

A s.e.s. of the form $0 \rightarrow L \rightarrow M \rightarrow P \rightarrow 0$ splits.

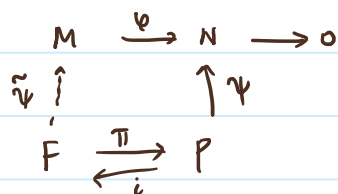
Q. Are the converses true?

① Suppose \exists a free module F and R -module Q s.t.
 $F = P \oplus Q$.

Is P projective?

Given $\varphi: M \rightarrow N$ onto and a map $\psi: P \rightarrow N$, does there exist $\tilde{\psi}: P \rightarrow M$ s.t. $\varphi\tilde{\psi} = \psi$?

First, we extend to $F \rightarrow N$ by composing φ with $\pi: F \rightarrow P$. Then, we can use lifting of F to get a map $F \rightarrow M$. We also have a natural $P \hookrightarrow F$. Verify it works.



In fact, the above shows that direct summands of projective modules are projective. It also shows Q is projective.

Also, the above shows that not all projective modules are free.
 Consider any direct summand of a free module which is not free. (E.g. $R = \mathbb{Z}/6\mathbb{Z} = F$. $F = \langle \bar{2} \rangle \oplus \langle \bar{3} \rangle$.)
 \downarrow not free but proj.

① + ①' : P is projective $\Leftrightarrow P$ is a direct summand of a free module

Q. Are direct sums of projective R -modules also proj.?

②' If every s.e.s. of the form $0 \rightarrow N \rightarrow M \rightarrow P \rightarrow 0$ splits, is P projective?

Yes! Take M to be a free R -module mapping onto P .
 Apply ①'.

We saw $\langle \bar{2} \rangle$ and $\langle \bar{3} \rangle$ are projective over $\mathbb{Z}/6\mathbb{Z}$ but not free.

More generally: If I, J are non-zero ideals s.t. $R = I \oplus J$, then I and J are projective but not free.

$$\begin{array}{c} \text{ann } I \supset J \\ \text{ann } J \supset I \end{array}$$

Q Suppose R is indecomposable as a module over itself.

Is every projective R -module free? **No! Find example(s)!**

Remarks ① Direct sum of projective modules are projective.

Let $\{P_i\}_{i \in I}$ be a family of proj. modules. Then,

$$F_i = P_i \oplus Q_i. \text{ Then,}$$

$$\bigoplus F_i = \left(\bigoplus P_i \right) \oplus \left(\bigoplus Q_i \right).$$

Thus, $\bigoplus P_i$ is proj., since $\bigoplus F_i$ is free.

② If P_1, P_2 are projective, then so is $P_1 \otimes_R P_2$.

$$F_1 = P_1 \oplus Q_1, \quad F_2 = P_2 \oplus Q_2, \quad \text{then}$$

$$F_1 \otimes F_2 = (P_1 \otimes P_2) \oplus (\) \oplus (\) \oplus (\)$$

\downarrow free \downarrow direct summand

③ Q. If P_1, P_2 are projective, is $\text{Hom}_R(P_1, P_2)$?

④ In general, quotients and submodules of projectives need not be projective. Find examples.

⑤ If $A \subset R$ is a m.r.s., P a proj. R -module, then
(Localisation) P_A is a projective R_A -module.

Same as before. Localising and \oplus commutes.

Localising free R -mod. gives free R_A -mod

⑥ If $I \subset R$ is an ideal, P is a projective R -module, then
(Quotient) P/IP is a projective R/I -module.

⑦ If $R \xrightarrow{\varphi} S$ is a ring map, P a projective R -module,
(Base change) then $S \otimes_R P$ is a projective S -module.

$$\begin{aligned}
 & F = P \oplus Q \\
 \hookrightarrow & S \otimes_R F = (S \otimes_R P) \oplus (S \otimes_R Q) \\
 \hookrightarrow & \text{free } S \qquad \qquad \qquad \hookrightarrow \therefore \text{this is projective}
 \end{aligned}$$

Proof works for ⑤ and ⑥ also!

Q. If $Q \subset P$ and $P, Q \rightarrow$ projective, is $Q|P$?

Nh. Take $R = P = \mathbb{Z}$ and $Q = 2\mathbb{Z}$.

Q. If P_p is a projective R_p -module for all $p \in \text{Spec } R$, is P a projective R -module?

Ans. Let $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$ be a s.e.s. of R -modules.
 Let P be s.t. P_p is proj. $\forall p \in \text{Spec } R$.
 (over R_p)

Consider

$$0 \rightarrow \text{Hom}_R(P, L) \xrightarrow{\psi^*} \text{Hom}_R(P, M) \xrightarrow{\psi^*} \text{Hom}_R(P, N) \rightarrow K \rightarrow 0$$

where $K = \text{coker } \psi^*$.

Since P_p is proj. over R_p , and

This is $\rightarrow (\text{Hom}_R(P, M))_p \cong \text{Hom}_{R_p}(P_p, M_p)$, we get $K_p = 0$
 not true in general. $\forall p$.

$$\therefore K = 0.$$

Thus, $\text{Hom}_R(P, -)$ is exact. $\therefore P$ is projective.
 (incomplete.)

⑧ (Projective module in local rings)

Let P be a f.g. projective R -module on a local ring (R, \mathfrak{m}, k) . Then, P is free. (consequence of NAK)

Proof. Let F be a free-module of rank $\mu(P)$. Then F maps onto P (minimally) which gives us $F = P \oplus K$.
 $\left(\begin{array}{l} P \text{ projective} \\ \downarrow \\ F \xrightarrow{\varphi} P \rightarrow 0 \text{ splits.} \end{array} \right)$

Here $K = \ker \varphi$ where $0 \rightarrow K \rightarrow F \xrightarrow{\varphi} P \rightarrow 0$.

Then,

$$F/\mathfrak{m}F \cong P/\mathfrak{m}P \oplus K/\mathfrak{m}K.$$

Both $F/\mathfrak{m}F$ and $P/\mathfrak{m}P$ have the same dim over k .

Thus, $K/\mathfrak{m}K = 0$ or $\mathfrak{m}K = K$.

By NAK, $K = 0$. \square

K is f.g. since $K \cong F/P$.

In fact, by Kaplansky, every projective module over (R, \mathfrak{m}, k) is free.

is free.

Next: (Schanuel's Lemma) Let P_1, P_2 be projective R -modules, K_1, K_2 and M R -modules s.t. we have the s.e.s.es

$$0 \rightarrow K_1 \rightarrow P_1 \rightarrow M \rightarrow 0 \quad \text{and}$$

$$0 \rightarrow K_2 \rightarrow P_2 \rightarrow M \rightarrow 0.$$

Then, $K_1 \oplus P_2 \cong K_2 \oplus P_1$.

Similar to free resolutions, one can define "projective resolutions".

Projective dimension of M as: $\text{pdr}(M) = \min$ length of proj. res'l of M , if it exists.

Lecture 19 (23-02-2021)

23 February 2021 11:36

Defⁿ. (Finitely presented) An R -module M is **finitely presented** if \exists free R -modules F_0 and F_1 of finite rank and an R -linear $\varphi: F_1 \rightarrow F_0$ s.t. $M \cong \text{coker } \varphi$.

$$(0 \rightarrow \varphi(F_1) \rightarrow F_0 \rightarrow M \rightarrow 0)$$

\downarrow f.g. \swarrow

φ can be thought of as a matrix, called a "presentation matrix" of M .

Remark If R is Noetherian, then M is f.g. $\Leftrightarrow M$ is finitely presented.

Ex. If M is finitely presented, N is an R -module, then for all $\mathfrak{p} \in \text{Spec } R$:

$$[\text{Hom}_R(M, N)]_{\mathfrak{p}} \cong_{R_{\mathfrak{p}}} \text{Hom}_{R_{\mathfrak{p}}}(M_{\mathfrak{p}}, N_{\mathfrak{p}}).$$

(With this assumption, the proof from yesterday goes through.)

((Proof of: $P_{\mathfrak{p}}$ proj. $\forall \mathfrak{p} \in \text{Spec } R \Rightarrow P$ proj.

If we assume P is f.p. then we are done.))

Q. If every localisation is free, then is the module free?

No. Take any f.g. proj. module which is not free. The localisation is f.g. + proj. over local ring. Thus, it is free.

Thm (Schanuel's Lemma) Let P_1, P_2 be projective R -modules, K_1, K_2 and M R -modules s.t. we have the s.e.s.es

$$0 \rightarrow K_1 \rightarrow P_1 \rightarrow M \rightarrow 0 \quad \text{and}$$

$$0 \rightarrow K_2 \rightarrow P_2 \rightarrow M \rightarrow 0.$$

Then, $K_1 \oplus P_2 \cong K_2 \oplus P_1$.

Proof.

Given:

$$\begin{array}{ccccccc}
 0 & \rightarrow & K_1 & \xrightarrow{\psi_1} & P_1 & \xrightarrow{\psi_1} & M \rightarrow 0 \\
 & & & & \uparrow \alpha & & \parallel id \\
 0 & \rightarrow & K_2 & \xrightarrow{\psi_2} & P_2 & \xrightarrow{\psi_2} & M \rightarrow 0
 \end{array}$$

Idea: Get an onto map $K_1 \oplus P_2 \rightarrow P_1$ and hope the kernel is K_2 .

Lift $id_M: M \rightarrow M$ to $\alpha: P_2 \rightarrow P_1$ using projectivity of P_2 .

Thus, $\psi_1 \alpha = \psi_2$. We already have $\psi_1: K_1 \rightarrow P_1$.

Thus, using the red maps, we get

$$\varphi: K_1 \oplus P_2 \rightarrow P_1 \quad \text{defined by}$$

$$\varphi(x, y) = \psi_1(x) + \alpha(y).$$

Is φ onto? Let $z \in P_1$. $\exists y \in P_2$ s.t. $\psi_2(y) = \psi_1(z)$.
(in M)

$$\text{But } \psi_2 = \psi_1 \alpha.$$

$$\therefore \psi_1 \alpha(y) = \psi_1(z) \quad \text{or} \quad \alpha(y) - z \in \ker \psi_1$$

" \parallel $\text{im } \psi_1$.

$$\therefore \alpha(y) - z = -\psi_1(x) \quad \text{for some } x \in K_1.$$

$$\Rightarrow z = \psi_1(x) + \alpha(y) = \varphi(x, y).$$

φ is onto. \square

Now, let $(x, y) \in \ker \varphi$. Then, $\psi_1(x) = -\alpha(y)$. ①
 $0 = -\psi_1 \alpha(y)$) apply ψ_1
 \parallel
 $\psi_2(y)$

$$\therefore y \in \ker \psi_2 = \text{im } \psi_2.$$

$$\text{Thus, } y = \psi_2(x') \quad \text{for some } x' \in K_2.$$

Put in ①: $\psi_1(x) = -\alpha \psi_2(x')$

In fact, α restricts to map $K_2 \rightarrow K_1$. We get:

$$\begin{array}{ccccccc}
 0 & \longrightarrow & K_1 & \xrightarrow{\psi_1} & P_1 & \xrightarrow{\psi_1} & M \longrightarrow 0 \\
 & & \alpha \uparrow & & \uparrow \alpha & & \parallel \text{id} \\
 0 & \longrightarrow & K_2 & \xrightarrow{\psi_2} & P_2 & \xrightarrow{\psi_2} & M \longrightarrow 0
 \end{array}$$

$\varphi(x) = -\alpha\psi_2(x') = -\varphi_1\alpha(x')$ and thus, injectivity gives
 $\alpha(x') = -x$.

Thus, $(x, y) = (-\alpha(x'), \psi_2(x'))$.

(Clearly, if $x' \in K_2$, then $(-\alpha(x'), \psi_2(x')) \in \ker \varphi$.)

Thus, the map

$$\begin{aligned}
 \psi: K_2 &\longrightarrow \ker \varphi \\
 x' &\longmapsto (-\alpha(x'), \psi_2(x'))
 \end{aligned}$$

is well-defined, onto, and one-one (since ψ_2 is 1-1).

$\therefore K_2 \cong \ker \varphi$.

We have $0 \rightarrow \ker \varphi \rightarrow K_1 \oplus P_2 \rightarrow P_1 \rightarrow 0$

Since P_1 is projective, the s.e.s. splits and we have

$$K_1 \oplus P_2 \cong P_1 \oplus \ker \varphi \cong P_1 \oplus K_2. \quad \square$$

Q. ① What information does Schanuel's lemma give about a finitely presented module?

② If every R -module M has a free-res'l of the form

$$0 \rightarrow R^n \rightarrow R^m \rightarrow M \rightarrow 0$$

what can you conclude about R ?

Next: ① (Lifting lemma) Let $P. \rightarrow M$ and $Q. \rightarrow N$ be projective resolutions and $f: M \rightarrow N$ be R -linear.

Then, \exists a chain map $\alpha. : P. \rightarrow Q.$ lifting f . $\alpha.$ is unique up to homotopy.

$$\begin{array}{ccccccc}
 0 & \leftarrow & M & \xleftarrow{\varphi_0} & P_0 & \xleftarrow{\varphi_1} & P_1 & \xleftarrow{\varphi_2} & P_2 & \leftarrow \dots \\
 & & f \downarrow & \wr & \downarrow \alpha_0 & \wr & \downarrow \alpha_1 & \wr & \downarrow \alpha_2 & \\
 0 & \leftarrow & N & \xleftarrow{\psi_0} & Q_0 & \xleftarrow{\psi_1} & P_1 & \xleftarrow{\psi_2} & P_2 & \leftarrow \dots
 \end{array}$$

defⁿ of
"lifts f"

commutes by defⁿ of chain map

(2) (Horseshoe Lemma) Gives proj. res'ls $P_i \rightarrow L$, $Q_i \rightarrow N$ and a seq.

$$0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0,$$

\exists a proj. res'l P'_i of M s.t. $(P'_i)_i = P_i \oplus Q_i$.

(Try to find an appropriate map $P_i \oplus Q_i \rightarrow P_i \oplus Q_i$.)

Lecture 20 (04-03-2021)

04 March 2021 09:30

Thm. Lifting Lemma: Let $P. \xrightarrow{\alpha_0} M \rightarrow 0$, $Q. \xrightarrow{\beta_0} N \rightarrow 0$ be projective resolutions and $f: M \rightarrow N$ be R -linear. Then, f can be lifted to a chain map $\gamma: P. \rightarrow Q.$, which is unique up to homotopy.

Proof. Want α_i s.t. $\dots \rightarrow P_2 \xrightarrow{\alpha_2} P_1 \xrightarrow{\alpha_1} P_0 \xrightarrow{\alpha_0} M \rightarrow 0$
 each square $\downarrow \gamma_2 \quad \downarrow \gamma_1 \quad \downarrow \gamma_0 \quad \downarrow f$
 commutes $\dots \rightarrow Q_2 \xrightarrow{\beta_2} Q_1 \xrightarrow{\beta_1} Q_0 \xrightarrow{\beta_0} N \rightarrow 0$

• γ_0 : we have $P_0 \xrightarrow{\alpha_0} M$ and $Q_0 \xrightarrow{\beta_0} N$.
 This gives γ_0 s.t. the rightmost square \square .

• γ_1 : Let $M_1 = \ker \alpha_0$ and $N_1 = \ker \beta_0$ and consider

$$\begin{array}{ccccccc} D & \rightarrow & M_1 & \rightarrow & P_0 & \rightarrow & M \rightarrow 0 \\ & & & & \downarrow \gamma_0 & \square & \downarrow f \\ 0 & \rightarrow & N_1 & \rightarrow & Q_0 & \rightarrow & N \rightarrow 0 \end{array}$$

$$\begin{array}{ccccccc} P_3 & \rightarrow & P_2 & \rightarrow & P_1 & \rightarrow & M_1 \rightarrow 0 \\ & & & & \downarrow \gamma_1 & \leftarrow \text{want} & \\ Q_3 & \rightarrow & Q_2 & \rightarrow & Q_1 & \rightarrow & N_1 \rightarrow 0 \end{array} \quad (1)$$

(Now, we want a map $M_1 \rightarrow N_1$. Then we can get a γ_1 .)
 Easy to check that $\gamma_0(M_1) \subset N_1$. Thus, γ_0 restricts to a map $\gamma_0|_{M_1}: M_1 \rightarrow N_1$.

Now, as earlier, we lift $\gamma_0|_{M_1}$ to γ_1 in (1).
 (Note that $Q_1 \xrightarrow{\beta_1} N_1$ is onto. $Q_1 \xrightarrow{\beta_1} Q_0$ need not have been.)

Thus, we get γ_1 as $\begin{array}{ccc} P_1 & \xrightarrow{\alpha_1} & M_1 \\ \downarrow \gamma_1 & \square & \downarrow \gamma_0|_{M_1} \\ Q_1 & \xrightarrow{\beta_1} & N_1 \end{array}$

$$\begin{array}{ccc} \gamma_1 \downarrow & \supseteq & \downarrow \gamma_0 \\ Q_1 & \xrightarrow{\beta_1} & M \end{array}$$

The above γ_1 also makes the following commute:

$$\begin{array}{ccc} P_1 & \xrightarrow{\alpha_1} & P_0 \\ \gamma_1 \downarrow & \supseteq & \downarrow \gamma_0 \\ Q_1 & \xrightarrow{\beta_1} & Q_0 \end{array}$$

- Ex. (a) Complete the above proof using induction.
 (b) What is the "best statement" of this proof?

Now, we prove uniqueness, up to homotopy.

Suppose $\gamma, \delta : P_0 \rightarrow Q_0$ are two liftings of f .

(That is, $f \alpha_0 = \beta_0 \gamma_0 = \beta_0 \delta_0$ and $\gamma_{i-1} \alpha_i = \beta_i \gamma_i$ AND $\delta_{i-1} \alpha_i = \beta_i \delta_i \quad \forall i \geq 1$.)

Claim: γ and δ are homotopic, i.e.

$$\begin{array}{ccccccc} \dots & \rightarrow & P_2 & \rightarrow & P_1 & \xrightarrow{\alpha_1} & P_0 & \xrightarrow{\alpha_0} & M & \rightarrow & 0 \\ & & & & \swarrow \delta_1 & \downarrow \gamma_1 & \downarrow \delta_1 & \swarrow \gamma_0 & \downarrow \delta_0 & & \downarrow f \\ & & & & \swarrow \delta_1 & \downarrow \gamma_1 & \downarrow \delta_1 & \swarrow \gamma_0 & \downarrow \delta_0 & & \downarrow f \\ \dots & \rightarrow & Q_2 & \xrightarrow{\beta_2} & Q_1 & \xrightarrow{\beta_1} & Q_0 & \rightarrow & N & \rightarrow & 0 \end{array}$$

• σ_0 : Note that $\beta_0 \circ (\gamma_0 - \delta_0) = 0$. Thus, $\text{im}(\gamma_0 - \delta_0) \subset \text{im } \beta_0$.

Moreover β_1 maps Q_1 onto $\text{im } \beta_0$.

Thus, projectivity of P_0 gives us σ_0 lifting $\gamma_0 - \delta_0$.

$$\begin{array}{ccc} & P_0 & \\ \sigma_0 \swarrow & & \downarrow \gamma_0 - \delta_0 \\ Q_0 & \xrightarrow{\beta_0} & Q \end{array} \quad \text{commutes}$$

Take $\sigma_{-1} = 0$ map. Thus,
 $\gamma_0 - \delta_0 = \beta_1 \sigma_0 + \sigma_{-1} \alpha_0.$

• σ_1 : Want $\sigma_1: P_1 \rightarrow Q_2$ s.t.

$$\beta_2 \sigma_1 + \sigma_0 \alpha_1 = \gamma_1 - \delta_1.$$

Or
$$\beta_2 \sigma_1 = \underbrace{\gamma_1 - \delta_1 - \sigma_0 \alpha_1}_{\text{Thus try to lift this!}}$$

Note that
$$\begin{aligned} \beta_1 \circ (\gamma_1 - \delta_1 - \sigma_0 \alpha_1) &= \gamma_0 \alpha_1 - \delta_0 \alpha_1 - \beta_1 \sigma_0 \alpha_1 \\ &= (\gamma_0 - \delta_0 - \beta_1 \sigma_0) \alpha_1 = \sigma_{-1} \alpha_0 \alpha_1 = 0. \end{aligned}$$

Thus, again $\text{im}(\gamma_1 - \delta_1 - \sigma_0 \alpha_1) \subset \text{im} \beta_2$ which gives us a lift $P_1 \rightarrow Q_2$.

Complete using induction. □

Two main applications:

Take $M = N$. (i) $f = \text{id}_M$, (ii) $f = \mu_\alpha$.

Then, we get maps $P \xrightarrow{\alpha} Q$ and $Q \xrightarrow{\beta} P$ which are lifts. $\alpha \circ \beta$ and id_Q both are lifts.

Thus, they are homotopic. Similarly, $\beta \circ \alpha \simeq \text{id}_P$.

$\therefore P \simeq Q$. (Homotopic)

An interesting consequence: If F is an additive functor of R -modules, then $F(P)$ and $F(Q)$ are also homotopic.

In particular, the homologies are isomorphic even if $F(P)$ and $F(Q)$ are not exact anymore.

This allows one to define "right derived functors" of a left exact additive functor, et cetera.

Lecture 21 (08-03-2021)

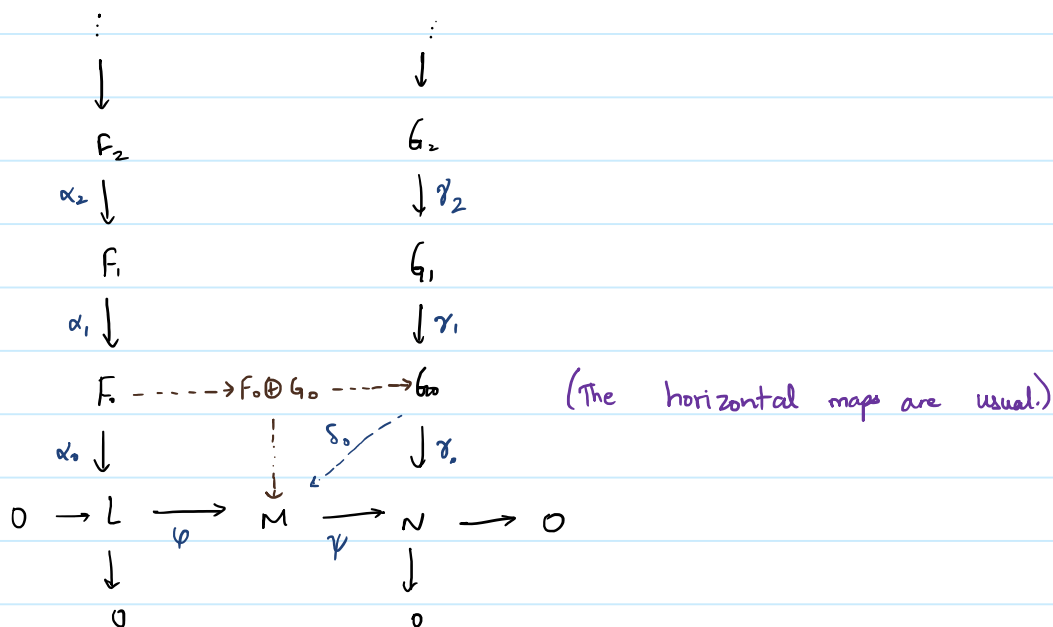
08 March 2021 10:40

Horseshoe lemma. Let F_i and G_i be projective resolutions of L and N . If

$$0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0 \text{ is a s.e.s.}$$

of R -modules, then there exists a resolution of M where the i^{th} projective module is $F_i \oplus G_i$.

Proof. Consider the diagram.



Claim 1. \exists an onto map $F_0 \oplus G_0 \rightarrow M$ making the diagram commute.

Lift $\gamma_0: G_0 \rightarrow N$ to $\delta_0: G_0 \rightarrow M$.

For $(x, z) \in F_0 \oplus G_0$ define

$$\beta_0(x, z) = \varphi \alpha_0(x) + \delta_0(z). \quad (\text{Clearly } R\text{-linear.})$$

Q. Is β_0 surjective?

Answer. Yes!

Let $y \in M$.

$$\exists z \in G_0 \text{ s.t. } \gamma_0(z) = \psi(y). \quad (\gamma_0 \text{ is onto!})$$

$$\text{Let } y' = \delta_0 z.$$

$$\begin{aligned} \text{Then, } \psi(y - y') &= \psi y - \psi \delta_0 z \\ &= \gamma_0(z) - \gamma_0(z) = 0. \end{aligned}$$

$$\therefore y - y' \in \ker \psi = \text{im } \varphi.$$

$$\Rightarrow y - y' = \varphi(x') \text{ for some } x' \in L.$$

$$\text{Since } \alpha_0 \text{ is onto, } \exists x \in F_0 \text{ s.t. } \alpha_0 x = x'.$$

$$\therefore y = \varphi(x') + y' = \varphi \alpha_0(x) + \delta_0(z) = \beta_0(x, z). \quad \square$$

Note: Could have also used first form of Snake Lemma for this \square .

$$0 \rightarrow F_1 \rightarrow F_1 \oplus G_1 \rightarrow G_1 \rightarrow 0$$

$$\downarrow \quad \downarrow \quad \downarrow \gamma_1$$

$$0 \rightarrow F_0 \rightarrow F_0 \oplus G_0 \rightarrow G_0 \rightarrow 0$$

$$\downarrow \beta_0 \quad \downarrow \delta_0 \quad \downarrow$$

$$0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$$

$$\downarrow \quad \downarrow \quad \downarrow$$

$$0 \quad 0 \quad 0$$

Claim 2. $\exists \beta_1 : F_1 \oplus G_1 \rightarrow F_0 \oplus G_0$
making the squares commute s.t.

$$\text{im } \beta_1 = \ker \beta_0.$$

Attempt 1. We know $\exists \delta_1 : G_1 \rightarrow F_0 \oplus G_0$.

In fact, we have the lift

$$\delta_1 = (0, \gamma_1)$$

But

$$\beta_0 \beta_1(x_1, z_1) = \varphi \alpha_0 \cdot \alpha_1(x_1) + \delta_0 \delta_1(z_1).$$

However, we only know $\delta_0 \delta_1(z_1) \in \ker \psi = \text{im } \varphi$,
not that it is 0.

Attempt 2. Observe that $\delta_0 \gamma_1(G_1) \subset \varphi(L)$ and

$$\varphi \alpha_0 : F_0 \rightarrow \varphi(L) \text{ is onto.}$$

Use projectivity of G_1 to get a lift

$$\begin{array}{ccc} & G & \\ \delta_1 \nearrow & \downarrow \delta_0 \gamma_1 & \\ F_0 & \rightarrow \varphi(L) & \rightarrow 0 \\ \varphi \alpha_0 \searrow & \uparrow \eta & \end{array}$$

$$\delta_1 : G_1 \rightarrow F_0$$

Now, define

$$M \quad \beta_1(x_1, z_1) = (\alpha_1(x_1) - \delta_1(z_1), \gamma_1(z_1)).$$

$$\begin{aligned} \text{Now, } \beta_0 \beta_1(x_1, z_1) &= \varphi \alpha_0(\alpha_1(x_1) - \delta_1(z_1)) + \delta_0(\gamma_1(z_1)) \\ &= 0 - \varphi \alpha_0 \delta_1(z_1) + \delta_0 \gamma_1(z_1) \\ &= 0 \end{aligned}$$

(In fact, we chose the sign $-$ precisely so that $\beta_0 \beta_1 = 0$.)

Now, we show $\text{im } \beta_1 \supseteq \text{ker } \beta_0$.

Let $(x, z) \in F_0 \oplus G_0$ be s.t. $\varphi \alpha_0(x) + \delta_0(z) = 0$.

Claim. $\exists (x_1, z_1) \in F_1 \oplus G_1$ s.t. $\beta_1(x_1, z_1) = (x, z)$.

Ex. (1) Prove $\text{ker } \beta_0 \subseteq \text{im } \beta_1$.

(2) Complete the proof of Horseshoe Lemma by induction.

Tor and Ext modules

Defⁿ. Let K be a fixed R -module, M an R -module with projective resolution F_\bullet .

$$F_\bullet : 0 \leftarrow F_0 \xleftarrow{\varphi_1} F_1 \xleftarrow{\varphi_2} \dots \quad \text{with} \quad H_0(F_\bullet) = M.$$

Then, $F_\bullet \otimes K : 0 \leftarrow F_0 \otimes K \xleftarrow{\varphi_1 \otimes \text{id}} F_1 \otimes K \xleftarrow{\varphi_2 \otimes \text{id}} F_2 \otimes K \leftarrow \dots$
may not be exact.

Define $\text{Tor}_i^R(M, K) := H_i(F_\bullet \otimes K)$.

Q. (1) Is it independent of the resolution chosen?

(Note that we didn't write $\text{Tor}_i^R(M, K, F_\bullet)$. We often suppress notations in homological alg.)

- ols.
- (0) $\text{Tor}_0^R(M, K) \cong M \otimes K$ since $-\otimes K$ is right exact.
- (1) If K is flat, we see that $\text{Tor}_i^R(M, K) = 0 \quad \forall i \geq 1$
- (2) (Answering Q) Suppose $f. \rightarrow M$ and $g. \rightarrow M$ are proj resolutions. Then, we had seen that $f.$ and $g.$ are homotopic. Then, so are $f. \otimes K$ and $g. \otimes K$.
 ($-\otimes K$ is additive.)
 Thus, $H_i(f. \otimes K) = H_i(g. \otimes K) \quad \forall i$.
 Therefore, there is no abuse of notation!
 That is, $\text{Tor}_i^R(M, K)$ depends only on M and K .
 (More precisely, $-\otimes K$.)

- Q.
- (2) If $\text{Tor}_i^R(M, K) = 0 \quad \forall i \geq 1 \quad \forall M$, then is K flat?
- (3) For a fixed M , what are the K s.t. $\text{Tor}_i^R(M, K) = 0 \quad \forall i \geq 1$.
 (fix n) $\forall i \geq n$.

(4) (Functoriality of Tor) Given $f: M \rightarrow N$ R linear, \exists an R -linear

$$f_i: \text{Tor}_i^R(M, K) \rightarrow \text{Tor}_i^R(N, K)$$

constructed from f such that $(f \circ g)_i = f_i \circ g_i \quad \forall i \geq 0$,
 $(\text{id}_M)_i = \text{id}_{\text{Tor}_i^R(M, K)} \quad \forall i \geq 0$,
 $D_i = 0 \quad \forall i \geq 0$,

$$(M=N \text{ and } a \in R) \quad (\mu a)_i = \mu a \quad \forall i \geq 0.$$

$$\mu a: M \rightarrow M \quad \hookrightarrow \mu a: \text{Tor}_i^R(M, K) \rightarrow \text{Tor}_i^R(M, K)$$

• Construction of f_i :

Given proj res'ls $f. \rightarrow M$ and $g. \rightarrow N$.

By the lifting lemma, we may lift $f: M \rightarrow N$ to a chain map $\gamma: f. \rightarrow g.$

This induces a chain map $\gamma \otimes K: f. \otimes K \rightarrow g. \otimes K$.

This, in turn, induces a map on the homologies, which were the Tor's.

the Tor's.

By lifting lemma, we also know that y is unique up to homotopy. Thus, f is well-defined.

Lecture 22 (09-03-2021)

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(5) If M is projective, then $\text{Tor}_i^R(M, K) = 0 \quad \forall i \geq 1, \forall K$.

Proof Choose the resolution

$$0 \rightarrow M \rightarrow M \rightarrow 0 \quad \text{for } M.$$

Then, if first co-ordinate is proj. or second is flat, we get 0.

(6) Let P be projective and L be such that

$$0 \rightarrow L \rightarrow P \rightarrow M \rightarrow 0 \quad \text{is a s.e.s.}$$

Then, $\text{Tor}_i^R(M, K) \cong ?$

Let $Q \rightarrow L \rightarrow 0$ be a proj. res'l of L/R .

Then, $Q \rightarrow P \rightarrow M$ is a proj. res'l of M/R .

$$\begin{array}{ccccccc} \cdots & \rightarrow & Q_2 & \rightarrow & Q_1 & \rightarrow & Q_0 \rightarrow L \rightarrow 0 \\ & & & & & & \\ \cdots & \rightarrow & Q_1 & \rightarrow & Q_0 & \rightarrow & P \rightarrow M \rightarrow 0 \end{array}$$

Then, $\text{Tor}_i^R(M, K) \cong \text{Tor}_{i-1}^R(L, K) \quad \text{for } i \geq 2.$

For $i=0$, $\text{Tor}_i^R(M, K) \cong M \otimes K.$

For $i=1$, $\text{Tor}_i^R(M, K) \cong ?$

(7) Let $0 \rightarrow L \xrightarrow{\varphi} M \xrightarrow{\psi} N \rightarrow 0$ be exact. What can one say about $\text{Tor}_i^R(L, K)$, $\text{Tor}_i^R(M, K)$, $\text{Tor}_i^R(N, K)$.

Let F and G be projective resolutions of L and N , resp.

Then, we construct a resolution H of M , where $H_i = F_i \oplus G_i$.

(by Horseshoe Lemma)

Tensor this with K to get:

$$\begin{array}{ccccccc}
& & \vdots & & \vdots & & \downarrow \\
0 & \rightarrow & F_{i+1} \otimes K & \rightarrow & F_{i+1} \otimes K \oplus G_{i+1} \otimes K & \rightarrow & G_{i+1} \otimes K \rightarrow 0 \\
& & \downarrow & & \downarrow & & \downarrow \\
0 & \rightarrow & F_i \otimes K & \rightarrow & F_i \otimes K \oplus G_i \otimes K & \rightarrow & G_i \otimes K \rightarrow 0 \\
& & \downarrow & & \downarrow & & \downarrow \\
0 & \rightarrow & F_{i-1} \otimes K & \rightarrow & F_{i-1} \otimes K \oplus G_{i-1} \otimes K & \rightarrow & G_{i-1} \otimes K \rightarrow 0 \\
& & \downarrow & & \downarrow & & \downarrow \\
& & \vdots & & \vdots & & \vdots
\end{array}$$

(Note rows were (and are) split exact.)

Apply Snake lemma to get a l.e.s. on homology:

$$\begin{array}{ccccccc}
\cdots & \rightarrow & \text{Tor}_{i+1}^R(N, K) & \rightarrow & \text{Tor}_i^R(L, K) & \rightarrow & \text{Tor}_i^R(M, K) \rightarrow \text{Tor}_i^R(N, K) \\
& & & & & & \downarrow \\
& & & & & & \cdots \leftarrow \text{Tor}_{i-1}^R(L, K) \\
& & & & & & \downarrow \\
& & & & & & \text{Tor}_i^R(M, K) \leftarrow \cdots \\
& & & & & & \downarrow \\
& & & & & & \text{Tor}_i^R(N, K) \rightarrow L \otimes K \rightarrow M \otimes K \rightarrow N \otimes K \rightarrow 0.
\end{array}$$

natural maps

Consequences:

(a) $0 \rightarrow L \rightarrow P \rightarrow M \rightarrow 0$ s.e.s. P projective

$$\Rightarrow \text{Tor}_i^R(M, K) \cong \begin{cases} \text{Tor}_{i-1}^R(L, K) & ; i \geq 2 \\ \ker(L \otimes K \rightarrow P \otimes K) & ; i = 1 \\ M \otimes K & ; i = 0. \end{cases}$$

(Since $\text{Tor}_i^R(P, K) = 0 \ \forall i \geq 1$.)

(b) If two out of L, M, N have the property that $\text{Tor}_i^R(-, K)$ is 0 for $i \gg 0$, then so does the third.

Lecture 23 (11-03-2021)

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General
Comment.

If F is a covariant, right-exact, additive functor on R -modules, then given an R -module M , we construct the "left derived functors" of F as follows:

(left derived functors)

- Take a projective resolution $P_\bullet \rightarrow M \rightarrow 0$ and define

$$L^i F(M) := H_i(F(P_\bullet)) \quad \forall i \geq 0$$

(One functor for each i .)

Project 2.
20th Nov

(a) Discuss properties. (Including what it does to maps as well as functoriality.)

Q: Does the "flatness" go through?

(b) What happens if F is contravariant left-exact?

Example: $\text{Hom}_R(-, K) \leftarrow$ what we get is $\text{Ext}_R^i(-, K)$.

Back to Tor:

- ⑧ Let (R, \mathfrak{m}, k) be a Noetherian local ring, M a f.g. R -module and $F_\bullet \rightarrow M \rightarrow 0$ a min'l free resolution of M .

Recall that the maps ψ_i in F_\bullet have entries in \mathfrak{m} . (The matrices.)

Thus, the maps in $F_\bullet \otimes k$ are 0. ($\because k = R/\mathfrak{m}$)

Thus, $H_i(F_\bullet \otimes k) \cong F_i \otimes k$.

Thus, $\text{Tor}_i^R(M, k) \cong k^{\beta_i(M)}$ or $\dim_k(\text{Tor}_i^R(M, k)) = \beta_i(M)$.

Consequence: (a) $\text{pdim}_R(M) = \max \{ i : \text{Tor}_i^R(M, k) \neq 0 \}$.

(b) [Two out of three] Given a s.e.s. $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$,

if two out of three have finite projective dimension, then so does the third.

(Consequence of (a) + l.e.s. of Tors.)

Ex.) Get bounds on $\text{pdim}_R(-)$ of the third.

Aside. Q. In the Horseshoe lemma, if the res'l of L and N are min'l, is the constructed res'l of M also min'l?
No. Take M s.t. $\text{pdim}_R(M) < \max(\text{pdim}_R(L), \text{pdim}_R(N))$.

Also, we have $\beta_i^R(M) \leq \beta_i^R(L) + \beta_i^R(N)$.

(And $\text{pdim}_R(M) \leq \max(\text{pdim}_R(L), \text{pdim}_R(N))$.)

(c) Assume symmetry of Tor: In particular, assume

$$\text{Tor}_i^R(M, K) \cong \text{Tor}_i^R(K, M).$$

Then, $\text{pdim}_R(M) \leq \text{pdim}_R(K)$.

In particular, if $\text{pdim}_R(K) < \infty$, then so is $\text{pdim}_R(M)$ for all f.g. M .

Proof. Compute $\text{Tor}_i^R(M, K)$ using a projective (free) resolution of K and use $\beta_i^R(M) = \dim_K(\text{Tor}_i^R(M, K))$.

Lecture 24 (15-03-2021)

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(Regular local ring)

Defⁿ. Let (R, \mathfrak{m}, k) be a Noetherian local ring. We say that R is a **regular local ring (RLR)** if $\text{pdim}_R(k) < \infty$.

• If R is an RLR, then \forall f.g. module M , then $\text{pdim}_R(M) < \infty$.

Ex. If R is an RLR, then so is $R_p \forall p \in \text{Spec } R$.

Q. If R is an RLR, and $a \in \mathfrak{m}$ is a non-zero-div, when is $\bar{R} = R/aR$ an RLR?

Back to Tor:

① Let R be a ring, N, M be R -modules.

Then, $\forall i \geq 0$, we have $\text{Tor}_i^R(M, k) \cong \text{Tor}_i^R(k, M)$.

(For $i=0$, we already know.)

Strategy: Prove it for free modules and use a free presentation of k (or M).

Proof Step I. Let M be free.

If we show that M is flat, then both sides are 0 $\forall K \forall i \geq 1$.

Now, a free module is flat since it is a direct sum of copies of R . Tensor product distributes over \oplus and R is flat as an R -module with $R \otimes N \cong N \forall N$.

Thus, $\text{Tor}_i^R(M, k) \cong \text{Tor}_i^R(k, M) \forall K \forall i \geq 0$.
($i=0$ is $k \otimes M$)

Step II. Let $0 \rightarrow L \rightarrow F \rightarrow M \rightarrow 0$ be a s.e.s. where F

is free and M is an arbitrary R -module
 Let $\dots \rightarrow G_2 \rightarrow G_1 \rightarrow G_0 \rightarrow K \rightarrow 0$ be a free resolution
 of K .

Tensor the first with G_i to get

$$\begin{array}{ccccccc}
 & & \downarrow & & \downarrow & & \downarrow \\
 0 & \rightarrow & G_2 \otimes L & \rightarrow & G_2 \otimes F & \rightarrow & G_2 \otimes M \rightarrow 0 \\
 & & \downarrow & & \downarrow & & \downarrow \\
 0 & \rightarrow & G_1 \otimes L & \rightarrow & G_1 \otimes F & \rightarrow & G_1 \otimes M \rightarrow 0 \\
 & & \downarrow & & \downarrow & & \downarrow \\
 0 & \rightarrow & G_0 \otimes L & \rightarrow & G_0 \otimes F & \rightarrow & G_0 \otimes M \rightarrow 0 \\
 & & \downarrow & & \downarrow & & \downarrow \\
 & & 0 & & 0 & & 0
 \end{array}$$

exact since G_i are free and hence flat

(Columns are complexes.)

By Snake Lemma, we get

$$\text{Tor}_i^R(K, F) \rightarrow \text{Tor}_i(K, M) \rightarrow K \otimes L \rightarrow K \otimes F \rightarrow K \otimes M \rightarrow 0$$

since F flat

Thus, $\text{Tor}_i(K, M) = \ker(K \otimes L \rightarrow K \otimes F) \cong \text{Tor}_i(M, K)$
 by consequence 7(a).

Furthermore, since $\text{Tor}_i^R(K, F) = 0 \quad \forall i \geq 1$, we get

$$\text{Tor}_i^R(K, M) \cong \text{Tor}_{i-1}^R(K, L) \quad \forall i \geq 2.$$

Since we know $\text{Tor}_i^R(M, K) \cong \text{Tor}_{i-1}^R(L, K) \quad \forall i \geq 2$,

we are done by induction. \square

(10) Let $a \in R$ be a non-zero-div. This gives a natural s.e.s

$$0 \rightarrow R \xrightarrow{a} R \rightarrow \bar{R} \rightarrow 0 \quad \text{which is } \bar{R} = R/aR$$

a free res'l. Given an R -module M , what is

$$\text{Tor}_i^R(M, \bar{R})?$$

$$\text{Tor}_i^R(M, \bar{R}) \cong \begin{cases} 0 & ; i \geq 2 \\ M/aM & ; i = 1 \\ 0 & ; i = 0 \end{cases}$$

$$\text{Tor}_i^R(M, \bar{R}) \cong \begin{cases} 0 & ; i \geq 2 \\ M/aM & ; i = 0 \\ \text{ann}_M(a) & ; i = 1 \\ \text{"} \\ \text{ker } \mu_a & \end{cases}$$

Special case: (a) if a is also a non-zero-divisor on M ,
i.e., $\text{ann}_M(a) = 0$, then $\text{Tor}_i^R(M, \bar{R}) = 0 \quad \forall i \geq 1$.

Consequence: Let a be a n.z.d on R and M
and $F. \rightarrow M$ be an R -free resolution of M .
Then, $0 = \text{Tor}_i^R(M, \bar{R}) \cong H_i(F. \otimes R) \quad \forall i \geq 1$.

$$\text{Thus, } H_i(F. \otimes R) = \begin{cases} \bar{M} & \text{for } i=0, \\ 0 & \text{for } i \geq 1. \end{cases}$$

($\bar{M} = M/aM$)

That is, $F. \otimes \bar{R}$ is a free-res'l of \bar{M} over \bar{R} .

In particular, if (R, \mathfrak{m}, k) is Noe. local, then

$$\beta_i^R(M) = \beta_i^{\bar{R}}(\bar{M}) \quad \forall i \geq 1, \text{ and}$$

$$\text{pdim}_R(M) = \text{pdim}_{\bar{R}}(\bar{M}).$$

Lecture 25 (16-03-2021)

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Recall. Here, assume that $a \in R$ is not a unit.

10. (a) If a is a nzd on R and M , and if $F. \rightarrow M$ is an R -free resolution, then $F. \otimes_R \bar{R}$ is a free res'l of \bar{M} .
In particular, $\text{pdim}_R(\bar{M}) < \text{pdim}_R(M)$.

(b) Furthermore, if (R, \mathfrak{m}, k) is Noetherian local, then

(i) $\beta_i^R(M) = \beta_i^{\bar{R}}(\bar{M})$, } Reason: If $F. \rightarrow M$ is min'l / R ,
(ii) $\text{pdim}_R(M) = \text{pdim}_{\bar{R}}(\bar{M})$, } then $F. \otimes_R \bar{R} \rightarrow \bar{M}$ is min'l / \bar{R} .

(iii) in particular,
 $\text{pdim}_{\bar{R}}(\bar{M}) < \infty \iff \text{pdim}_R(M)$.

(ii) Suppose a is a nzd on M . Then, what can we say about $\text{pdim}_R(\bar{M})$?

We have a natural s.e.s.

$$0 \rightarrow M \xrightarrow{a} M \rightarrow \bar{M} \rightarrow 0.$$

Given an R -module K , tensoring the s.e.s. with K gives a l.e.s. of Tors.

$$\begin{array}{ccccccc} \dots & \rightarrow & \text{Tor}_{j+1}^R(\bar{M}, K) & \rightarrow & \text{Tor}_j^R(M, K) & \xrightarrow{1 \otimes a} & \text{Tor}_j^R(M, K) & \rightarrow & \text{Tor}_j^R(\bar{M}, K) \\ & & & & & & & & \downarrow \\ & & & & & & & & \dots & \leftarrow & \text{Tor}_{j-1}^R(\bar{M}, K) \\ \hookrightarrow & \text{Tor}_j^R(\bar{M}, K) & \rightarrow & M \otimes K & \xrightarrow{1 \otimes a} & M \otimes K & \rightarrow & \bar{M} \otimes K & \rightarrow & 0. \end{array}$$

(a) If $\text{Tor}_j^R(M, K) = 0$ for $j \gg 0$, then same is true for \bar{M} .

In fact, if $\text{Tor}_j^R(M, K) = 0 \forall j > t$, then $\text{Tor}_j^R(M, K)$

$$\forall j > t+1.$$

Q. Suppose $\text{Tor}_i^R(\bar{M}, K) = 0 \quad \forall j \gg 0$, then is $\text{Tor}_j^R(M, K) \forall j \gg 0$?

(b) Suppose R is Noetherian, M, K are f.g. and $a \in \text{Jac}(R)$, then the above Q has a positive answer.

Proof. First observe that, with the given hypothesis, we can choose $F. \rightarrow M \rightarrow 0$ to be a free res'd where $\text{rank}(F_i) < \infty \forall i$. Hence, $F_i \otimes K$ is f.g. $\forall i$.

Thus, each $\text{Tor}_i^R(M, K)$ is f.g. (A subquotient.)

Now, we have

$$\text{Tor}_{i+1}^R(\bar{M}, K) = 0 \rightarrow \text{Tor}_i^R(M, K) \xrightarrow{\mu_a} \text{Tor}_i^R(M, K) \rightarrow 0 = \text{Tor}_i^R(M, K).$$

Use NAK to get $\text{Tor}_i^R(M, K) = 0$ for $i \gg 0$.

Thus, if $\text{Tor}_i^R(\bar{M}, K) = 0 \quad \forall i \geq t$, then $\text{Tor}_i^R(M, K) = 0 \quad \forall i \geq t$.

(c) Special case of (b): $(R, \mathfrak{m}, \mathbb{k})$ is Noetherian local, $a \in \mathfrak{m}$ is a nzd on $M \leftarrow$ f.g., $K = \mathbb{k}$.

Recall: $\text{Tor}_j^R(M, \mathbb{k})$ is a \mathbb{k} vector space. Hence, $\mu_a: \text{Tor}_j^R(M, \mathbb{k}) \rightarrow \text{Tor}_j^R(M, \mathbb{k})$ is the 0-map. (since $a \in \mathfrak{m}$.)

Thus, $\forall j: \text{Tor}_j(M, \mathbb{k}) \cong \text{Tor}_{j+t}^R(\bar{M}, \mathbb{k})$.

In particular, $\text{pdim}_R(M) < \infty \Leftrightarrow \text{pdim}_R(\bar{M}) < \infty$.

In this case, $\text{pdim}_R(M) + 1 = \text{pdim}_R(\bar{M})$.

General comments on associated primes:

Assume R is Noetherian. Then, $\mathfrak{p} \in \text{Spec } R$ is associated

to an R -module M if $R/\mathfrak{p} \hookrightarrow M$ as an R -module.

This naturally gives an s.e.s

$$0 \rightarrow R/\mathfrak{p} \hookrightarrow M \rightarrow N \rightarrow 0$$

which is useful in homological computations.

(Recall: $M \neq 0 \Leftrightarrow \text{Ass}(M) \neq \emptyset$.)

Moreover, every f.g. R -module M has a prime filtration, i.e.,

\exists a chain of submodules $0 = M_0 \subsetneq M_1 \subsetneq \dots \subsetneq M_n = M$, where $M_i/M_{i-1} \cong R/\mathfrak{p}_i$ for some $\mathfrak{p}_i \in \text{Spec}(R)$.

Any series as above gives a series of s.e.s es:

$$0 \rightarrow M_1 \hookrightarrow M_2 \rightarrow M_2/M_1 \rightarrow 0$$

\vdots

$$0 \rightarrow M_{n-1} \hookrightarrow M \rightarrow M/M_{n-1} \rightarrow 0$$

Observe: $M_1 \cong R/\mathfrak{p}_1$ and $M_i/M_{i-1} \cong R/\mathfrak{p}_i \quad \forall i \geq 2$.

Consequence: By the "boot strapping" process, prove that if R is Noetherian, K is an R -mod, then \forall f.g. M

$$\text{Tor}_i^R(M, K) = 0 \quad \text{if} \quad \text{Tor}_i^R(R/\mathfrak{p}, K) = 0$$

$\forall \mathfrak{p} \in \text{Spec}(R)$.

Recall: If R Noe, M f.g., then

① $\text{Ass}_R(M)$ is a finite set,

② $\bigcup_{\mathfrak{p} \in \text{Ass}_R(M)} \mathfrak{p} =$ set of zero-div. of M .

Finding a nzd on M is precisely an element in $R \setminus \bigcup_{\mathfrak{p} \in \text{Ass}_R(M)} \mathfrak{p}$.

Q. If (R, \mathfrak{m}, k) is Noeth. local, M f.g., when does
 $\exists a \in \mathfrak{m}$ which is a n.z.d. on $M \otimes R$?

$\mathfrak{m} \notin \text{Ass}(M) \cup \text{Ass}(R)$. (Prime avoidance.)

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Q. Let (R, \mathfrak{m}, k) be Noetherian local, $I \subsetneq R$.

When does \exists a minimal generator which is a n.z.d.?

• If r is such an element, then $r \notin \bigcup_{\mathfrak{p} \in \text{Ass}(R)} \mathfrak{p}$ and $r \in I \setminus \mathfrak{m}I$.

Thus, $r \notin \left(\bigcup_{\mathfrak{p} \in \text{Ass}(R)} \mathfrak{p} \right) \cup (\mathfrak{m}I)$. Can we use the general prime avoidance.

Thus, existence is guaranteed if $I \not\subseteq \mathfrak{p} \ \forall \mathfrak{p} \in \text{Ass}(R)$ and $I \not\subseteq \mathfrak{m}I$.

Ex. Let (R, \mathfrak{m}, k) be Noe. local, $I \subsetneq R$. If I contains a n.z.d., then it has a min'l gen. which is a n.z.d. (Prime avoidance + NAK.)

Regular sequences and depth

Def. Let R be a ring, M an R -module. We say that a sequence $\underline{a} = (a_1, \dots, a_n)$ is an M -regular sequence (or M -sequence) if

① a_1 is a n.z.d. on M .

② a_i is a n.z.d. on $\frac{M}{\langle a_1, \dots, a_{i-1} \rangle M}$ $\forall i \geq 2$ and

③ $\langle \underline{a} \rangle M \neq M$.

(in particular, no a_i is a unit.)

• If $a \in R$, then $\underline{a} = (a)$ is a regular sequence, then $\mu_a: M \rightarrow M$ is one-one but not onto.

If $M = R$, then a is a n.z.d. and non-unit.

(converse as well.)

Defⁿ Let $I \subset R$ be an ideal and M an R -module.
 The length of a maximal M -sequence on I , of maximum length, is called **depth** (I, M). (depth)

Q. Is this finite? Are there some maximal sequences which are finite?

Notation ① If $M = R$, $\text{depth}(I, R)$ is often denoted $\text{grade}(I)$.
 ② If (R, \mathfrak{m}, k) is local, then $\text{depth}(M) := \text{depth}(\mathfrak{m}, M)$.

Q. Let (R, \mathfrak{m}, k) be Noetherian local, M f.g. When is $\text{depth}(M) = 0$?

Obs. $\text{depth}(M) = 0 \Rightarrow \mathfrak{m} \subset \mathbb{Z}(M) = \bigcup_{\mathfrak{p} \in \text{Ass}(M)} \mathfrak{p}$ (Ass(M) < ∞, use prime avoidance)

↓

$\mathfrak{m} = \mathfrak{p}$ for some $\mathfrak{p} \in \text{Ass}(M) \Leftrightarrow \mathfrak{m} \subset \mathfrak{p}$ for some $\mathfrak{p} \in \text{Ass}(M)$

↓

$\mathfrak{m} \in \text{Ass}(M)$.

The converse is also true. $\mathfrak{m} \in \text{Ass}(M) \Rightarrow \mathfrak{m} \subset \mathbb{Z}(M)$
↓
no reg. sequence. □

Thus, $\text{depth}(M) = 0 \Leftrightarrow \mathfrak{m} \in \text{Ass}_R(M) \Leftrightarrow \mathbb{k} \hookrightarrow M$ as an R -module
 $\Leftrightarrow \exists x \in M$ s.t. $(0 :_R x) = \mathfrak{m}$.
 $\Leftrightarrow \text{ann}_M(\mathfrak{m}) = 0 :_M \mathfrak{m} \neq 0$

Defⁿ Given an R -module M , define $\text{soc}(M)$ (called the **socle of M**) as $\text{ann}_M(\mathfrak{m}) = 0 :_M \mathfrak{m} \cong \text{Hom}_R(k, M)$.
 ((R, \mathfrak{m}, k) local above)

Thus, $\text{depth}(M) = 0 \Leftrightarrow \mathfrak{m} \in \text{Ass}_R(M) \Leftrightarrow \text{soc}(M) \neq 0$

A. If (R, \mathfrak{m}, k) is local, M an Artinian R -module, then $\text{soc}(M) \neq 0$ (and thus, $\text{depth}(M) = 0$).

Q. Let $a \in \mathfrak{m}$ be M -regular. What can you say about depth of M/aM .

Take a reg. sequence in M/aM . Put a at beginning.

We get an M -sequence.

Thus if $\text{depth}(M/aM) < \infty$, then $\text{depth}(M) \geq \text{depth}(M/aM) + 1$.

Auslander-Buchsbaum Formula: Let (R, \mathfrak{m}, k) be Noetherian local, M f.g.

s.t. $\text{pdim}_R(M) < \infty$.

Then, $\text{pdim}_R(M) + \text{depth}(M) = \text{depth}(R)$.

What can we say about RLR and depth .
reg. $\overset{\uparrow}{\text{local ring}}$

Defⁿ. A Noetherian ring R is regular if $R_{\mathfrak{p}}$ is a RLR $\forall \mathfrak{p} \in \text{Spec}(R)$.

Remind project on Monday.

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- Ex.
- ① A permutation of a regular sequence need not be regular.
Find examples.
 - ② If R is Noetherian, and M an R -module, is $\text{depth}(I, M) < \infty$?
(WJK local help?)
 - ③ Let R be Noetherian local. Is $\text{depth}(I, R) \geq 1$, then
 - Ⓐ $\text{depth}(R) \geq 1$,
 - Ⓑ I contains a n.z.d. and hence, $\text{ann}_R(I) = 0$.
 - Ⓒ I has a min'l generator which is a n.z.d. by prime avoidance (modified).
 - ④ Let R be Noetherian local, $\text{depth}(I, R) = r$.
Then, does \exists a min'l gen. set a_1, \dots, a_r of I s.t.
 (a_1, \dots, a_r) is I -regular?

Tensor Product of Complexes

Let (C, ∂) and (D, δ) be complexes of R -modules.
We define the complex $(C \otimes D)$ as the complex with

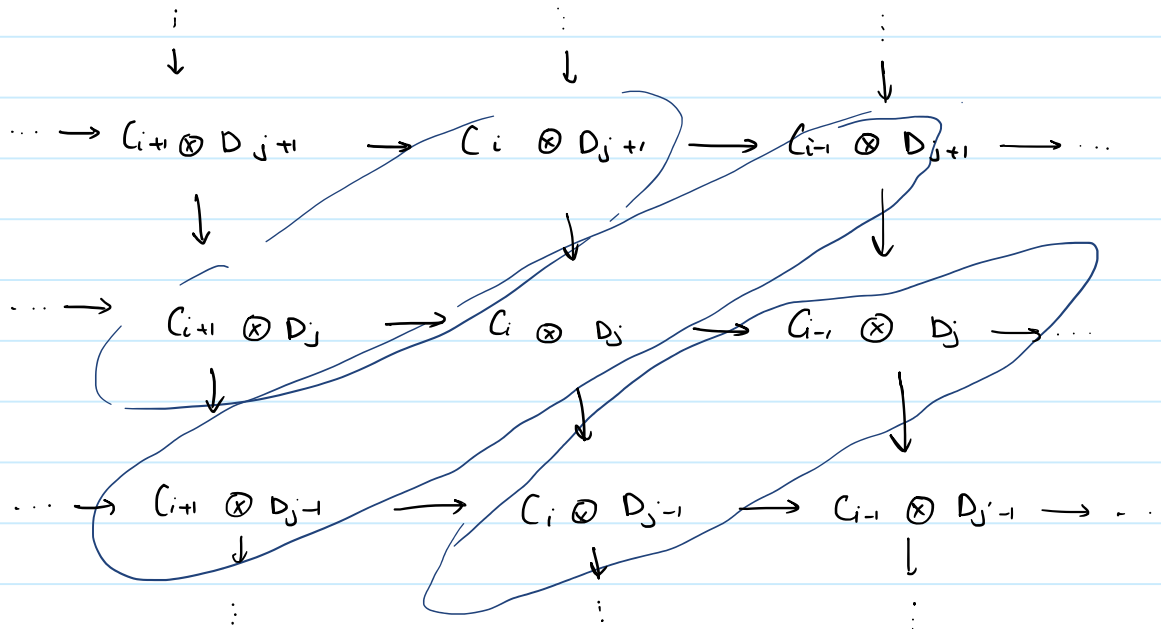
$$(C \otimes D)_n = \bigoplus_{i+j=n} C_i \otimes D_j \quad \text{with differential defined as:}$$

$$\text{for } x_i \in C_i, y \in D_j, \text{ set } x \otimes y \mapsto \begin{matrix} \partial x \otimes y & + & (-1)^i x \otimes \delta y \\ \cap & & \cap \\ C_{i-1} \otimes D_j & & C_i \otimes D_{j-1} \end{matrix}$$

Ex. Verify that $(C \otimes D)$ is a complex.

$$\begin{array}{ccccccc}
 \cdots & \rightarrow & C_{i+1} & \rightarrow & C_i & \rightarrow & C_{i-1} & \rightarrow & \cdots & \otimes & D_{j+1} \\
 & & & & & & & & & & \downarrow \\
 \cdots & \rightarrow & C_{i+1} & \rightarrow & C_i & \rightarrow & C_{i-1} & \rightarrow & \cdots & \otimes & D_j \\
 & & & & & & & & & & \downarrow \\
 \cdots & \rightarrow & C_{i+1} & \rightarrow & C_i & \rightarrow & C_{i-1} & \rightarrow & \cdots & \otimes & D_{j-1} \\
 & & & & & & & & & & \downarrow \\
 & & & & & & & & & & \vdots
 \end{array}$$

The above gives



Diagonal gives the modules in $C \otimes D$, called the total complex of the given double complex. Working through the double complex also gives a proof of symmetry of Tor.

Q. If $F \rightarrow M \rightarrow 0$, $G \rightarrow N \rightarrow 0$ are free resolutions, what is $H_0(F \otimes G)$. Does $F \otimes G$ give a free resolution of H_0 ? If not, identify conditions when it is.

Example Let (C, δ) be a complex and $D: 0 \rightarrow M \rightarrow 0$ where M is an R -module.

$$\begin{array}{ccc}
 \text{Then, } (C \otimes D)_n & = & C_n \otimes M \\
 \downarrow & & \downarrow \delta \otimes \text{id} \\
 C_n & & \cdots
 \end{array}$$

$$\begin{array}{ccc} \downarrow & \downarrow \delta \otimes 1 \otimes & \downarrow \\ (C. \otimes D.)_{n-1} & = C_{n-1} \otimes M & \delta x \otimes y \end{array}$$

Q: What is $H_n(C. \otimes D.)$? Not exactly $H_n(C.) \otimes M$ since $\delta x \otimes y = 0 \not\Rightarrow \delta x = 0$.

For example, if $C.$ is a resolution, then we get $Tors$, not necessarily 0.

Remark Note that if we took $D. \otimes C.$, the maps would have a sign.

② Let $(C., \delta.)$ be a complex and $D. = 0 \rightarrow R \xrightarrow{a} R \rightarrow 0$.

$$\begin{array}{ccc} \text{Then, } (C. \otimes D.)_n & = & C_n \otimes R \oplus C_{n-1} \otimes R \\ & & \begin{array}{ccc} \begin{array}{c} \downarrow \delta \otimes 1 \\ \delta x \otimes 1 + 0 \end{array} & & \begin{array}{c} \begin{array}{c} y \otimes e \\ \downarrow \\ \delta y \otimes e + (-1)^{n-1} y \otimes a \end{array} \end{array} \end{array} \\ (C. \otimes D.)_n & = & C_{n-1} \otimes R \oplus C_{n-2} \otimes R \end{array}$$

In general, identify image of $(x \otimes b, y \otimes ce)$.

$$(x \otimes b, y \otimes ce) = (\delta x \otimes b + (-1)^{n-1} y \otimes a, \delta y \otimes ce)$$

$$\text{or } (b \delta x + (-1)^n c a y, c \delta y)$$

Note that the second 6-ord. is just δ .

Notation Shift of a complex.

Given $(C., \delta.)$, define $\Sigma^{-1}C.$ to be the complex with no dules $(\Sigma^{-1}C.)_n = C_{n-1}$ with differential $(\Sigma^{-1}\delta.)_n = \delta_{n-1}$.

$$\begin{array}{ccccccc} C. & : & \dots & \rightarrow & C_{i+1} & \xrightarrow{\delta_{i+1}} & C_i & \xrightarrow{\delta_i} & C_{i-1} & \xrightarrow{\delta_{i-1}} & \dots \\ \Sigma^{-1}C. & : & \dots & \rightarrow & C_i & \xrightarrow{\delta_i} & C_{i-1} & \xrightarrow{\delta_{i-1}} & C_{i-2} & \xrightarrow{\delta_{i-2}} & \dots \end{array}$$

δ_i δ_{i-1} δ_{i-2}

More generally, $(\Sigma^j C)_i$ is defined.

Back to example: $(C \otimes D)_n = C_n \otimes R \oplus C_{n-1} \otimes R_{\neq 1} = C_n \oplus C_{n-1}$

Thus, we have $0 \rightarrow C \rightarrow C \otimes D \rightarrow \Sigma^{-1} C \rightarrow 0$
 (The maps on $C \otimes D$ are actually $\Sigma^{-1} \delta$.)

The above is a s.e.s of complexes which gives a l.e.s. on homologies as

$$\dots \rightarrow H_{n+1}(C \otimes D) \rightarrow H_{n+1}(\Sigma^{-1} C) \rightarrow H_n(C) \rightarrow H_n(C \otimes D) \rightarrow H_n(\Sigma^{-1} C) \rightarrow \dots$$

What is $H_i(\Sigma^{-1} C)$? What is the connecting homomorphism?

"

$$H_{i-1}(C)$$

$$\begin{array}{ccccccc} \vdots & & \vdots & & \vdots & & \\ \downarrow & & \downarrow & & \downarrow & & \\ 0 \rightarrow C_n & \rightarrow & C_n \oplus C_{n-1} & \rightarrow & C_{n-1} & \rightarrow & 0 \\ \delta \downarrow & & \downarrow & & \downarrow \delta & & \\ 0 \rightarrow C_{n-1} & \rightarrow & C_{n-1} \oplus C_{n-2} & \rightarrow & C_{n-2} & \rightarrow & 0 \\ \downarrow & & \downarrow & & \downarrow & & \\ \vdots & & \vdots & & \vdots & & \end{array}$$

middle map: $(x, y) \mapsto (\delta x + (-1)^{n-1} ay, \delta y)$

Verify the connecting homom. is $\pm \partial$.

Thus, the l.e.s. is

$$\dots \rightarrow H_{n+1}(C \otimes D) \rightarrow H_n(C) \xrightarrow{\pm \partial} H_n(C) \rightarrow H_n(C \otimes D) \rightarrow H_{n-1}(C) \xrightarrow{\pm \partial} \dots$$

Q. What can you conclude?

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$$\dots \rightarrow H_{n+1}(C \otimes D) \rightarrow H_n(C) \xrightarrow{\pm a} H_n(C) \rightarrow H_n(C \otimes D) \rightarrow H_{n-1}(C) \xrightarrow{\pm a} \dots$$

Conclusions:

① If $a \in \mathcal{U}(R)$, then $H_n(C \otimes D) = 0$.

(In fact, we only need each $\pm \mu_a$ to be an isomorphism on $H_n(C)$.)

② If $H_i(C) = 0 \forall i \geq n$, then $H_i(C \otimes D) = 0 \forall i \geq n+1$.

Moreover,, $H_n(C \otimes D) = \ker(\pm \mu_a : H_{n-1}(C) \rightarrow H_{n-1}(C))$.

In fact if $H_n(C) = H_{n-1}(C) = 0$, then $H_n(C \otimes D) = 0$.

③ $H_i(C \otimes D) = 0 = 0 \forall i \geq n \Rightarrow \mu_a : H_i(C) \rightarrow H_i(C)$ is an iso. for all $i \geq n$.

④ The converse of ② is not true. Find examples

⑤ If each C_i is Noetherian and $a \in \text{Jac}(R)$, then by NAK, if $H_n(C \otimes D) = 0$, then $H_n(C) = 0$. (since μ_a becomes on to.)

This acts as a converse to ③.

Special case of example ②:

$$C : 0 \rightarrow R_{e_1} \xrightarrow{a} R \rightarrow 0 ; D : 0 \rightarrow R_{e_2} \xrightarrow{a} R \rightarrow 0$$

What is $C \otimes D$?

$$\begin{aligned} 0 &\rightarrow R_{e_1} \otimes R_{e_2} \rightarrow R_{e_1} \oplus R_{e_2} \rightarrow R \rightarrow 0 \\ e_1 \otimes e_2 &\mapsto a_1 e_2 - a_2 e_1 \\ e_1 &\mapsto a_1 \\ e_2 &\mapsto a_2 \end{aligned}$$

One can now explicitly construct an isomorphism (of complexes) with $K(a_1, a_2)$.

↳ Koszul on two elements

More generally, by induction, one can show that if $a_1, \dots, a_n \in R$, then

$$K.(a_1) \otimes \dots \otimes K.(a_n) \cong K.(a_1, \dots, a_n).$$

[Show $K.(a_1, \dots, a_{n-1}) \otimes K.(a_n) \cong K.(a_1, \dots, a_n)$.]

$\begin{array}{ccc} \uparrow & & \uparrow \\ C_0 & & D_0 \end{array}$

By the observations in example ②, we can relate Koszul homology of (a_1, \dots, a_n) with Koszul homology of (a_1, \dots, a_{n-1}) and $\mathbb{H}a_n$.

Defⁿ Let $a_1, \dots, a_n \in R$, M an R -module. We define the Koszul complex on $\underline{a} = (a_1, \dots, a_n)$ w.r.t. M as

$$K.(\underline{a}, M) := K.(\underline{a}) \otimes_R M.$$

Special case: If S is an R -algebra via φ ; $a_1, \dots, a_n \in R$, then

$$K.(\underline{a}, R) \otimes_R S \cong K.(\varphi(\underline{a}), S).$$

Based on the observations following Example ②, we have the following result:

Thm. Let $a_1, \dots, a_n \in R$, M an R -module. Consider the following statements:

- ① a_1, \dots, a_n is M -regular.
- ② $H_i(K.(\underline{a}, M)) = 0 \quad \forall i \geq 1$
- ③ $H_1(\underline{a}, M) = 0$.

Then, ① \Rightarrow ② \Rightarrow ③. Moreover, if M is Noetherian and $a_1, \dots, a_n \in \text{Jac}(R)$, then ③ \Rightarrow ①.

Proof. Note that ① \Leftrightarrow ② \Leftrightarrow ③ for $n=1$. Observe that $K(a_1, \dots, a_n, M) \cong K(a_1, \dots, a_{n-1}, M) \otimes K(a_n)$ and a_1, \dots, a_n is M -regular $\Leftrightarrow a_1, \dots, a_{n-1}$ is M -regular and a_n is a n.z.d. on $M / \langle a_1, \dots, a_{n-1} \rangle M$.

These observations allow us to use induction.

- Assume ①. Then a_1, \dots, a_{n-1} is M -reg $\Rightarrow H_i(a_1, \dots, a_{n-1}, M) = 0 \quad \forall i \geq 1$

\downarrow

$$H_i(a_1, \dots, a_n, M) = 0 \quad \forall i \geq 2$$

Moreover, $H_1(a_1, \dots, a_n, M) = \ker(\mu_{a_n}: H_0(a_1, \dots, a_{n-1}, M) \rightarrow H_0(a_1, \dots, a_n, M))$
 Since a_n is a n.z.d. on $H_0(a_1, \dots, a_{n-1}, M) \cong \frac{M}{\langle a_1, \dots, a_{n-1} \rangle M}$, $H_1 = 0$.

Thus, ① \Rightarrow ②.

• ② \Rightarrow ③ obvious.

Assume M is Noetherian and $\underline{a} \subset \text{Jac}(R)$. Use induction for ③ \Rightarrow ①.

$H_1(\underline{a}, M) = 0 \Rightarrow \mu_{a_n}$ is injective on $H_0(a_1, \dots, a_{n-1}, M) \cong \frac{M}{\langle a_1, \dots, a_{n-1} \rangle M}$.

By obs. ①, $H_1(a_1, \dots, a_{n-1}, M) = 0$.

By induction, a_1, \dots, a_{n-1} is M -reg and a_n is a n.z.d. on $\frac{M}{\langle a_1, \dots, a_{n-1} \rangle M}$.
 $\therefore a_1, \dots, a_n$ is M -reg. □

Consequences:

① If a_1, \dots, a_n is R -regular, then $K_0(a_1, \dots, a_n)$ is a free resolution of $H_0 = \frac{R}{\langle a_1, \dots, a_n \rangle}$.

In particular, $H_i K_0(\underline{a}, M) \cong \text{Tor}_i^R(R/\langle \underline{a} \rangle, M)$.

② If (R, \mathfrak{m}, k) is Noetherian local and \mathfrak{m} is generated by a regular sequence, then $\text{pdim}_R(k) < \infty$ and hence, R is an RLR.
 Moreover, if $\mu(\mathfrak{m}) = n$, then $\text{pdim}_R(k) = n \geq \text{pdim}_R(M) \forall f.g. M$.

E.g. $k[[x_1, \dots, x_n]]$ is an RLR. What is $\beta_i^R(k)$? $\binom{n}{i}$.

③ If $\underline{a} \subset \text{Jac}(R)$, then any permutation of \underline{a} is regular.

Lecture 29 (25-03-2021)

25 March 2021 09:34

Observations.

① Let $a_1, \dots, a_n \in R$ and $I = \langle a_1, \dots, a_n \rangle$ (and $\underline{a} = a_1, \dots, a_n$).

Recall : $H_0(\underline{a}, M) = M/I_M$ and $H_n(\underline{a}, M) = \text{ann}_M(I)$.

$$\begin{array}{ccc} \partial \rightarrow & R & \rightarrow R \\ & \otimes^M & \otimes^M \\ & & \otimes^n \\ \rightarrow M & \rightarrow & M^n \end{array}$$

In both cases, the homology is annihilated by I .

In fact, we now show that $\forall i, I \cdot H_i(\underline{a}, M) = 0$, i.e., I annihilates every homology.

Proof.

Let $S := R[T_1, \dots, T_n]$. Then, R is an S -algebra via $\varphi: S \rightarrow R$ is defined by $T_i \mapsto a_i$ (and id_R on R).

Let M be an R -module. Then, M is an S -module via φ .

$$K^R(\underline{a}; M) \cong K^S(\underline{T}; M) \text{ as } S\text{-complexes.}$$

Note T_1, \dots, T_n is a regular sequence (in S) and hence,

$$H_i^S(\underline{T}; M) \cong \text{Tor}_i^S(S/\langle \underline{T} \rangle, M)$$

Fact: If M, N are S -modules, then $\forall i$
 $\text{ann}_S(M) + \text{ann}_S(N) \subset \text{ann}_S(\text{Tor}_i^S(M, N)).$ (Ex.)

$$\text{Thus, } \text{ann}_S(H_i^S(\underline{T}; M)) \supset \underset{\langle T_1, \dots, T_n \rangle}{\text{ann}_S(S/\langle \underline{T} \rangle)} + \underset{\text{ann}_S(R) = \ker(\varphi)}{\text{ann}_S(M)}$$

$$\text{Thus, } \text{ann}_S(H_i^S(\underline{I}; M)) \supseteq \langle T_1, \dots, T_n \rangle + \langle T_1 - a_1, \dots, T_n - a_n \rangle$$

$$\cong \langle a_1, \dots, a_n \rangle = I$$

ideal \vec{a} gen. in R

This forces $\text{ann}_R(H_i^S(I, M)) \supseteq I$. But $H_i^S(I, M) \cong H_i^R(\vec{a}, M)$. \square

(2) Let (R, \mathfrak{m}, k) be local. Let $\{a_1, \dots, a_n\}$ be a min'l gen. set of I .
 If $I = \langle b_1, \dots, b_n \rangle$, then is $H_i(\vec{a}) \cong H_i(\vec{b})$? $\begin{bmatrix} a_1 \\ \vdots \\ a_n \end{bmatrix} = A \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$

$$\begin{array}{ccccc} \rightarrow & \bigoplus_{i=1}^n R e_i & \rightarrow & R & \rightarrow R/\langle \vec{a} \rangle \rightarrow 0 \\ & \downarrow A^T & & \downarrow \text{ide} & \\ & \bigoplus_{i=1}^n R f_i & \rightarrow & R & \rightarrow R/\langle \vec{b} \rangle \rightarrow 0 \end{array}$$

$e_i \mapsto a_i$ $f_i \mapsto b_i$

As in the proof of the splitting lemma, A^T is invertible.

$$\bigoplus R e_i = K_i \xrightarrow{\varphi} \bigoplus R f_i = K'_i$$

Focus now on:

Given φ , we construct $\bigwedge^i \varphi: \bigwedge^i K_i \rightarrow \bigwedge^i K'_i$

K_i K'_i

modules in Koszul complex

Suppose φ is given by $A = [a_{ij}]$.

Then, $e_1 \wedge e_2 \mapsto (a_{11} f_1 + a_{21} f_2 + \dots + a_{m1} f_n) \wedge (a_{12} f_1 + a_{22} f_2 + \dots + a_{n2} f_n)$

$\sum_{1 \leq i < j \leq n} (a_{ii} a_{jj} - a_{ij} a_{ji}) (f_i \wedge f_j)$

\hookrightarrow 2×2 minor formed by rows i & j , columns i and j .

By induction, use the $k \times k$ minor of A to define $\bigwedge^k \varphi: \bigwedge^k(K_i) \rightarrow \bigwedge^k(K'_i)$

$$e_{i_1} \wedge \dots \wedge e_{i_k} \mapsto \sum_{j_1, \dots, j_k} a_{(L, j)} (f_{j_1} \wedge \dots \wedge f_{j_k}).$$

(And this gives a chain map.)

In particular, the map $\wedge^n \varphi : \wedge^n(K) \rightarrow \wedge^n(K')$ is $\cdot \det(A)$.

If φ is an isomorphism, one would like to say that $\wedge^i \varphi$ is an isomorphism $\forall i$. (And that it gives an isomorphism of complexes.)

Idea: φ is an iso $\Rightarrow \exists a_{ij} \notin \mathfrak{m}$, i.e., a_{ij} is a unit.

know, $a_{ii} \notin \mathfrak{m}$.

We show the isomorphism by changing basis elements one by one.
 $\{e_1, \dots, e_n\} \rightarrow \{f_1, e_2, \dots, e_n\} \rightarrow \dots \rightarrow \{f_1, \dots, f_n\}$.

$$\begin{aligned} \varphi(e_i) &:= a_1 e_1 + \dots + a_n e_n, & a_i \notin \mathfrak{m} \\ \varphi(e_i) &:= e_i & \forall i \geq 2 \end{aligned}$$

$$\wedge^k \varphi(e_{i_1} \wedge \dots \wedge e_{i_k}) = e_{i_1} \wedge \dots \wedge e_{i_k} \quad \text{for } 1 < i_1 < \dots < i_k \leq n.$$

$$\wedge^k \varphi(e_1 \wedge e_{i_2} \wedge \dots \wedge e_{i_k}) = a_1 (e_1 \wedge e_{i_2} \wedge \dots \wedge e_{i_k}) + \text{other terms}$$

Can choose basis s.t. matrix of $\wedge^k \varphi$ is upper triangular, with units on diagonal. Thus, $\wedge^k \varphi$ is invertible.

This will show $K(a_i; R) \cong K(b_i; R)$. □

Remark

Can be extended:

① If $a_1, \dots, a_n \in R$. $b_1 = r_1 a_1 + \dots + r_n a_n$, where $r_i \in U(R)$, then
 $K(a_1, \dots, a_n) \cong K(b_1, a_2, \dots, a_n)$.

② If $a_{n+1} \in \langle a_1, \dots, a_n \rangle$, then

$$K(a_1, \dots, a_{n+1}) \cong K(a_1, \dots, a_n, 0).$$

③ If $\langle a_1, \dots, a_n \rangle = \langle b_1, \dots, b_n \rangle$ (not nec. min'l),

then $K_0(A) \cong K_0(B)$

Q What can we do in the non-local case?

(Notes by Sean Sather-Wigstaff on Kaszdan complexes.)

Lecture 30 (30-03-2021)

30 March 2021 11:32

Let $C. = K.(a_1, \dots, a_{n-1}; M)$, $D. = K.(a_n; R)$.

Recall that we have a l.e.s. of Koszul homology.

$$\begin{array}{ccccccc} \rightarrow H_{k+1}(C. \otimes D.) & \rightarrow & H_k(C.) & \xrightarrow{\cdot(-a_n)} & H_k(C.) & \rightarrow & H_k(C. \otimes D.) \\ & & & & & & \downarrow \\ & & & & \dots \leftarrow & H_{k-1}(C.) & \xleftarrow{\cdot(-a_n)} & H_{k-1}(C.) \end{array}$$

Here, $C. \otimes D. = K.(a; M)$.

Consider the following statements:

- ① a_{n-r+1}, \dots, a_n is M -regular.
- ② $H_i(a; M) = 0 \quad \forall i \geq n-r+1$
- ③ $H_{n-r+1}(a; M) = 0$.

(last r regular
 \Downarrow
 last r homologies are 0.)

Then, ① \Rightarrow ② \Rightarrow ③.

Proof

Induct as before on length of sequence.

If ① is true, $a_{n-r+1}, \dots, a_{n-1}$ is M -regular and hence

$$H_i(C.) = 0 \quad \forall i \geq (n-1) - (r-1) + 1 = n-r+1.$$

Conclude ②. □

Note (Rigidity of Koszul).

Let M be Noe. and $a_1, \dots, a_n \in \text{Jac}(R)$.

If $H_j(a; M) = 0$ for some $j \in \mathbb{N}$, then $H_i(a; M) = 0 \quad \forall i \geq j$.

(Consequence of ③ \Rightarrow ①.)

Let $(R, \mathfrak{m}, \mathbb{k})$ be a Noetherian local ring and M a f.g. R -module. If $I = \langle \underline{a} \rangle$, then

$$\text{depth}(I; M) = \min \{ j \mid H_j(\underline{a}, M) \neq 0 \}.$$

Consequences: ① $\text{depth}(I, M) < \infty$ and is $\leq \mu(I)$.

② Every maximal M -regular sequence in I has the same length.

③ (Depth lemma) Let $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$ be a s.e.s. of f.g. R -modules.

How do the depths compare?

Idea: $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$ gives a s.e.s. of complexes.
 $0 \rightarrow K(\underline{a}; L) \rightarrow K(\underline{a}; M) \rightarrow K(\underline{a}; N) \rightarrow 0.$

Lecture 31 (01-04-2021)

01 April 2021 09:40

① Rigidity of Koszul

② Let R be Noetherian, M.f.g., $a_1, \dots, a_r, b_1, \dots, b_s \in R$ be s.t. a_1, \dots, a_r is M -regular. Then,

$$H_j(\underline{a}, \underline{b}; M) = \begin{cases} 0 & ; j > s \\ \text{ann}_{M/\underline{a}M}(\langle \underline{b} \rangle) & ; j = s \end{cases}$$

Proof Induction on s .

Base case: $s = 1$. We have $H_i(\underline{a}; M) = 0$ for $i \geq 1$.

Hence, the l.e.s. gives $H_i(\underline{a}, b_1; M) = 0$ for $i \geq 2$

and $H_1(\underline{a}, b_1; M) = \ker \left(H_0(\underline{a}; M) \xrightarrow{b_1} H_0(\underline{a}; M) \right) = \text{ann}_{M/\underline{a}M}(b_1)$.

Complete the proof. □

Assume $I \neq M$.

3. Let $a_1, \dots, a_r = \underline{a} \in \langle \underline{b} = b_1, \dots, b_s \rangle = I$ be a max'd M -regular sequence in I .

Then,

$$(*) \quad H_i(\underline{a}, \underline{b}; M) = \begin{cases} 0 & ; i > s \\ \text{ann}_{M/\underline{a}M}(\underline{b}) & ; i = s \end{cases}$$

Note that $\text{ann}_{M/\underline{a}M}(\underline{b}) \neq 0$. Indeed, I consists of zero

divisors of $M/\underline{a}M$ and hence is contained in some $\mathfrak{p} \in \text{Ass}_R(M/\underline{a}M)$ by prime avoidance. Now, $\mathfrak{p} = \text{ann}_R(\bar{\alpha})$ for some $\bar{\alpha} \in M/\underline{a}M$ and $\underline{b} \cdot \bar{\alpha} = 0$ in $M/\underline{a}M$.

Thus, $0 \neq \bar{\alpha} \in \text{ann}_{M/\underline{a}M}(\underline{b})$.

On the other hand, recall that $K(\underline{a}, \underline{b}; M) = K(\underline{0}, \underline{b}; M)$

\swarrow
 r -many b_s
 \parallel
 $K(\underline{b}; M) \otimes K(\mathcal{O})$

$\therefore H_n(\underline{a}, \underline{b}; M) \cong H_n(\underline{b}; M) \otimes \wedge^n R^r$
 (Think of homologies as complex with zero maps)

\swarrow homologies of $K(\mathcal{O})$ is just the Koszul itself

let $n = r + s$.

$H_n(\underline{a}, \underline{b}; M) = H_s(\underline{b}, M) \otimes \wedge^r R^r$ since $H_j(\underline{b}, M) = 0 = \wedge^j R^r$ if $j > s$ or $i > r$.

$H_{n-1}(\underline{a}, \underline{b}; M) \cong H_s(\underline{b}; M) \otimes \wedge^{r-1} R^r \oplus H_{s-1}(\underline{b}; M) \otimes \wedge^r R^r$

⋮

$H_{n-r+1}(\underline{a}, \underline{b}; M) \cong H_s(\underline{b}; M) \otimes \wedge^1 R^r \oplus \dots \oplus H_{s-r+1}(\underline{b}; M) \otimes \wedge^r R^r$
 $= s+1$

$H_s(\underline{a}, \underline{b}; M) \cong H_s(\underline{b}; M) \otimes \wedge^0 R^r \oplus \dots \oplus H_{s-r+1}(\underline{b}; M) \otimes \wedge^{r-1} R^r \oplus H_{s-r}(\underline{b}; M) \otimes \wedge^r R^r$

Comparing with (*) gives $H_j(\underline{b}; M) \begin{cases} = 0 & j > s-r \\ \neq 0 & = s-r \end{cases}$

Also shows $s \neq r$.

Thus, $r = \min \{ j : H_{s-j}(\underline{b}; M) \neq 0 \}$.

\swarrow
independent of \underline{a}

Use the $\wedge^i R^r$ are free and compare from top.

Consequences:

① All maximal M -regular sequences have the same length and

$\text{depth}(\underline{I}; M) = \min \{ j \mid H_{s-j}(\underline{b}, M) \neq 0 \}$,
 where $\underline{I} = \langle \underline{b} \rangle$.

② If $\underline{I} = \langle b_1, \dots, b_s \rangle$, then $\text{depth}(\underline{I}; M) \leq s$.

③ Let $a \in I$ be M -regular.

Then, $\text{depth}(I, M/aM) = \text{depth}(I, M) - 1$.

Idea: $a_1, \dots, a_r \in I$ is a max'l M/aM -reg. seq. in I
 $\Leftrightarrow a, a_1, \dots, a_r \in I$ is a max'l M -reg. seq. in I .

Eg. Let (R, \mathfrak{m}, k) be Noe. local, $\text{depth}(R) = r$, $a_1, \dots, a_r \in \mathfrak{m}$
a max'l R -reg. seq. Then
$$\text{depth}\left(\frac{R}{\langle a_1, \dots, a_r \rangle}\right) = r - r = 0$$

In particular, $\text{depth}\left(\frac{R}{\langle a \rangle}\right) = 0$.

④ (Depth lemma)

Let $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$ be a s.e.s. of f.g. R -modules
and $I \subset R$ an ideal. What can we say about depths?

Idea: If $I = \langle b \rangle$, we have a s.e.s. of complexes
$$0 \rightarrow K_0(b; L) \rightarrow K_0(b; M) \rightarrow K_0(b; N) \rightarrow 0$$

Lecture 32 (05-04-2021)

05 April 2021 10:40

④ (Depth lemma)

let $0 \rightarrow L \rightarrow M \rightarrow N \rightarrow 0$ be a s.e.s. of f.g. R -modules and $I \subset R$ an ideal.

Then,

$$\text{depth}(I, M) \geq \min \{ \text{depth}(I, L), \text{depth}(I, N) \}$$

$$\begin{array}{ccccccc} H_{n-(j-1)}(\underline{b}, N) & \rightarrow & H_{n-j}(\underline{b}, L) & \rightarrow & H_{n-j}(\underline{b}, M) & \rightarrow & H_{n-j}(\underline{b}, N) \\ & & & & & & \downarrow \\ & & & & & & H_{n-(j+1)}(\underline{b}, L). \end{array}$$

$$\text{depth}(I, N) \geq \min \{ \text{depth}(I, M), \text{depth}(I, L) - 1 \}$$

$$\text{depth}(I, L) \geq \min \{ \text{depth}(I, M), \text{depth}(I, N) + 1 \}$$

An observation about the Koszul complex (mainly K_1)

let $\underline{a} = a_1, \dots, a_r \in R$ and $\dots \rightarrow K_2 \xrightarrow{\varphi_2} K_1 \xrightarrow{\varphi_1} K_0 \rightarrow 0$
be $K(\underline{a})$

What is $\ker(\varphi_1)$? $\ker(\varphi_1) = \left\{ \begin{bmatrix} b_1 \\ \vdots \\ b_r \end{bmatrix} \in R^r \cdot [a_1 \ \dots \ a_r] \begin{bmatrix} b_1 \\ \vdots \\ b_r \end{bmatrix} = 0 \right\}$
set of relations on (a_1, \dots, a_r)

$$m(\varphi_2) = \left\langle \underbrace{a_j e_i - a_i e_j}_1 \mid 1 \leq i < j \leq r \right\rangle$$

Koszul relations

$$H_1(\underline{a}) = 0 \Leftrightarrow \text{Koszul relations generate } \ker(\varphi_1)$$

(Hence, being a regular sequence is the closest to linear independence in the ring.)

Suppose $H_1(\underline{a}) = 0$ and $\begin{bmatrix} b_1 \\ \vdots \\ b_r \end{bmatrix} \in \ker(\varphi)$ Then, $\begin{bmatrix} b_1 \\ \vdots \\ b_r \end{bmatrix} \in \ker(\varphi_2)$

$$\Rightarrow \exists C_{ij} \in R \text{ s.t. } \begin{bmatrix} b_1 \\ \vdots \\ b_2 \end{bmatrix} = \sum_{1 \leq i, j \leq r} C_{ij} (a_j e_i - a_i e_j)$$

$$= C \begin{bmatrix} a_1 \\ \vdots \\ a_r \end{bmatrix},$$

where $C \in M_r(R)$ is skew-sym

Thm. (Auslander-Buchsbaum Formula) Let (R, \mathfrak{m}, k) be a Noe local ring, M a f.g. R module. If $\text{pdim}_R(M) < \infty$, then $\text{pdim}_R(M) + \text{depth}(M) = \text{depth}(R)$

[Consequence: In this setup, $\text{depth}(M) \leq \text{depth}(R)$
 Q. Is this true if $\text{pdim}_R(M) = \infty$?

Proof. We prove this by induction on $\text{depth}(R)$.
 Suppose $\text{depth}(R) = 0$ Then, $\mathfrak{m} \in \text{Ass}(R)$ and hence, $\exists a \in R \setminus \{0\}$
 s.t. $\mathfrak{m} = \text{ann}_R(a)$

Let $F. : 0 \rightarrow F_n \rightarrow \dots \rightarrow F_0 \rightarrow 0$ be a min'l free resolution of M over R ($\text{pdim}_R(M) = n$)

Either $n = 0$ or $n \geq 1$ and $\varphi_n(F_n) \subset \mathfrak{m} F_{n-1}$

If $n \geq 1$, then $\varphi_n \begin{bmatrix} a \\ 0 \\ \vdots \\ 0 \end{bmatrix} = 0$ and $\begin{bmatrix} a \\ 0 \\ \vdots \\ 0 \end{bmatrix} \neq 0$. Thus, φ is not injective $\rightarrow \leftarrow$

Hence, $n = 0$ which forces M to be free, i.e., $M = R^{\oplus r}$ for some r
 Hence, $\text{depth}(M) = 0$

(Here $\mathfrak{m} \in \text{Ass}_R(M)$ lets us conclude.)

[Ev. $\text{depth}(M \otimes_R k) \leq \min\{\text{depth}(M), \text{depth}(k)\}$]

$$\left[\begin{array}{l} \text{Ex } \text{depth}(M \otimes N) \leq \min \{ \text{depth}(M), \text{depth}(N) \}. \\ (R, \mathfrak{m}, \mathbb{K}) \text{ local, } M, N \text{ f.g.} \end{array} \right]$$

(Here $\mathfrak{m} \in \text{Ass}_R(M)$
lets us conclude.)

Thus, we are done if $\text{depth}(R) = 0$

Assume $\text{depth}(R) > 0$ We finish the case of $\text{depth}(M) \neq 0$
as follows

Suppose $\text{depth}(M) > 0$ by prime avoidance, \exists a n.z.d a on R and M
since $\mathfrak{m} \notin \text{Ass}(R) \cup \text{Ass}(M)$

Now, first observe that $\text{depth}(M/aM) = \text{depth}(M) - 1$ and
 $\text{depth}(R/aR) = \text{depth}(R) - 1$

Recall: $\text{pdim}_R(M) = \text{pdim}_{R/aR}(M/aM)$ (Use $\mathfrak{m} \rightarrow M \xrightarrow{a} M \rightarrow M/aM \rightarrow 0$)

Thus, by induction, this case is done

Now we look at the case $\text{depth}(R) > 0$ and $\text{depth}(M) = 0$

Note that $\mathfrak{m} \in \text{Ass}_R(M)$ and hence, $\exists x \in M \setminus \{0\}$ s.t.
 $\mathfrak{m} = \text{ann}_R(x)$

Let $\text{pdim}_R(M) = p$ and $\text{depth}(R) = r$ and $a_1, \dots, a_r \in \mathfrak{m}$
be a maximal regular sequence (on R) Then, if
 $\bar{R} = R/\langle a_i \rangle$, we have $\text{depth}(\bar{R}) = 0$

Idea of proof: Compute $\text{Tor}_i^R(M, \bar{R})$ in two different ways.

Since a is regular, we see that $K.(\underline{a})$ resolves $\bar{R} = R/\langle a \rangle$
over R .

Hence, $\text{Tor}_i^R(M, \bar{R}) \cong H_i(\underline{a}, M)$

Thus, $H_i(\underline{a}, M) = \begin{cases} 0 & \text{for } i > r \\ \dots & \dots \end{cases}$

$$\text{Thus, } H_i(\underline{a}, M) = \begin{cases} 0 & \text{for } i > r \\ \text{ann}_M(\underline{a}) & \text{for } i = r \end{cases}$$

Since $\alpha \in \text{ann}_M(\underline{a})$, we see that $\text{Tor}_r^R(M, \bar{R}) \neq 0$ and $\text{Tor}_i^R(M, \bar{R}) = 0 \quad \forall i > r$.

OTOM, since $\text{pdim}_R(M) = p$, we see that $\text{Tor}_i^R(M, \bar{R}) = 0 \quad \forall i > p$.
Hence, $p \geq r$. To show equality, we prove $\text{Tor}_p^R(M, \bar{R}) \neq 0$.

Note that $\text{Tor}_p^R(M, K) \neq 0$ since $\text{pdim}_R(M) = p$.

Now, since $\text{depth}(R) = 0$, $\eta \in \text{Ass}(R) \Rightarrow k \hookrightarrow \bar{R}$ let C be the \mathfrak{m} -kernel. Note that

$$\text{Tor}_{p+i}^R(M, C) = 0 \quad \text{since } \text{pdim}_R(M) = p$$

Hence, the l.e.s of homology after tensoring

$$0 \rightarrow K \rightarrow \bar{R} \rightarrow C \rightarrow 0 \quad \text{with } M \text{ yields}$$

$$0 \rightarrow \text{Tor}_p^R(M, K) \rightarrow \text{Tor}_p^R(M, \bar{R}) \rightarrow \dots$$

Since $\text{Tor}_p^R(M, K) \neq 0$, we get $\text{Tor}_p^R(M, \bar{R}) \neq 0$,
completing the proof □

Obs. Let (R, \mathfrak{m}, K) be Noe. local, M a f.g. R -module.
Recall that $\text{pdim}_R(M) = \sup \{ i \mid \text{Tor}_i^R(M, K) \neq 0 \}$.

$$(\dim_K(\text{Tor}_i^R(M, K)) = \beta_i^R(M).)$$

In the proof above, we showed that $\text{pdim}_R(M) = \sup \{ i \mid \text{Tor}_i^R(M, \bar{R}) \neq 0 \}$

Q Can you generalise this? Can we replace \bar{R} with something else?

Ex. $\text{pdim}_R(M) = \sup \{i : \text{Tor}_i^R(M, N) \neq 0\}$ for any R -module N with $\text{depth}(N) = 0$

Lecture 33 (06-04-2021)

06 April 2021 11:37

Essence of the splitting lemma

① Let (R, \mathfrak{m}, k) be Noe local; F, G be free R -modules of finite rank and $\varphi: F \rightarrow G$ be R -linear

If $\tilde{\varphi}: F/\mathfrak{m}F \rightarrow G/\mathfrak{m}G$ is an isomorphism, then so is φ

② Let R, F, G, φ be as above

If $\tilde{\varphi}: F/\mathfrak{m}F \rightarrow G/\mathfrak{m}G$ is one-one, then φ induces a splitting $\psi: G \rightarrow F$. (In particular, φ is one-one.)

Proof. Let $\{e_1, \dots, e_n\}$ be a basis of F . By hypothesis, $\{\tilde{\varphi}(e_1), \dots, \tilde{\varphi}(e_n)\}$ is lin indep. Extend to a basis $\{\tilde{\varphi}(e_i), \bar{f}_1, \dots, \bar{f}_m\}$

Then, $G = \langle \varphi(e_1), \dots, \varphi(e_n), f_1, \dots, f_m \rangle$ where $\text{rank}(G) = n+m$

Now, we claim that $\{\varphi(e_i), f_j\}$ is a basis of G over R , and this completes the proof. [Define ψ by $\varphi(e_i) \mapsto e_i$ and $f_j \mapsto 0$]

But that is simply true in general \square

Thm

(Serre)

Let (R, \mathfrak{m}, k) be Noe local. If $\mu(\mathfrak{m}) = s$, then

$$\beta_i^R(k) \geq \binom{s}{i} \quad \forall i$$

Proof

Let $\mathfrak{m} = \langle a_1, \dots, a_s \rangle$, $k = k(\underline{a})$ and F be a min'l free resolution of k over R .

Key point k splits off F .

Remaining proof Presentation later \square

This says that k satisfies the Buchsbaum-Eisenbud-Horrocks conjecture.

① Let (R, \mathfrak{m}, k) be a Noe. local ring with $\text{depth}(R) = d$
TFAE

- (i) $\text{pdim}_R(k) < \infty$,
- (ii) $\text{pdim}_R(M) < \infty \quad \forall \text{ f.g. } M$,
- (iii) \mathfrak{m} is generated by a regular sequence,
- (iv) \mathfrak{m} is generated by d elements,
- (v) $\text{pdim}_R(k) = d$.

(Auslander-Buchsbaum-Serre Thm)

Graded objects (and arrows)

Def ① Let R be a ring. We say that R is \mathbb{N}_0 -graded (or just graded) if \exists subgroups R_i of R such that

$$(i) \quad R = \bigoplus_{i \geq 0} R_i \quad \text{as (abelian) groups,}$$

$$(ii) \quad R_i R_j \subset R_{i+j} \quad \text{for all } i, j \in \mathbb{N}_0$$

② An element $a \in R_i$ is said to be homogeneous of degree i

③ An R -module M is graded if $\textcircled{i} M = \bigoplus M_i$ as groups,
 $\textcircled{ii} R_i M_j \subset M_{i+j} \quad \forall i \in \mathbb{N}_0, \forall j \in \mathbb{Z}$.

④ Let M be a graded R -module. We say that the submodule $N \subset M$ is graded (or homogeneous) if it is generated by homogeneous elements

Observations

- (1) $1 \in R_0$ $R_0 \subset R$ is a subring
- (2) Every M_i is an R_0 -module.
- (3) $N \subset M$ graded $\iff N = \bigoplus_{i \in \mathbb{Z}} (N_i \cap M_i)$



$\forall z \in N$ if $z = z_0 + \dots + z_d$ is a homogeneous decomposition,
then $z_i \in N \ \forall i$

Lecture 34 (08-04-2021)

08 April 2021 09:26

② Let $N \subset M$ be a graded submodule. Then, M/N is also graded with $(M/N)_i = M_i/N_i$.

Observe ① Let $\pi: M \rightarrow M/N$ be the natural projection. Then, $\pi(M_i) \subset M_i/N_i = (M/N)_i \forall i$.

We say that π is a "graded map" of "degree 0".

② If $I \subset R$ is homogeneous, then R/I is a graded ring with $(R/I)_j = R_j/I_j$.

③ Let $\varphi: M \rightarrow N$ be R -linear. We say φ is graded of degree i if $\varphi(M_j) \subset N_{i+j} \forall j$.

- Q ① Are compositions of graded maps graded?
 ② Are there non-graded R -linear maps?
 ③ Are $\ker(\varphi)$ and $\text{im}(\varphi)$ graded?

④ Let M, N be graded. What about $M \otimes N$?

It is graded with

$$(M \otimes N)_k = \bigoplus_{i+j=k} (M_i \otimes N_j)$$

Is $\bigoplus_{i \in \mathbb{Z}} (M_i \otimes N_i)$ a graded module?

⑤ $M \oplus N$ is graded

$$\text{as } \bigoplus_{i \in \mathbb{Z}} (M_i \oplus N_i)$$

Examples. ① $R = k[x_1, \dots, x_d]$ is graded with

$$R_0 = k,$$

$$R_n = k \langle x_1^{a_1} \cdots x_d^{a_d} \mid \sum a_i = n \rangle \text{ for } n \geq 1$$

\downarrow
 k -vec. space spanned by all monomials of degree n (in d -variables)

Q What is $\dim_k(R_n)$? $\binom{n+d-1}{d-1}$

The above is called the standard grading

Note that there is always the trivial grading with $R_0 = R$ and $R_n = 0 \forall n > 1$

(2) Let $R = k[x, y]$ Is $x^2 - y^3$ homogeneous?
No (In standard grading)
Can we make it?

Standard ways to make it homogeneous:

① Homogenization Embed $R \hookrightarrow k[x, y, z]$ $f = x^2 - y^3$
 $f_h = x^2 z - y^3$

② Weighted variables Set $\deg(x) = 3, \deg(y) = 2$
Then, f is homogeneous of degree 6